

TOP SECRET  
ULTRA

65/4/7A

I. A description of the machine.

We begin by describing the 'unsteckered enigma'. The machine consists of a box with 26 keys labelled with the letters of the alphabet and 26 bulbs which shine through etecelle on which letters are marked. It also contains wheels whose function will be described later on. When a key is depressed the wheels are made to move in a certain way and a current flows through the wheels to one of the bulbs. ~~XXXXXXXXXXXX~~ The letter which appears over the bulb is ~~called~~ the result of enciphering the letter on the depressed key with the wheels in the position they have when the bulb lights.

To understand the working of the machine it is best to conceive in our minds

The electric circuit of the machine without the wheels.

The circuit through the wheels.

The mechanism for turning the wheels and for describing the positions of the wheels.

The circuit of the machine without the wheels.

Fig 1



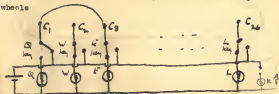
Eintrittszylinder

The machine contains a cylinder called the Eintrittszylinder (E.W.) on which are 26 contacts  $C_1, C_2, \dots, C_{26}$ . The effect of the wheels is to connect these contacts up in pairs, the actual pairings of course depending on the positions of the wheels. On the other side the contacts  $C_1, C_2, \dots, C_{26}$  are connected each to one of the keys. For the moment we will suppose that the order is ~~XXXXXXXXXXXX~~ QWERTZUIOASDFGHJKLZXCVBNM, and we will say that Q is the letter associated with  $C_1$ , W that associated with  $C_2$  etc. This series of letters associated with  $C_1, C_2, \dots, C_{26}$  is called the diagonal, for reasons which will appear in Chap

The particular order we have chosen is known as QWERTZU order.

The diagram shows the connections when the key Q is depressed and supposing that  $C_1$  is connected to  $C_3$  through the wheels.

Fig 2



The only outlet for the positive of the battery is through the Q key of  $C_1$  hence to  $C_3$  and then through the E bulb. The result is that the E bulb lights. More generally we can say

If two contacts  $C, C'$  of the Eintrittswalz are connected through the wheels then the result of enciphering the letter associated with  $C$  is the letter associated with  $C'$ .

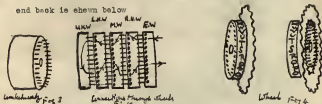
Notice that if P is the result of enciphering G, then G is the result of enciphering P at the same place, also that the result of enciphering G can never be G.

Henceforward we may neglect all of the machine except what affects the connections between the contacts of the E.W., and the turn-over mechanism which affects the positions of the wheels.

#### Connections through the wheels.

The wheels include one which is seldom removed from the machine, and which may or may not be rotatable. It is called the Umkehrwalz (U.K.W.). This wheel has 26 spring contacts which are connected together in pairs. There are three or more other wheels which are removable and rotatable; they have 26 spring contacts on the right end 26 plate contacts on the left (left and right with ~~same~~ positions when in the machine). Each spring

contact is connected to one end only one plate contact. On the wheels are rings or tyres carrying alphabets, and rotatable with respect to the rest of the wheel; more about this under 'turnovers'. When the machine is being used three of the wheels are put in between the U.K.W. and the E.W. in some prescribed order. The way that the current might flow from the E.W. through the wheels and back is shown below



Turnovers. Ringstellung. Window position. rest position.

From the point of view of the legitimate decipherer, the position of the wheels is described by the letters on the tyres ~~which show~~ which show through the three (or 4 if the U.K.W. rotates) windows in the casing of the machine. This sequence of letters we call the 'window position'. When a key is depressed the window position changes, but does not change further <sup>when</sup> the key is allowed to rise. We will say that the position changes into the 'following' position. The position which follows a given one depends only on the order of the wheels and on the original window position. This is because the mechanism for changing the positions is carried on the tyres.

The turning mechanism consists of

Three ~~pells~~ operated by the keys, one lying just to the right of the right hand wheel, one between the R.H.W. and M.W. and one between the M.W. and the L.H.W.

26 catches fixed on ~~the right~~ each wheel on the right.

One (or ~~possibly more~~ possibly more, here we will always assume it is only one) catch on each tyre <sup>on</sup> the left.

The effect of the right hand pell is to move the ~~right~~ R.H.W. forward one place every time a key is depressed. The middle pell

normally comes into contact with the smooth surface of the tyre which prevents of the R.H.W., ~~preventing~~ it from ~~moving~~ engaging with the catches of the M.W. If however it is able to slip in to the catch on the tyre of the R.H.W. it will reach the catch on the M.W. and will push both R.H.W. and M.W. forward: of course the R.H.W. is being pushed forward by the right hand pawl in any case. The occurrence of such a movement of the M.W. is called a 'turnover'. Owing to the fact that the catch is on the tyre the position at which the turnover occurs depends only on what wheel is in the right hand position, and on the window position of that wheel. For instance with German service wheels, wheel I turns over between Q and R, i.e. if I is in the R.H. position then the M.W. will move forward whenever the window position of the R.H.W. changes from Q to R. The left hand pawl operates similarly to the middle pawl, but in this case it is essential to remember that both M.W. and L.H.W. move forward.

Typical examples of consecutive window positions with middle wheel ~~turn~~ turnover E-F, at R.H.W. T.O. Q-R

AWD	BDO	MEW	PEQ
AwP	BDP	NFX	QFM
AWQ	BDQ	NFY	QFS
AXE	BER	NFZ	QFT
AXS	CFS		
AXT	CFT		

Fig 6

The effect of enciphering a letter depends only on the wheel order (Walzenlage) and the position (i.e. amount rotated) of the wheel proper (i.e. not the tyre). To describe this position we could imagine that there was a set of letters attached to the business part of each wheel, and that these letters could ~~also~~ be seen through the windows as well as the letters on the tyres. The letters seen would give the 'absolute' or 'real' position of the wheel (the point of the expression 'real position' will be seen in Chap ). The position of the tyre relative to the business part is fixed by means of a clip on the business part which can drop into holes near the letters. When the clip is in the

hole near the letter C we say that the Ringstellung is C for that wheel. It is clear that some equation of the form

Window position = Rot position + Ringstellung + a constant must hold (it being understood that A,B,C,... are regarded as interchangeable with 1,2,3,...). ~~Normally~~ Normally ~~one arranges~~ one arranges that this constant is zero (see also )

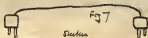
The steckered encipherment.

In some enciphers the association of the contacts of the Eintrittswalz with the keys and bulbs can be varied. There are 26 pairs of sockets labelled with the letters of the alphabet one of each pair leading to a contact of the Eintrittswalz and the other to one of the keys. Normally the two sockets are connected together by a hidden spring, if however a 'Stecker' is plugged into two pairs of sockets, W and R say, these springs are forced away and new connections are made through the Stecker, the W key being connected to the contact which would otherwise be connected to the R key, and vice-versa. That W and R are connected by such a plug is expressed in the form 'W/R' or 'R/W'. The effect of the Stecker on the encipherment is quite simple. If at a certain position of the wheels A enciphered gives N, (abbreviated to AN) then at the same position with Stecker A/V, N/O, and perhaps others, we have VO; if instead we have the Stecker A/V but none involving N, we should have VN (or as we sometimes say the 'connection' VN). Thus if a possible encipherment without any Stecker were

~~XXXXXXXXXXXXXXXXXXXX~~ ~~XXXXXXXXXXXX~~  
~~XXXXXXXXXXXX~~ ~~XXXXXXXXXXXX~~

then a possible encipherment starting from the same positions of the wheels (or as we say, from the same place) ~~xxx~~ with the Stecker B/S, R/N, B/X, V/Y would be

SIEDENKE  
 BVMYBEVO



Each short plug is connected to a long one

### Conventions for electricians

For the purpose of describing the wiring of wheels to electricians one works from a 'spot' on the right hand (spring contact bearing) side of the wheel, or if there is no spot, from the contact which is uppermost when any writing on the face is horizontal



Upper half of wheel.

Fig 8

The contact which is uppermost or nearest to the spot is called 1 and then the numbering is continued in a clockwise direction. One then makes out a scheme like this

Spring contacts	1	2	3	4	5	6	7	8	...
Fixed contacts	6	3	16	14	...				

Fig 9

From the point of view of the cryptograph or the most natural way of naming the contacts is rather different. One would put the Ringstellung to zero, then put zero (Z) in the window, and name any contact on the right of the R.H.W. <sup>after</sup> by the letter associated with the contact of the E.W. which it touches, there being assumed to be no Stecker. To connect these two notations it would be necessary to take into consideration the relative positions of the contact  $C_{24}$  of the E.W. and the windows, and also the positions of the clip and spot on the wheel. Here is a rule of thumb for obtaining electricians data from the cryptographic data, illustrated by Railway Wheel I. W

Write down the first upright of the inverse square for the wheel <sup>unstacked</sup> and above it the diagonal. Use the top two lines to 'transpose' the

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30
Q	W	E	A	T	Z	U	I	C	A	S	J	F	C	A	S	K	Y	X	C	V	E	N	T	L					
Z	E	J	A	B	X	P	W	L	D	I	T	K	C	Y	H	F	R	V	A	S	M	N	I	C					
6	3	16	14	23	19	18	2	26	12	9	6	17	21	19	15	13	4	22	13	1	11	25	24	7	9				

Fig 10

third line into numbers. Then rub out the second and third lines.

This rule is not ~~an~~ absolutely reliable because of possible variations of designs of wheels and machines.

### The comic string.

For demonstration purposes it is best to replace the machine by a paper model. We replace each wheel by a ~~paper strip~~ strip of squared paper 32 squares by 5 squares. The squares in the right hand column of the strip represent the spring contacts of the wheel in natural order (to make the squares of the strip agree with the contacts of the wheel one must wrap the strip round the wheel with the writing on the strip inwards). The squares on the left represent the plate contacts. In the right hand column is written the diagonal twice over, these being the 'cryptographers names' of the contacts as explained in the last section; in the left hand column letters are also written, and in such a way that squares containing the same letter ~~are~~ represent contacts which are the same letter ~~are~~ connected together. Down the centre column may be written the numbers 1,...,26,1,..., 26. These numbers serve to describe the position of the wheel, either the rod position or the window position according to how they are used. The Umkehrwalz is represented by a strip three squares wide, containing in one column the diagonal repeated (this is not entirely essential) in another the numbers 1,...,26 repeated. The third column represents the contacts; and squares representing contacts which are connected contain the same number (which does not exceed 13). The machine itself is represented by a sheet of paper with slots to hold the 'wheels'. In a column on the right is written the unsteckered diagonal ~~times~~ ~~times~~ to represent the Eintrittswalz. It is convenient to repeat ~~this~~ this alphabet between each pair of wheels. The square bearing the letter Q between the R.H.W. and the M.W. will be called R.H.W.'rod point Q' or M.W.'output point Q'. Between the wheels we also write 1,...,26 repeated. These ~~xxxxxxxxxxxx~~ ~~xxxxxxxx~~ are used for describing the position of the wheel when the Ringstellung is given. To understand how this can be done we need only notice that the same effect as a moveable type





could be obtained by having windows and pawls which could be rotated round the wheels in step. To use this Ringstellung device on the comic strips we make pencil marks against the numbers on the fixed sheet and read off the window positions on the strips opposite these marks. We also make permanent lines on the strips to show where the turnover occurs. When these lines pass the Ringstellung merke a turnover occurs.

If the machine has Stecker we may leave a column on the right for the keys to which the contacts of the E.W. are connected through the Stecker.

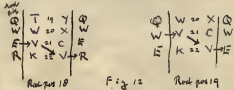
The rule of thumb for the making of comic strips is to take the last upright of the rod square for the left hand columns of the strips.

It may appear rather strange that the letters written on the fixed sheet between the strips should be in the order of the diagonal, rather than say ABCD... ; the point of writing the letters in this order is that wherever a strip is put into the machine there will be the same arrangement of letters on either side of it. If this were not so it would be necessary to have one 'rod square' for the wheel when in the R.H. position and another for the other positions.

Chapter II. Elementary use of code.

The rod square and inverse rod square

It is convenient to have a table giving immediately the effect of a wheel in any position. We can make this out in the form of a square measuring 26x26 small squares, the columns being labelled with the numbers 1,...,26, and the rows labelled with the letters of the diagonal, say qwertz... If we want to know the output letter which is connected to a given rod point we look in the row named after the rod point and the column named after the rod position for the wheel. Thus in column 18 and row e of the purple square we find R, and looking on the fixed comic stripe (Fig 11) where the purple wheel is in rod position 18 we find the rod point E connected to output point R



This square is known as the 'rod square' for the wheel; its rows are known as 'rods' and its columns as 'uprights'.

We can make out a rather similar square in which the rows are named after the output letters and the letters in the squares are the rod points. This is called the inverse square.

It should be noticed that in both squares as one proceeds diagonally from top to bottom and from right to left the letters are in the order of the diagonal. Hence the name. That this must happen is obvious from the fact that if one proceeds steadily round the E.W. as the wheel moves forward one will always be in contact with the same point of the R.H.W. and therefore connected to the same point on the left hand side of the R.H.W. This point is moving steadily round and therefore the rod points describing its position move backward along the diagonal.

Encoding on the rods

For the purpose of decoding without a machine, and in connection with many methods of finding keys it is convenient to have the

V	Y	R	P	S	D	M	T	K	W			q
M	A	Y	T	N	T	C	W	Z	K	O		o

Z	X	F	G	Q	U	Y	R	J	M	P	n	e
J	W	C	A	E	K	G	M	F	Z	U		v

E	V	S	R	P	H	L	G	U	I	C		b
S	B	Z	M	Z	V	E	U	P	A			a

G	M	K	H	L	S	B	J	T	N			k
A	L	U	S	V	G	D	B	I	X	P		g

N	U	L	U	B	R	I	Y	S	O	X		s
P	E	G	L	Y	I	Q	H	D	R			z

X	T	Y	D	F	L	Z	P	E	H			w
R	H	Q	X	O	W	J	F	N	D	T		c

F	Z	X	K	W	J	O	A	B	B	H		u
O	F	N	J	G	M	A	V	H	R			j

H	J	E	O	C	Z	P	Q	X	Q			t
I	Q	I	N	T	O	X	D	A	C			d

L	P	J	Q	D	N	K	Z	M	F			p
C	G	H	W	I	X	T	K	L	Y	L		y

W	O	M	Z	A	C	F	S	V	U			e
D	N	O	C	R	B	R	X	A	T	J		i

Y	K	W	F	M	P	U	L	G	S			y
T	S	B	I	X	E	V	E	Y	L			u

K	R	A	Y	U	Y	W	C	W	P	M		z
U	C	P	E	K	A	S	N	R	J			l

J	U	S	R	S	Q	I	D	B	H	F	N	O	C	G	E	V	A	U	P	Z	N	L	P	Y	T	L
N	V	U	P	R	B	S	T	Y	J	Q	H	I	O	V	W	A	N	S	X	P	X	A	K	C	Z	N

1st 2nd 3rd  
 4th 5th 6th  
 7th 8th 9th  
 10th 11th 12th

Set up of N.W. code for U.K.W. code per 10 L.H.W. (green, 12) 14

12 13 14 15 16 17 18 19 20 21 22 23 24 25 26

12.

PAUGKJCVL... 5  
WVRCSDUPA... 11

NAXHACPOWV... 9  
VKDBHBRNK... 6

EUPAIMVJN... 7  
LPGUQKTVQ... 2

ZRPYGOKDIQ... 6  
DKLXWYUHA... 5

RQCEXASZKQ... 6  
IY SOLBKMSZ 8

OJTPBNUTVX... 1  
ARBITDQGYX... 0

● HCSVGAEB... 5  
KZYNMGAZR... 8

● FSWNYQPIGAR... 6  
GDKFZLYRQ... 6

SIHRKVABDRW... 4  
IOEFVURSLPPY 6

TNOZRWHXCPA... 2  
PFNSJTPPEL... 5

● BTVDPRFIYSUL... 11  
● XDAQNP LKUP... 8

● YH:WCZEAHDK... 11  
● QVZKOUNCHR... 11

UXNLRHUTVVF... 11  
CLQTSIZBYG... 5

Deciphering message 0

First taking up R.H.W. code.

12 13 14 15 16 17 18 19  
Q S Z V  
D E U T

Page 14

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26	
O M T L C Z M X F U L S I H R K V B Q D R W Y J A P	4
T B C D O E S W P N Y G M K F Z L Y R Q I B A H J X	5
E G S T F R X L C J Q Y H I W C Z U A M D K P V U M	6
S C M P D T J F K O V Q W Z K O U N C H B G R A T L	7
A N O Q B J A E A S P U X M L R H U T V Y F C D W I	8
A C P M G Q Z U S V A D K L X W Y O H A T J I R B K	9
K I N A M A T B Y D Z N A U A E J C V L S E Q P H O	10
U W B I L V H O U F A K Z Y N M E G Z R C P D I S Q	11
Y R V X S I W A Q K B T F N J D T N P E L U V O G H	12
H Z Z G T C Q V K W X J O E V U R S L F P Y B I L N	13

P T A E H L P T U R N Z B H Y G O X D I Q C G N V E	14
D F K V X P O Z L I M H C J S V A A E B T Q W T X U	15

M A W N K J R H D Y I C L Q T J I Z B X A V F E O R	16
V K R O A G H K H Z E D T V D P F I Y S U L X F N C	17

N V A B B D E Y M X H L P G U Q X T W O N S J K C Z	18
P U H Z V U A P E C K A R B I T D Q G Y X N O Q M J	19

X Q J E I O A N J G T W V R C S K D U P A Z N Y D B	20
C H I Y S B D O R T O X D A Q N O L K U F S E G Z V	21

J O X K N F L T Z A C F S W N Y Q P I G O R H U E V	22
V Y L H W U I D B H A R Q C E X A T S Z K O T N P S	23

R L Y F Z K G C A B U V U P N I N V J N H T S Z Q B	24
G P E C Y A V Q O L J Y K D B H S R N K W M Z C I P	25

Q X G U T H V S N E R I Y S O L B K N J Z D U W E T	26
I J U B E N Y R V P S T N O Z F V H X C M A W L K G	27

L Z Q V U M C G I Q B O J T P B N W E T E X K S Y A	28
Z S D I U Y K I F O U N G X H A C F O W V H L B R Y	29

Send set up of A. H. W. rolls.

Roll set up - 10

F K S J T T U N Y L Q S Z V I D V M P N E X	19 20 21 22 23 24 25 26
D O H I U F A N N H S Q E T R U P P	

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20	← Think along
A C M R W W X U I Y O T Y N G V V X D Z ...	
C N S I N D J E T Z T A N E K A L A	

rows of the red square written out on actual cardboard rods, in  
guage with squared paper. Let us suppose that we wish to decode a  
~~the following~~ message beginning

~~xxxx~~ QSZVI DMFFN EXACM RWWXU JYUTY NGVVX DZ...

of not more than 30 groups, that we know the wheel order to  
be III I II (Green, Red, Purple), the Ringstellung to be  
26 17 16 13, and the Spruch schlussel to be 10 5 26 1  
i.e. that the ~~windxxxxxxx~~ machine should be set to <sup>the window position</sup> 10 5 26 1  
~~xxxx~~ and the deciphering then begun. We first work out the  
turnovers in terms of red positions. Wheel II has window T.O.  
E-F i.e. 5-6, and since the Ringstellung for this wheel is 13  
the red T.O. is 18-19. The middle wheel window T.O. is M N-O  
and the red T.O. is 24-25. Next we transform the Spruchschlüssel  
10 5 26 1 into red values by subtracting the Ringstellung. We  
obtain 10 14 10 14, and we can now write ~~the xxxxxxxxxxxx~~  
~~xxxxxxxx~~ over the letters of the message the red positions of the  
R.H.W. at which they are to be enciphered, remembering that the  
~~xxxxxxx~~ window position at which the first letter is enciphered  
is not the Spruchschlüssel but its successor. We can also mark in  
the turnovers. Over each section between turnovers we can mark the  
position of the middle wheel. As the message is ~~only~~ not more than  
150 letters no double T.O. will be reached and the U.K.W. will  
be at 10 and the L.H.W. at 14 throughout. We can work out the  
effect of these two wheels for this message once and for all.  
We set up the comic strips for the U.K.W. and L.H.W. to this  
position and read off the pairs of M.W. red points which are  
connected through them. (The fixed comic strips Fig 11 have the  
U.K.W. and M.W. set to this position) They are <sup>q</sup>q, <sup>v</sup>ev, <sup>b</sup>be, <sup>k</sup>ke,  
<sup>s</sup>se, <sup>w</sup>we, <sup>m</sup>mf, <sup>t</sup>td, <sup>p</sup>pr, <sup>f</sup>fi, <sup>y</sup>yu, <sup>x</sup>xi, <sup>b</sup>bw. From these we wish to  
obtain the connections between the right hand wheel red points  
for all relevant positions of the M.W. If we set up the red rods

14	15	16	17	18	19	20	21	22	23	24	25	26	1	2	3
Q	S	Z	V	I	D	W	F	F	N	E	X	A	C	M	...

runner with red points + T.O.s

according to the pairs 40, 47, ... (see Fig 13). In any column of  
 the resulting ~~matrix~~<sup>set-up</sup> will be found the letters of the alphabet in  
 pairs; these pairs are the R.H.W. rod points which are connected  
 together through the U.K.W., L.H.W. and M.W. with the U.K.W. and  
 L.H.W. in the position 10 14 and the M.W. in the position given  
 at the head of the column in question: this can be verified from  
 Fig 11 in the case of column 10. In order to decipher the part ~~for~~  
 the message before the first turnover we set up the purple rods  
 according to the pairs in column 10 of Fig 13. This set of pairs  
 is called the 'coupling of the R.H.W. rods' or simply the  
 'coupling'. The pairs of letters in the various columns of the  
 purple set-up are the possible constations when the U.K.W. ~~is in the~~  
~~matrix~~ ~~position~~ L.H.W., and M.W. have the positions 10 14 10 and the  
 R.H.W. has the positions given at the head of the column. We can  
 therefore use ~~this~~ the set up for decoding up to the first T.O.  
 Afterwards we have to rearrange the rods with the coupling in the  
 11th column of the red rod set-up (FIGS 15 )

Chapter III. Methods for finding the connections of a machine.  
Alphabets and boxes

For any position of the wheels of a machine the letters of the alphabet can be put into 13 pairs so that the result of enciphering one member of a pair is the other member. These pairs are usually written one under the other and called 'the alphabet' at the position in question. Thus the alphabet for the wheel order Green Red Purple and red position 10 14 11 17 is

Y  
 MS  
 VL  
 ZU  
 HY  
 JE  
 TR  
 OG  
 IF  
 XD  
 ED  
 AQ  
 BW  
 NP

The order in which these are written is immaterial.

When we have two alphabets to deal with it is sometimes helpful to describe both alphabets simultaneously in the form of a 'box'. Take for instance the two alphabets

VM	VU
ZJ	ON
ES	JW
GA	HI
KP	TM
XR	FG
OF	EZ
HI	LR
LB	QB
DW	XP
YT	YK
UK	AC
QC	SD

To form a box from these we choose a letter at random, say T, and ~~immediately~~ write it down with its partner in the first alphabet, Y, following it, thus TY; we then look for Y in the second alphabet and find it in YK; we write the K diagonally downwards to the left from Y, thus 

TY	
	K

; now we look for K in



the first and finding it in KU write TY. From this we get to KU  
 TY and TY, but now if we were to continue the process we should  
 KU KU  
 V VM

get TY  
 KU  
 VM  
 TY  
 KU  
 VM  
 TY  
 .  
 .

We therefore draw a line, select a new letter, Hasy, and start again, writing our results below what we have already written. Thus we get

TY  
 KU  
 VM  
 RK  
 PN  
 OF  
 GA  
 CQ  
 HL

Eventually when there are no letters left we stop with the completed 'box' (a box)

TY  
 KU  
 VM  
 RK  
 PN  
 OF  
 GA  
 CQ  
 HL  
 SE  
 ZJ  
 WD  
 HI

There are various remarks to be made about boxes. A box completely determines the alphabet from which it was made. Also it can be written in various forms depending on the choices of letters which are made during the process, but two different boxes made from the same alphabet can always be transformed into one another by a combination of the processes

i) Rearranging the order of the compartments

ii) Moving a number of lines from the top of the compartment to the bottom, the order of the lines remaining the same

iii) Rotating a compartment through  $180^\circ$  about its centre, and then rotating each letter  $\nearrow$  through  $180^\circ$  about its centre.

At first sight it would seem possible that in making a box one might reach a state of affairs like this

AB
CD
E.

and that EA occurs in the first alphabet, and one would not then know what to do. This is not actually possible as EA in the first alphabet would contradict AB. For the same reason it is not possible to have E coupled with any other letter which has already occurred.

If we think of the columns in a compartment of a box we see that the effect of going down the left hand column of a box compartment, or up the right hand column gives the result of enciphering a letter with the first alphabet and then enciphering the result with the second. Consequently if ~~xxxxxxxxxx~~ instead of being given the alphabets we have the result of this double encipherment we shall almost have the box. We shall not know how much to slide the opposite sides of a compartment relative to one another, and in the case of compartments of equal size we shall not know how to pair off the sides.

The effect of enciphering first with  $\alpha$  then with  $\beta$  I shall call 'the permutation  $\beta\alpha$ ', likewise the effect of enciphering with  $\alpha$  then  $\beta$  then  $\gamma$  will be called  $\gamma\beta\alpha$ . For these permutations there is a notation similar to the boxes. However this kind of 'general box' does not enable one to recover the original alphabets. It is also more convenient to write them horizontally (the same applies to ordinary boxes, but the tradition there is firmly established). As an example of the notation

$\gamma\beta\alpha \cdot (\text{QKLAISUHFP})(\text{TCWMBZ})(\text{DKVRN})(\text{J})(\text{O})(\text{Q})$

This means that G enciphered at  $\alpha$  (giving A), and then at  $\beta$  (giving C) and then at  $\gamma$  gives K, likewise K enciphered at  $\alpha$  with  $\gamma\beta$  gives L, P enciphered gives G, and J enciphered gives J. With the same notation the alphabet  $\alpha$  could be expressed in the form (VM)(ZJ)(ES)(GA)(NP)(XR)(OF)(HI)(LB)(DW)(YT)(UK)(QC).

If the letters of a pair of alphabets are subjected to a substitution, and a new box is made up from the resulting alphabets the sizes of the compartments of this box will be the same as in the original box: in fact this box can be obtained from the first box by subjecting it to the same substitution, (except possibly for order of compartments etc.): e.g. if we subject the alphabets  $\alpha, \beta$  to the substitution

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z  
Z D G Y T N B H F I K O L U E M S R Q C J A V X W P

(Z to replace A etc.) then we get the alphabets

AL	AJ	and the box	CW
PI	EU		KY
TQ	IV		AL
BZ	<del>MB</del> <i>MB</i>		RI
KR	NB		MU
EN	TP		EN
HF	OR		BZ
OD	SD		GS
YV	XM		DO
WC	WK		QT
JK	ZO		PI
SG	QY		VY
UM	MX		HF

Conversely if we are given two pairs of alphabets  $\lambda, \mu$  and  $\rho, \sigma$  such that the sizes of the compartments in the  $\lambda, \mu$  box are the same as in the  $\rho, \sigma$  box, then it is possible to find a substitution which will transform  $\lambda$  into  $\rho$  and  $\mu$  into  $\sigma$  (in fact usually a great many such substitutions). We have only to write the boxes in decreasing compartment size(say), and then a substitution with the required property will be the one which transforms letters in corresponding positions into one another.

The sizes of the compartments in a box, and the lengths of the ~~my xxxxx~~ brackets (cycles) are important, as they remain the same if all the letters involved are subjected to the same substitution, (which might be a Steckerling). ~~xxxxxx~~ If we write down the lengths of the cycles of a substitution in decreasing order we obtain what we call the 'class' or the 'shape' of the substitution. e.g. the class of  $\gamma\beta$  above is 11,6,6,1,1,1; with boxes there are two ways of describing the shape, either by the lengths of the compartments or by the numbers of letters in them. It is always obvious enough which is being used. The following information about frequencies of box shapes may be of interest.

26	25%
24,2	13%
22,4	7.3%
20,6	5.4%
18,8	4.5%
16,10	4.0%
14,12	3.9%
22,2,2	3.7%

---

66.8

# The phenomena involved

Before trying to explain the actual methods used in finding the connections of a machine it will be so well to show the kind of phenomena on which the solution depends.

The most important of the phenomena is this, Suppose we are given the alphabets at the positions ~~XXXXXXXX~~ REA FKA WMA and also at REB FKB WMB then there is a substitution which will transform the alphabet REA into REB, FKA into FKB etc. The substitution is that which transforms the letters of the column of the red square corresponding to position A into the letters on the same red square in column B. When we are given complete alphabets we can box REA with FKA and REB with FKB, and the substitution will have to be one which transforms the first box into the second. As an example of this phenomenon we may take the alphabets and boxes

REA	REB	FKA	FKB	WMA	WMB	REA FKA	REB FKB	REA WMA	REB WMB
KX	HO	KH	ZJ	TW	XI	KX	HO	KX	HO
UL	FU	JQ	NP	QD	PG	UL	FU	UL	FU
ED	JM	NL	SU	ZF	HB	ED	KZ	KN	ZE
OD	AG	GC	MA	EN	VE	OD	JM	RT	VI
YV	KL	ZR	HV	VJ	IN	OD	AG	WI	ID
FS	BY	IO	DC	OC	CA	FS	SP	FS	TW
RT	VX	PA	TQ	KL	ZU	RT	NH	YV	KL
QM	PS	BW	WI	CS	MY	RT	VI	JZ	NH
WI	ID	TV	XL	BY	WK	WI	LK	FS	BY
BP	WT	SY	YK	IP	DT	BP	YB	GM	WT
AO	QC	MD	SG	BM	JS	AO	ID	MQ	SP
JZ	NH	EF	RB	AU	QF	JZ	CQ	DC	CA
NK	EZ	UX	FO	LE	OR	NK	TM	QA	CQ

The substitution which will transform REA into REB, FKA into FKB, WMA into WMB, the box FKA into FKB and WMA into WMB is

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z  
Q W A C R B M J D N Z U S E C T P V Y X F L I O K H

In this example the alphabets have been written out in such a way that the letter and the result of applying the substitution occupy corresponding positions. Of course if our alphabets were data from which the substitution was to be found this would not generally be the case. Our problem would be to arrange them in ~~XXXXXXXXXX~~ or the boxes made from them, in such an order.



We make use of a third phenomenon when we have found some parts of the rods. Suppose we find the substitution which transforms the first column of the purple rods into the third

1	34	6
Z	DJ	Y
D	KW	P
G	EC	A
Y	VX	I
T	CD	E
N	FA	D
B	ST	R
H	ZG	O
F	HZ	W
I	UB	N
K	NR	X
O	TL	Z
L	QV	M
U	BI	V
E	OQ	J
M	WN	S
S	MP	T
R	YF	K
Q	GN	H
C	IX	F
J	XX	P
A	OM	Q
V	IG	O
X	ED	O
W	RS	G
B	NS	L

It is

(ZDKNFH)(GEOTCIUBSMWRXYELQ)(JX)(AP)

and the substitution which transforms the <sup>nd</sup> third column into the <sup>th</sup> fourth is

(JYBNSLZWPTRXIVMQ)(CADEOG)(HU)(FK)

These two substitutions are of the same 'shape', and if we write them like this

(YVLQGEOTCIUBSMWR)(NFEZDK)(PA)(JX)  
(JYBNSLZWPTRXIVMQ)(CADEOG)(HU)(FK)

each letter in the lower line is below the letter which is three places further on along the (QWERTZU) diagonal. We can see that this must happen because if we replace the letters of the first and third columns of the rod square by those which are three places further along the diagonal and then move the ~~xxxx~~ result three places to the right and <sup>back</sup> three upwards we get the fourth and sixth columns.

A rather similar phenomenon is useful when we know the diagonal of the machine. In such a case we can make a correction to our constetations transforming them into connections between the ~~constetations~~ on the right of the R.H.W. instead of between constetes of the Eintrittswalz. The constetations when so transformed are described as 'added up' or 'buttoned up'. The process can be carried out with two strips of cardboard with the diagonal written on them, and in one case repeated. As an example to make quite clear what this edding up process is take the fixed comic stripe Fig 11. The alphabet for this position of the machine is (CG)(FR)(TV)(XG)(JK)(WQ)(AG)(FY)(EZ)(HM)(IL)(IM)(US). The added up alphabet can be obtained either by tracing through the wheels from the purple column on the right back to this column again, or by applying the substitution

Q W E R T Z U I O A S D F G H J K P Y X C V B N M L  
Y X C V B N M L Q W E R T Z U I O A S D F G H J K P

to the ordinary alphabet. It is

~~YXCVBNMLQWERTZUIOASDFGHJKP~~

(FR)(TV)(BG)(DQ)(IO)(XY)(WZ)(AS)(EN)(UK)(LP)(CJ)(ME)

Instead of tracing the current through from the right hand purple column in Fig 11 we can of course trace it through from the left hand purple column back to this column again. ~~This is exactly the same as the first method, but with the addition of a slide on the left hand purple column.~~ This gives us a very simple picture of how the added up alphabets between turnovers are related; one is obtained from another simply by a slide on this left hand purple column, i.e. a slide on the ~~fixed~~ last upright of the rod square. For instance ~~the~~ if on the fixed comic stripe Fig 11 we move the R.H.W. to rod position 15 we have the added up alphabet

(EA)(RD)(VM)(IO)(PN)(UB)(LF)(GW)(YX)(CT)(QJ)(KZ)(HS)(

which can be obtained from the added up alphabet at rod position 16 by the substitution

T W V K S B C E Y U F H X Z M N J G O P A Q I R L D  
R L D T W V K S B C E Y U F H X Z M N J G O P A Q T



### The maze

Suppose that one was left alone with an enigma for half an hour, the lid being locked down and the Umkehrwalz not moveable, what data would it be best to take down, and how would one use the data afterwards in order to find out the connections of the machine? Can one in this way find out all about the connections? This problem is unfortunately one which one cannot often apply, but it helps to illustrate other more practical methods.

It is best to occupy most of one's half hour in taking down complete alphabets. At least nine of these are necessary, as follows from this argument. ~~Experiments in writing down all possible solutions~~  
~~find the connections by writing down all possible solutions~~  
~~connections and comparing with the connections of the machine~~  
~~with the data~~ If the solution is completely determined by the data, the number of possible different data must be at least equal to the number of possible different ~~in~~ solutions. Now the number of possible different diagonals is ~~approximately~~ 261, the number of ways in which one can wire up a wheel is also 261, and the number of ways in which one can wire an Umkehrwalz is approximately  $(261)^{\frac{1}{2}}$ , so that the number of possible solutions is about  $(261)^{9/2}$ . The number of possible variations of an alphabet is about  $(261)^{\frac{1}{2}}$ , so that the number of possible variations of nine alphabets is about  $(261)^{9/2}$  which is the number of solutions.

The practical minimum amount of data is surprisingly close to this theoretical minimum. It is possible to find the connections with 9 properly chosen alphabets and 10 other constetions properly chosen. However in order to shorten the work I shall take an example where we are given 11 alphabets and 10 constetations.

Data for page

AAA AAC ABA ABC CAA CAD  
 AAB AAD ABB ACA BAA  
 XXXXXXXX

ACB ADA CAC  
 DAA

GAD

AL AD AI AM AK AE AW AM AS AQ AZ  
 BS BO BY BS BO BS BV BP BO BV BN  
 CE EK OT CH CF CR CZ CR CP CH CO  
 DE EV DM DR DE DQ DX DW DJ DU DF  
 FM GE EV EO OQ FL EJ FG EU EP EI  
 GR HN FN FQ HW GV FO HL FQ FM GL  
 IK IT GK GP IX HK GU IZ GV GM HK  
 JN JY BJ LJ JP IN HI JO HY IZ JR  
 OZ LU JO EK LS JP KR EQ IL JO EP  
 PV OQ KZ LT MY MO LQ NU XT KN MY  
 QW PS LW UZ NR UT MT RS MX RN QV  
 TY RL PQ VY TZ WZ NS TI NR ST ST  
 UX MW RS NW UV TX PY VY WZ XY UW

SO UQ MJ HK  
 ZJ LB IL VS

MA  
 EU

BR  
 TZ

There will be a substitution which transforms AAA into AAB,  
 for finding such a substitution  
 ABA into ABB and ACA into ACB. Following the method explained  
 in the last paragraph we form the boxes AAA, AAB and also AAO  
 which will be needed later

AAA AAB AAO  
 ABA ABB ABC

ACA ACB

CAA CAC

AL AD AI  
 SE BU HU  
 CE EK OT  
 DE EV DM  
 ME VF TC  
 FM GE EV  
 GR HN FN  
 IK IT GK  
 JN JY BJ  
 OZ LU JO  
 PV OQ KZ  
 QW PS LW  
 TY RL PQ  
 UX MW RS  
 YV ZW

AM (SO)  
 BP  
 CE  
 DW  
 EQ  
 IZ  
 JO  
 KQ  
 NU  
 RS  
 TK  
 VY

A S HK  
 BO VS  
 CP  
 DJ  
 EU  
 FQ  
 GV  
 HY  
 IL  
 KT  
 MX  
 NR  
 WZ

AAB

We want to rearrange the box AAB in the way that was done at  
 the bottom of p. The substitution which transforms ABA  
 into ABB must also transform two constetations of ACA into SO  
 and ZJ. The only constetations ~~for~~ of ACA from which SO could  
 have arisen are LH, ~~??~~, VY. If OS arises from ~~within~~ LH we  
 should have to have a substitution which involves ZJ arising  
 from OE in ACA, and this does not exist.

~~A similar objection applies to 48.~~ However if we rearrange it so that OS arises from VY we find ZJ arising from IZ. We can similarly arrange <sup>AAC</sup>ABC to fit with them and agree with GAA and CAC, <sup>A AA</sup> and fit <sup>AAD</sup>GAA to fit onto CAD agreeing with BAA and BAD.

Rearranged			Rearranged	
AAA	AAB	AAC	AAA	AAD
A BA	ABB	ABC	CAA	CAD
AL	VF	GK	AL	UZ
SB	IU	DM	IK	AM
OZ	FJ	TC	TY	IV
TY	RS	ZK	HD	QF
MF	BC	ES	JN	DR
CE	EK	NP	RG	FI
<del>TH</del>	FI	OJ	VP	EO
WQ	NE	EV	CE	CH
GE	KE	BY	UX	IK
NJ	AD	PQ	MF	PG
PV	GO	LN	QW	LT
UX	<del>GW</del>	AI	ZO	SB
<del>IK</del>	<del>ZG</del>	<del>HU</del>	<del>ES</del>	<del>NW</del>

We can now write down the parts of the rods which are in the columns corresponding to the window positions A, B, C, D though we do not know the correct order. They are

AVGT	YSEK	WNEU	UMAV
LFCL	MHRM	QHVZ	KWIY
SLDS	PCSA	GKBT	IZHG
BUMB	GRNF	REYI	ECUP
OYTN	EKFO	NAPE	<del>JDQO</del>
ZICH	DTCC	<del>PLD</del>	
ITZE	HLJH	VOMR	

The substitution which transforms the letters in the first column of these rods into those on the same rods in the second column is (AVOYSLFCREKWN)(BUM)(ZJMFQHI)(GE)

That which transforms the second into the third is (VQUMAPZH)(FK)(LDQ)(YTOWIJCSKBRME)

and that which transforms the third into the fourth (GINFQDZHWGJE)(XLSAVFK)(EUP)(YI)

These three substitutions have now to be arranged one under the other in such a way that the substitution which transforms the third into the second is the same as that which transforms the second into the first, this substitution being a slide of one on the diagonal. Clearly <sup>(FN)</sup> in the <sup>second</sup> has to fit under either



Our set of rods is

TEHG	e
HLJH	h
XWYI	i
EXFQ	x
GKBJ	g
VQWR	v
REYI	r
LCFL	l
KGUP	k
JDQO	j
MBRM	m
OYIN	o
FCBA	f
NAPE	n
PQLD	p
BUMB	b
DTOC	d
CHNF	c
YSKK	y
AVUT	a
ZJGW	z
WNEU	w
SILDS	s
UMAV	u
TPZK	t
QHVZ	q

and we can now transform all our data about other alphabets into this form of data about rod couplings. The ones we need first are

AA AB AC AD  
 eh ar yw fv  
 be bl eh es  
 eu ow bd  
 dt dg px  
 fi ey gt  
 gw fj ki  
 jn ko zo  
 kq ms vl  
 lz nu ns  
 no pt um  
 px qv jg  
 rv rh fe  
 sz iz ar

From these we can get the upright of the middle wheel. The first step is of course to add up the alphabets. Here they are added up with Z as standard

AA\*AB\*AC\*AD\*

pi qj vu hx  
 ol rd dg mb  
 nd sl ye  
 ms to ow  
 kx uk es  
 je ve jh  
 ws zy xf  
 vb op im  
 uh en pq  
 tr bn tn  
 qg wh lr  
 y z fx ze  
 of gi bk

We now box AA\* with AB\* and AB\* with AC\*, and then re-rrange  
 AB\*  
 A C\* so that es to find the substitution which transforms  
 AA\* AB\*  
 A B\* into AC\* and AC\* into AD\*

AA\* AB\* AB\*  
 AB\* A C\* AC\* rearranged

pl	qj	ix	fx	jq
qr	hw	xx	cmx	po
je	et	xf	xf	yz
ve	nb	sk	sk	ed
n d	ku	mh	mh	ig
rt	ve	yx	yx	dr
ol	sl	ks	ks	ls
sw	rd	tx	tx	ev
hu	gi	dx	dx	uk
fx	ne	yx	bw	en
fe	zy	xf	xf	to
ed	op	xx	xx	wo
cm	xi	xf	xi	en

~~This upright of the middle wheel~~ This substitution sends each  
 letter of the upright of the middle wheel into the next on the  
 upright; hence the upright is  
 is leexfrdgpjyxniqohukbmrvao

As we added up to position Z as standard this upright is the  
 upright for position Z. We can make out part of the red square  
 from it, there being  
 from it, difficulties about where to begin as before

ZABCD

LNJHE	z
SNKOL	h
KVHUP	i
WYDOW	x
FMREZ	g
TOLEH	j
KUINE	w
DESIG	s
GAKBV	l
PGOZD	u
JHREK	d
YITVN	m
KOZSU	t
NHADF	e
IFMEZ	n
QTVWO	s
CZERS	e
HLIYAE	p
UFPLI	q
KQUXR	b
BDGYT	o
MTFCA	r
WKNPC	y
VSQJM	f
ABWIX	k
QEOGY	v

We can now transform our remaining data into information about couplings of the middle wheel rods. By sliding the diagonal up the side of the rod square we can get the couplings immediately into added up form

A\* B\* C\* D\* A\* B\* C\* A\* rearranged  
C

ra	aa	ay	kd	ra	es	wl
bt	ba	ba	ox	sl	gz	or
ee	or	el		wj	eq	do
di	do	dr		kg	jk	vf
fo	eq	ez		xv	tx	nb
gk	fv	fa		fo	ph	iu
h y	gz	gs		di	wl	my
fw	hp	hw		un	or	es
le	iu	xp		bt	do	gz
mx	jk	jq		xs	vf	eq
n u	lw	kt		yh	nb	jk
pq	my	mu		pq	iu	tx
vz	tx	vo		ss	my	ph

The left hand wheel upright is

rwdmqreptznschkvbgfiyjouel  
zhixgfwaludmtensapqboryfkv

and under it has been written the diagonal. This serves to transform A or A\* into the Umkehrwala connections. They are  
yv,fs,ee,zw,oi,mu,rj,qx,pk,nd,ht,bg,el

?

'Adding up' method

Most practical methods of finding the connections of the machine depend on getting a long crib, either by 'reading on depth' (see Colonel Tiltman's paper) or by pinching. In many cases we expect the diagonal to have some special value, (e.g. quertzu because the original commercial machine had such a diagonal). In this case the amount of crib necessary is not very much. To estimate the amount of material that we have it is best to work out

(Length - 215) X square of average corrected depth'

Calling this the 'material measure'. By corrected depth we mean the ~~xxxx~~ actual number of constatations, so that this can never exceed 13. As regards the amount of material necessary, it will almost always be impossible to get the wheel out with less than a measure of 90, from 90 to 140 it will be a matter of chance whether it comes out or not. From 140 onwards it will always come out, but with increasing excess the material measure mounts up. With a material measure of 300 it is so easy that the trouble of adding up further material would be more than would be gained in shortening the further work. The method is ~~xxxxxxxx~~ essentially the same as we used for finding the middle wheel in the case of the sage. Here however we have to do with partial alphabets or even single constatations instead of complete alphabets. We cannot therefore do any boxing. After we have added the material up we take some hypothesis about the upright, e.g. that P immediately follows K and work out its consequences. If for instance we find the (added up, I shall ~~not~~ omit to mention this in future) constatations  $\begin{smallmatrix} K \\ 2 \\ K \end{smallmatrix}$  and  $\begin{smallmatrix} P \\ 2 \\ P \end{smallmatrix}$  immediately following one another we can infer that P immediately follows R on the upright. This we may express in the form

~~xxxxxx~~

$$K P - R T$$

the dash denoting logical equivalence. We follow out the consequences until we reach a confirmation or a contradiction. When there is

the K's mean 'P follows K on the upright'.  $K P^2$  would mean 'K + P has no effect on the upright'.



plenty of material we do not usually start to work a hypothesis unless there is going to be an immediate confirmation, e.g. if TC implies RJ from two different parts of the crib. This will mean to say that the constatations  $\overset{R}{R}$  and  $\overset{J}{C}$  occur ~~twice~~ consecutively twice over. Alternatively we can say that  $\overset{J}{R}$  occurs twice over at a certain distance, and that C also occurs twice over at the same distance. In order therefore to find these profitable hypotheses we have only to look for repetitions of constatations (half-bombes as they are rather absurdly called). For this reason and ~~another~~ also because later we will want to be able to spot occurrences of a given letter at a glance, we put our material as we add it up into the form in Fig 19.

Now to take a particular problem. We are given material six deep and 100 long, and we expect that the diagonal is qwertzu. Our material is

MYC..

NGJ..

ROA..

YID..

DAS..

TTV..

YON..

RMI..

OFL..

VQO..

MUX..

NJQ..

(I must apologise for it not making sense) . ~~XXXXXXXXXXXXXXXXXXXX~~

~~LSW.~~~~MTY.~~~~TRF.~~

We decide to try out the hypothesis that there is no T.O. in the first seven columns, and therefore we add up the columns 1-7, 27-33, 53-59, getting

LGN..

MJY..

TRF..

XAH..

FEG..

ZUM..

...

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

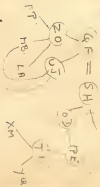
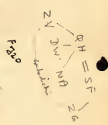
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z



NC  
 (Range from 100 to 1000)  
 with oxygen (1000)  
 (1000)

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100  
 A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

However we put the material directly into the form of Fig 19 . We see numerous half- bombes and do not need to make any analysis of their lengths in order to find a profitable start. The half bombes <sup>Q</sup> S and <sup>F</sup> H suggest the two possible starts  $Q F = SH$  and  $Q H = SF$  (the two strokes meaning a double implication, not equality!). The consequences of the second of these are shown in Fig 20 . A contradiction is quickly reached. The consequences of  $QF$  in Fig 19 . The loop  $QF-ZO-MB-UJ-QF$  gives a second confirmation, and our hypothesis is now a virtual certainty. We now abandon the tree figure for an alphabet with consecutives written against them (FIG 21 ). All goes smoothly except that there is clearly an error in our data as we have a few contradictions. We sort out the good from the bad by using pairs of letters two apart on the upright. Thus  $JO^2 - AF^2$  confirming ~~the~~ JZ,ZO,AQ,QF. When we have checked them all we can write out the upright of the R.H.W.

AQFFEVKYNOCUJZODXMBSHITRWL

We then have to find the upright of the M.W. To do this we use the same process as we did with the sage. We have to find the added up couplings of the middle wheel. This can actually be done without either adding up separately or writing out the rod square, simply by having <sup>two</sup> moveable strips with the upright and qwertzu written out on each, and sliding these above the added up crib till the constatations agree with pairs of letters on the strips directly above. We then read off the coupling from the row of qwertzu letters, taking the pair of letters in column 1 for columns 1-7 of the crib column 2 for 27-33 etc. Under Fig 19 is shown the strips as ~~xxxx~~ set for reading off <sup>one</sup> ~~xxxx~~ of the added up couplings for 53-59, viz  $eq$  . The added up

couplings that we get are

1-7	27-33	53-59	79-85	105-111
qp	hx	qs	fn	zn
wb	qs	wj	xv	ti
ef	wu	eg	tr	
ry	ek	th	fb	
tn	rn	rv	ql	
zu	to	zx	up	
ix	zy	un	oy	
os	is	io	ds	
eg	ov	sk	wb	
dm	dj	db	oi	
hw	fn	fy	gz	
jo	gb	pn	en	
kl	pl	ol	kn	

(some of these being compared obtained from material not yet given)

Boxing these together we get

1-7	27-33	53-59
27-33	53-59	79-85
qp	hx	qs
lk	xy	ks
ef	fm	db
nd	uw	wj
je	jd	np
tn	bg	um
ry	ek	eg
zu	sq	zx
wb	ei	vr
gs	ov	th
ix	rn	fy
hv	pl	oi
os	<u>at</u>	<u>al</u>

---

When we fit these boxes together we fail miserably, and so we have to assume that there is a double F.O. somewhere in spite of the boxes all turning out the same shape. We find that this is between the first and second alphabets, and that the remainder can be fitted together with the upright

wbncovrtirlyazqgpfkmsudj

I will give a second example of the 'adding up' method for a case where it is only just possible to get the problem out. The material is given in Fig 23 all ready added up. There are no 'equidistances' (half-bombas with equal distances) and so we have to make an analysis showing all the consequences of any hypothesis that one letter follows another on the upright (Fig 24). For instance from the analysis we see that ~~AV, HT, NF, ZA~~, are all consequences of IM. The pencil letters round the outside ~~xxx~~ were put in to help with the making of the analysis and were used in connection with columns 32, 33 of the material. Of course some of the consequences will be false owing to turnover, but as we are dealing only with distances of 1 we can hope to neglect this without harm. We now pick out large squares with a large number of entries in them and follow out the further consequences of them, making trees as before, and hoping to find confirmations. When we get contradictions we leave the tree for the present but have to remember the T.O. possibility. <sup>Fig 25-30</sup> When we get stuck we can sometimes continue using consequences which are of the form that two letters are at distance 2 on the upright. For this purpose an analysis of positions at which letters occur is useful (Fig 14). ~~At~~ In particular we need <sup>the</sup> this at Fig 20. Now VW and WY imply  $VY^2$  and PR and RS imply  $PS^2$  and these imply on another from columns 19, 21. We also get  $GL^2$  which starts off another train of consequences involving another confirmation (Fig 11). Eventually we get stuck with the bits of ~~xxxx~~ upright

VWY  
N.Q. PRS  
UHSK  
PGIL.O  
B.E

We might try putting in KA as a hypothesis, ~~Thixxxxx~~ afterwards try KB etc. (KA appears at first to give confirmations, but these are bogus. The only reliable rule about confirmations is to <sup>try to verify</sup> see if one can leave a constitation out and then see if it can be inferred from the hypothesis). We might also try

putting in as many new constatastions as possible which are consequences of those we have and our available information about the upright, and then start off afresh with some new distance on the upright, say 5. But there is a quicker road to success. Note the constation  $\bar{J}$  in 1 and  $\bar{I}$  in 17. Since we have J following H and I following G on the upright it seems highly probable that we have  $HG^{16}$  and  $JI^{16}$ . If this is so we have this as part of the upright

FGIL.O....UHF $\bar{K}$

Hence  $GH^6$  which implies  $FK^6$  giving us this as upright

FGILNQPSRUH $\bar{K}$

From this we get many confirmations and are able to fill in the whole of the upright (except K which goes in the one remaining place)  $N^{0-}$  so that the T.O. which actually occurs between 24 and 25 has not troubled us at all.

1 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23  
 T W G J C N I U J X C S L S K K G N V G W I T S U G I U J A U L J  
 R I D E G I Z W F S N Z E Y H T H U T K R J R R G U O D V T P M I  
 H B E V J V M A O H A G T P D G L O Y S S L A P I M J S W G I U  
 J U S P M M P P Y G T N G H I I K H Y K C L M T R K K S Y G

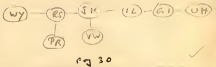
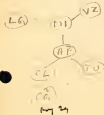
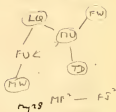
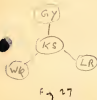
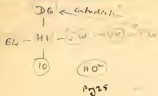
Fr 23

4 -

9

A 5 8 14 20  
 B 4 11 17  
 C 5 6 24  
 D 3 15 22  
 E 3 13 15  
 F 7 9 25 31  
 G 3 11 20 27 30 32  
 H 8 10 16 17 23 33  
 I 4 6 7 18 19 26 27 30 33  
 J 1 4 5 9 22 29 33  
 K 12 15 20 23 29 30  
 L 13 18 25 32 34  
 M 6 7 8 17 32  
 N 6 11 14 19  
 O 9 17  
 P 14 9 16 19 23  
 Q 5 22 34 35  
 R 1 20 23 24 35  
 S 5 10 14 20 22 23 29 31  
 T 1 10 16 16 17 30  
 U 2 8 18 25 26 28 31 33  
 V 6 11 19 29  
 W 5 24 30  
 X 10  
 Y 10 14 20 22 32  
 Z 7 10

Fr 24



WY	JK	GI
RS	VW	UH
PR	IL	

$P_3^2$  — ant

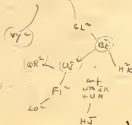


Fig 31



A blank 26x26 grid on a yellow background. The top edge is labeled with numbers 1 through 26, and the left edge is also labeled with numbers 1 through 26. The grid is composed of small squares, with the first row and first column being slightly larger than the others.

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
A						CL	CU	CV																		
B						UN	VS	WT	UX	UY	UZ															
C						BL			QU																	
D																										
E																										
F																										
G																										
H																										
I																										
J																										
K																										
L																										
M																										
N																										
O																										
P																										
Q																										
R																										
S																										
T																										
U																										
V																										
W																										
X																										
Y																										
Z																										

Page 32  
 Program of hypothesis as to successive letters in upright.

Clicks at twenty-six-distances

This is a method for finding the connections when we do not know the diagonal. It is very similar to the ~~xxxxxx~~ beginning of the cage, in principle. It depends on making hypotheses ~~xxxxxx~~ about pairs of letters being on the same rod, and drawing conclusions ~~xxx~~ to the effect that other pairs of letters are on the same rod. Suppose for example that in our crib were the following combinations

E	G	31	32	57	58	83	84
A	E	F	E	T	U	P	U
F	G	T	R	P	R	A	G

We might make the hypothesis that on the rod which has A in column 5 there is G in column 6. We could then infer that there was another rod with F and E in columns 5, 6, and likewise rods TR, PU and this confirms our hypothesis that there was a rod AG. Proceeding in this way we can with sufficient material find sufficiently much of some of the rods to be able to find the diagonal by the cage method. The amount of material needed is very great. We adopt a measure similar to the one for 'adding up' viz

(length-39) X square of average corrected depth

I believe it is practically impossible to solve any problem with this measure less than 2000. It ~~xxxxxxxxxxxx~~ should be possible for 3000 but might sometimes involve a great deal of labour. With the example given here the measure is 4400.

When the material is sufficient we avoid taking hypotheses at random, and choose ones which we can see ~~xxxxxxxxxxxx~~ without very much analysis, to lead to a confirmation. This would be the case for example with these combinations

E	G	31	32
E	E	R	E
V	D	V	D

Either the hypothesis that E follows R or that D follows it on a rod would be immediately confirmed. In the absence of other information the probability that one or other of these

hypotheses is correct is about 79%. Our first job therefore is to look for such configurations of letters. All that we have to do is to analyse the ~~material~~ constations which have ~~existed~~ the same right hand wheel position, and ring round any repetitions. We then write out the ringed constations on a separate sheet (Fig 14). With the first occurrence of each constation we give a number showing how far on the other occurrence is. This ~~material~~ plan also shows us where the T.O. is likely to be. It should be mentioned that in the case of this material there were two turnovers. <sup>known to be 13 apart</sup> The principle of spotting the turnover is this. Consider for example the constations HE at b,II and b,I and JE at i,II and i,I. The first pair of these constations shows that there must have been a ~~coupling~~ pair in common between the coupling at b,II and b,I. Likewise there must be one in common between those at i,II and i,I. It is therefore fairly likely that there is no turnover between b,II and ~~at~~ i,II, as if there had been it would have been quite likely that after the T.O. there would no longer have been a pair in common in the couplings. The evidence from a single such instance is rather slight, but with ~~enough~~ material as we have in our present problem we can fix it ~~with no doubt at all~~, as occurring between x and e and between n and n.

It is worth while writing down all the favourable hypotheses under the pairs of columns of the rod square involved (Fig 34'). We have done this only for the part e to m, and find that in five cases there are two favourable hypotheses viz. col. b with e col. b with h, col. d with j, col. e with i, and col. g with j. We hope that in some of these cases the favourable hypotheses will imply one another, making them both virtually certain. ~~materially~~ ~~in the case of d with j, the constations are the same~~. The consequences of these hypotheses are shown in Figs 34-40. The notation is this. An expression like OF under the head 'd into j' means that the rod with d in col. d has ~~at~~ F in col. j, and the strokes joining these mean that one can be deduced from the other. In the case of g into g the two hypotheses are essentially the same and we have an immediate

confirmation. With  $b \rightarrow h$  we find that both of the first alternatives of the one hypothesis contradict both alternatives of the other. With  $d \rightarrow j$  we manage to connect the two hypotheses together and with  $e \rightarrow i$  we fail to connect but ~~the~~ one of the hypotheses confirms itself. ~~The~~ The information we have obtained about the rods from this is expressed in the Fig. 4. In order to avoid bogus confirmations in what follows it is as well whenever we make a deduction to cross out ~~the~~ one of the constations used in the deduction. ~~Then~~ Up to this point the crossing out has been done with red strokes slanting up to the right. (Green vertical strokes were used to eliminate repetitions of a constation, red vertical strokes to remove contradicted constations.) From now on for a time we will use similarly slanting green strokes.

~~Then~~ Up to now we have simply been trying to 'get a start', and

so long as we could ~~get~~ <sup>most</sup> a fairly considerable bite of the rods square fixed we did not very much care what parts they were. But now we have got a fully adequate start, and we should consider a plan of campaign. In general what we want is ~~the~~ <sup>most</sup> ~~to~~ to have ~~some~~ of the letters of the rods in columns  $p, p+q, p+r, p+q+r, t, t+q, t+r, t+q+r$ , of which any number may coincide, provided  $q, r, t$  are none of them 0. ~~Then~~ If we then find the permutation which transforms col.  $p$  into col.  $p+q$ , expressed in cycles as on p 13 or p 26, and similarly for col.  $p+r$  and col.  $p+q+r$ . A slide of  $r$  on the diagonal will transform these into one another. We get further information about the slide of  $r$  on the diagonal by finding the substitutions that transform col.  $t$  into col.  $t+q$ , and col.  $t+r$  into col.  $t+q+r$ . Between the two sets of information we should have enough to reconstruct the diagonal (unless  $r=13$  and as long as the bits of rod are not too incomplete)

In the present case we can take the columns c,d,f,g,j,k; giving them the numbers 3,4,6,7,10,11 instead of the letters this corresponds to  $p+3$ ,  $q+3$ ,  $r+4$ ,  $s+4$ ,  $t+1$ . In order to get these columns we look on Fig 35 for suitable hypotheses to work in order to add in the extra columns. These hypotheses enable us to write in extra letters in the Fig 41 and we continue to write in letters in this figure until we reach a confirmation ~~xxxxx~~ or a contradiction. Until we reach a confirmation it is as well to differentiate the letters that are certain from the rest. The hypotheses that we actually used were: ~~xxxxx~~ c into g IQ=SE; g into k KE=ND. After a considerable amount of work our ~~xxxxx~~ code look like Fig 46. The lines crossed out are ones that have been amalgamated with others. We now think we can start to look for the diagonal, and therefore make up the permutations transforming c into f, d into g, f into j and g into k. The notation is that of p 1<sup>3</sup>, except that we are mostly unable to complete the brackets, and leave dots.

c into f

...DCYQFVJZTAXHIN...SCOPR...KE...LUB...M...N...

d into g

...EWOM...ANSY...GLIJ...TUQ...HEEXOR...FFEV...H...

f into j

...QOTK...UEWNGR...BSZW...PFA...CXIM...YD...E...L...V...

g into k

...IND... (KE) ...KP...TYRE...MQBLJWURG...PA...C...S...O...V...

We have now to write the c into f permutation over the <sup>d</sup> f into j permutation, and the f into j over the g into k in such a way that the a given letter in ~~wix~~ 'c into <sup>and, in</sup> f into j stands over the same letter in 'd into g' and 'g into k. To get a start on this observe the configuration of the ringed letters. This suggests that we arrange the permutations in this way

DCYQFVJZTAXHIN

HEEXOR

YD  
KE

This is further confirmed many times, and we get the permutations arranged like this

{DCTQFVJZTAJHN} MSGOFR  
{EBKORANSTELJD} KWCMTUQ

{YD} QOTK UELNGROXIMPTA  
{XE} OTYHE IDDMQBLJWURG

giving us the partial diagonal slide of 1

...BCSE...EDNJHK...LXYTOQRF...WMOAV...UP...

Z must be followed either by ~~E~~ E, L, W, or U. If it is followed by U we get

LUB  
FPZV

and the diagonal slide as

{BCSEUPLXYTOQRFEDNJHKWMOAV}

If Z is followed by L we have the bits

MSGOFR KE LUB W  
KWCMTUQ H FPZV

to fit together, which we find ~~impossible~~ impossible, can only be done like this

{KEMSGOFR} {BWUL} {KRWLUEMSGOFR}  
{HKWCMTUQ} {FPZV} or like this {HFPZVKWCMTUQ}

giving the diagonal slides

{EDNJHK}{...}

{UP}...

both of which are impossible. If Z is followed by W we have the bits

MSGOFR KE W LUB  
KWCMTUQ, H FPZV

which fit together only as

{KEMSGOFR} {LUHW}  
{HKWCMTUQ} {VFPZ}

and as before the K configuration makes this impossible. We cannot have Z followed by E because of the impossibility of fitting KE onto <sup>with E on Z</sup> H FPZV. The diagonal slide is therefore

BCSEUPLXYTOQRFEDNJHKWMOAV

46

After the previous examples that have been given it is hardly necessary to explain how to get the uprights of the various wheels after this point. The upright of the right hand wheel would be obtained by rearranging our bits of rod, and the middle wheel by the method described on p. 28. With luck we might find other messages on the same day with different L.H.W. positions and so find the L.H.W. upright. In the case that the Umkehrweiz is movable this may be rather tricky (~~and somewhat dangerous~~), but in such a case there are probably no Stecker, and we should be able to solve other days by single wheel processes, with the known wheels in the R.H.W. position, and hope for the unknown wheels to occur in the M.W. position.

In the example given above the diagonal is actually ABED... with Stecker. We might have ~~even~~ had a hatted fundamental diagonal with Stecker, and of course in such a case we could not have said what the fundamental diagonal was. We should then have had to ~~work our way~~ proceed to try to solve other days keys by spider methods, without diagonal board, and assuming temporarily some arbitrary diagonal <sup>as</sup> fundamental diagonal, and non reciprocal steckering. With two or three such keys we should be able to find the actual fundamental diagonal by comparison of the steckered diagonal.



a b c d e f g h i j k l m n o p q r s t u v w x y z

B (K) T C N E U K (Q) R F Y T H (P) Z Z X P N B P L C I  
 F N L U R T W P (E) N R U N N S T K L K O S L O X N

Q H V M K S I J Y X L G S F V J B P W K Z K  
 R N T R A N S P O R T C H L P S K V O N V O

J J X R H M Y B P N V N Q F N J P U Q W P R B R M  
 K E I N S Z V O N U L D R E V O N V N B S N N

I

F R Z V P Q U T Q S Z S Y N K V A B N G U A J P Y X  
 S Z W O V U A X V I K R X M A U P T A R E Q U E N

L (H) Y F Y B F A U Q B P A D I K P U L G Y X J D C  
 N E H A E M L I N S N U L S K E N S S T R I C H R

I I O B I O H J U K V W W S Y F Z Y T S X N Z  
 I N L A P E S O K U M R G D V U E S T L

II

Fig 33. Material for 'Clocks at 26 distance'.

Letters across top give relative right hand wheel position, and the places for  
 together in order of Roman numerals. Green vertical strokes eliminate  
 mostly confounding combinations. Red ones eliminate misleading ones.  
 Meaning of striking strokes explained in text.

This Fig showed no more further checks

a b c d e f g h i j k l m n o p q r s t u v w x y z  
K N C Q O / Y H E K E R K R H H N D Z A (J) (B) (W) (A) (V) (H)  
Z I N N A / R V I K A U S O U S C H A L T (U) (T) (X) (P) (U) (N)  
Y S / O V Z / (L) H / J D / N A (H) H A G E P B B  
O N / I N S / (N) S B E N / T O (G) S O L A J K I  
P H C G J G / A C (R) V M B (A) / D J S / (L) (I) P L H W Z  
I C H U M V / S T (F) R E M S / U E S / (U) B U A T A A

III

B D / A C U / (D) T D / U I D / O Z K I / (I) T B C (G) N  
K E P A U C / H E U / (B) E R / A R I N / (U) P U N K S T

~~W E X D A G C S V X J~~  
(W) X D U A G C S V X J  
(A) N T R A N S P O R T C

O / A C / Y M (X) V D / A Q W P H / L L A C X / C K V  
N / S F / S N U L N U L L V R / K S D R / (D) K (G)

IV

Q / I A Q / (Y) / Y / M J / (N) X Q / G E M Y  
T / U E D / (F) / F / A H / (S) / K R / H / L T

C (B) G V / A X / (T) N / S O U S / I K (D) I H  
M N O / D / R O / I H / I / V V V M / (O) W A A

a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

~~Q~~ ~~P~~ ~~U~~ ~~B~~ ~~Z~~ ~~U~~ ~~J~~ ~~X~~ ~~Z~~ ~~O~~ ~~V~~ ~~A~~ ~~L~~ ~~E~~ ~~W~~ ~~U~~ ~~V~~ ~~R~~ ~~V~~ ~~L~~ ~~Y~~ ~~X~~ ~~Z~~ ~~V~~ ~~K~~ ~~G~~  
~~K~~ ~~L~~ ~~A~~ ~~E~~ ~~A~~ ~~G~~ ~~S~~ ~~T~~ ~~Z~~ ~~K~~ ~~N~~ ~~X~~ ~~I~~ ~~A~~ ~~R~~ ~~E~~ ~~N~~ ~~B~~ ~~L~~ ~~Z~~ ~~I~~ ~~R~~ ~~K~~ ~~S~~ ~~U~~

$\begin{pmatrix} J & K & L & M & N & O & P & Q & R & S & T & U & V & W & X & Y & Z \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 & 15 & 16 & 17 \end{pmatrix}$

T A H I	B W Q B C O
U U U J	A G E R A N

W X ~~1~~ V ~~2~~ A G J ~~2~~ V ~~1~~ I ~~1~~ O ~~1~~ I ~~1~~ I ~~1~~ S R V ~~0~~ W J  
N U ~~1~~ U ~~1~~ Z F R ~~1~~ R ~~1~~ M ~~1~~ Y ~~1~~ ~~1~~ ~~1~~ ~~1~~ N Z I ~~2~~ ~~1~~

R O I T / X V U Z L N S O N Z T E  
V I K W / T N U L S C H S U L U M R

31 W 2 N 7 P R I N G S Y V T K R V S G D N (A) J K  
5 L L 4 B L 6 1 B T 2 U M V R K W I R M I T (12) L U

G A ~~N~~ ~~V~~ ~~Q~~ ~~S~~ ~~V~~ ~~S~~ ~~E~~ ~~K~~ ~~D~~ A G ~~Y~~ ~~S~~ J H H G L  
 K Z ~~O~~ ~~E~~ ~~N~~ ~~E~~ ~~L~~ ~~N~~ ~~F~~ ~~R~~ C H T T I C A C T P

H R C / A Z G N (2)    Q B T U L A D D C M Y (K) C F  
 T O V / V C H E (2)    U Q A T U L A D D C M Y (K) C F  
                                  N Z U O U U U F A D P

R A Z B C A U L T A V Y C E D V T A R J  
 N S V R V R L T R S O N S T W I E N Z W O

8 7 6 5 4 3 2 1

a b c d e f g h i j k l m n o p q r s t u v w x y z

~~L A Q V~~ ~~L V X W~~ ~~Y K M~~ ~~N~~  
~~G Z U~~ ~~A E B C S~~ ~~E T Z~~ ~~I~~

Z I ~~B~~ N U Y ~~T~~ G I A U ~~E~~ N H P V P ~~I~~ N ~~L~~ Q W ~~S~~ O D Z  
U K ~~N~~ P B R ~~T~~ N S T O P V S O L A ~~I~~ N U S H ~~M~~ I S W

A N Z S ~~P~~ W ~~T~~ Y K G ~~I~~ Z ~~A~~ N B N N ~~O~~ ~~I~~ K D Y C N F  
X S ~~T~~ A ~~N~~ D O R T M ~~L~~ D ~~E~~ N U N ~~A~~ ~~I~~ P A X T X N

~~O L A P N O A V I O E K Q C Y B C A H H C L L S O N~~  
~~M E K H T O A R C P R~~

O L ~~I~~ A N ~~I~~ ~~H~~ ~~I~~ ~~I~~ C X Q C Y B C ~~I~~ H N C ~~Z~~ L ~~O~~  
M O ~~K~~ ~~I~~ ~~O~~ ~~A~~ R ~~I~~ P E W K S T V ~~N~~ V O N A ~~I~~ M

~~L W~~ ~~T~~ ~~X~~ ~~I~~ ~~I~~ ~~I~~

Q ~~H~~ ~~V~~ ~~I~~ ~~S~~ ~~I~~ ~~N~~ ~~D~~ Y P ~~E~~ R A ~~I~~ Y Z ~~I~~ Z Z V ~~I~~ X A V ~~I~~  
Y ~~I~~ ~~S~~ ~~I~~ U L N U ~~L~~ V ~~I~~ R M ~~N~~ B B B Y R ~~I~~

VII

a b c d e f g h i j k l m n o p q r s t u v w x y z

~~A~~ ~~T~~ ~~S~~ ~~G~~ ~~M~~ ~~P~~ ~~Q~~ ~~Z~~ ~~X~~ ~~Y~~ ~~H~~ ~~V~~ ~~S~~ ~~K~~  
~~N~~ ~~O~~ ~~R~~ ~~L~~ ~~C~~ ~~H~~ ~~S~~ ~~E~~ ~~U~~ ~~L~~

~~Q~~ ~~L~~ ~~O~~ ~~T~~ ~~U~~ ~~T~~ ~~Z~~ ~~J~~ ~~B~~ ~~T~~ ~~Z~~ ~~A~~ ~~N~~ ~~H~~ ~~P~~ ~~W~~ ~~D~~ ~~G~~ ~~D~~ ~~Y~~ ~~R~~ ~~I~~  
~~O~~ ~~X~~ ~~K~~ ~~X~~ ~~W~~ ~~W~~ ~~Y~~ ~~P~~ ~~I~~ ~~E~~ ~~A~~ ~~E~~ ~~R~~ ~~K~~ ~~O~~ ~~R~~ ~~P~~ ~~S~~ ~~E~~ ~~I~~ ~~N~~ ~~S~~ ~~N~~

~~W~~ ~~T~~ ~~D~~ ~~S~~ ~~T~~ ~~T~~ ~~C~~ ~~G~~ ~~U~~ ~~I~~ ~~T~~ ~~N~~ ~~J~~ ~~K~~ ~~N~~ ~~B~~ ~~C~~ ~~O~~ ~~Y~~ ~~G~~ ~~M~~  
~~U~~ ~~L~~ ~~L~~ ~~U~~ ~~V~~ ~~E~~ ~~R~~ ~~F~~ ~~N~~ ~~C~~ ~~Y~~ ~~Y~~ ~~E~~ ~~U~~ ~~N~~ ~~I~~ ~~E~~ ~~B~~ ~~E~~

~~V~~ ~~I~~ ~~L~~ ~~E~~ ~~I~~ ~~Z~~ ~~A~~ ~~A~~ ~~O~~ ~~D~~ ~~J~~ ~~S~~ ~~X~~ ~~U~~ ~~O~~ ~~V~~ ~~G~~ ~~A~~ ~~T~~ ~~C~~ ~~M~~ ~~K~~ ~~N~~  
~~W~~ ~~S~~ ~~T~~ ~~K~~ ~~L~~ ~~E~~ ~~S~~ ~~T~~ ~~E~~ ~~X~~ ~~T~~ ~~R~~ ~~G~~ ~~I~~ ~~T~~ ~~D~~ ~~R~~ ~~A~~ ~~N~~ ~~A~~ ~~N~~ ~~D~~

~~P~~ ~~A~~ ~~N~~ ~~I~~ ~~A~~ ~~R~~ ~~N~~ ~~R~~ ~~I~~ ~~U~~ ~~F~~ ~~V~~ ~~P~~ ~~A~~ ~~I~~ ~~V~~ ~~A~~ ~~L~~ ~~K~~ ~~T~~ ~~E~~ ~~B~~  
~~C~~ ~~K~~ ~~U~~ ~~Q~~ ~~X~~ ~~X~~ ~~O~~ ~~I~~ ~~T~~ ~~O~~ ~~Y~~ ~~W~~ ~~O~~ ~~L~~ ~~K~~ ~~E~~ ~~H~~ ~~X~~ ~~E~~ ~~N~~ ~~S~~ ~~U~~

VIII

~~C~~ ~~B~~ ~~Y~~ ~~X~~ ~~D~~ ~~G~~ ~~A~~ ~~N~~ ~~Y~~ ~~G~~ ~~A~~ ~~U~~ ~~E~~ ~~U~~ ~~K~~ ~~T~~ ~~Z~~ ~~A~~ ~~H~~ ~~N~~ ~~A~~ ~~A~~ ~~O~~ ~~J~~  
~~U~~ ~~L~~ ~~N~~ ~~D~~ ~~N~~ ~~O~~ ~~N~~ ~~A~~ ~~R~~ ~~U~~ ~~P~~ ~~E~~ ~~N~~ ~~A~~ ~~R~~ ~~V~~ ~~K~~ ~~X~~ ~~A~~ ~~B~~ ~~W~~ ~~U~~ ~~A~~ ~~P~~

~~Z~~ ~~J~~ ~~K~~ ~~B~~ ~~L~~ ~~K~~ ~~G~~ ~~D~~ ~~Z~~ ~~N~~ ~~H~~ ~~A~~ ~~G~~ ~~P~~ ~~S~~ ~~G~~ ~~W~~ ~~H~~ ~~O~~ ~~G~~ ~~K~~ ~~T~~ ~~S~~  
~~P~~ ~~U~~ ~~P~~ ~~P~~ ~~E~~ ~~W~~ ~~I~~ ~~T~~ ~~V~~ ~~V~~ ~~I~~ ~~T~~ ~~A~~ ~~A~~ ~~E~~ ~~K~~ ~~D~~ ~~T~~ ~~X~~ ~~O~~ ~~S~~ ~~L~~ ~~N~~

~~X~~ ~~C~~ ~~R~~ ~~S~~ ~~E~~ ~~J~~ ~~B~~ ~~O~~ ~~G~~ ~~C~~ ~~C~~ ~~U~~ ~~E~~ ~~A~~ ~~A~~ ~~X~~  
~~N~~ ~~G~~ ~~E~~ ~~H~~ ~~N~~ ~~D~~ ~~A~~ ~~N~~ ~~T~~ ~~W~~ ~~O~~ ~~R~~ ~~T~~

~~A~~ ~~I~~ ~~C~~ ~~A~~ ~~Y~~ ~~A~~ ~~Z~~ ~~V~~ ~~O~~ ~~D~~ ~~H~~ ~~O~~ ~~W~~ ~~A~~ ~~L~~ ~~A~~ ~~F~~ ~~W~~ ~~T~~ ~~U~~ ~~H~~ ~~S~~ ~~E~~  
~~W~~ ~~O~~ ~~V~~ ~~E~~ ~~R~~ ~~X~~

~~S~~ ~~Z~~ ~~W~~ ~~O~~ ~~V~~

~~W~~ ~~A~~ ~~E~~ ~~A~~ ~~Y~~ ~~T~~ ~~E~~ ~~V~~ ~~G~~ ~~A~~ ~~L~~ ~~U~~ ~~S~~ ~~L~~ ~~O~~ ~~H~~ ~~J~~ ~~L~~ ~~T~~ ~~U~~ ~~H~~ ~~L~~  
~~B~~ ~~I~~ ~~S~~ ~~Z~~ ~~W~~ ~~O~~ ~~Y~~ ~~I~~ ~~K~~ ~~E~~ ~~I~~ ~~E~~ ~~R~~ ~~X~~ ~~E~~ ~~N~~ ~~S~~ ~~U~~ ~~E~~ ~~N~~ ~~A~~

~~G~~ ~~A~~ ~~N~~ ~~L~~ ~~E~~ ~~T~~ ~~T~~ ~~C~~ ~~L~~ ~~W~~ ~~D~~ ~~A~~ ~~L~~ ~~G~~ ~~I~~ ~~L~~ ~~A~~ ~~Y~~ ~~T~~ ~~O~~  
~~L~~ ~~N~~ ~~L~~ ~~E~~ ~~T~~ ~~T~~ ~~C~~ ~~L~~ ~~W~~ ~~D~~ ~~A~~ ~~L~~ ~~G~~ ~~I~~ ~~L~~ ~~A~~ ~~Y~~ ~~T~~ ~~O~~

IX

a b c d e f g h i j k l m n o p q r s t u v w x y z  
 L D T Y Z L S T N O L S K M P N I E V T V H Y K  
 N N A R I K W R A E N Z A H R E I C I R K R  
 J T G W N C Y W J T Q X U I A E J W G F Y Y B  
 O R D N O R D W R S T I N S Y U O L K N C I R  
 Y T O L Y U T A R T M U L P N D E H G C Y V A Z Z B  
 L N U L V H A R I N D E S T E N S H U N D F R F S C  
 Y H O X X Z O O Y Z R S I F T N N G I N U D  
 S E I S U L N U L N U L M N T E R Z S C H  
 C Y T O T Y W T G R J W P L Z W K B C Y Y H Z Z  
 U L U L U H A R I T E N G H U D T S  
 V K N G N A L U I A Z H B P O Y P A Q Y  
 S O P F E V O V O B D U U I P U N K P

X

Q Y P R L L K K P T P A Y S U Z K U W A H P D T M N  
 I E A S S H I P R R N D U N T E R B I N D U N G R  
 T O C Q V T X B R P A U A N A P R J P J X L  
 U S Y T A H T O Z W O Y Y H N B E K C H T  
 O N P P N V O Q J A E S T X P N E X V O L B H  
 H U I N S N U L Y R S P A E A G A A N  
 Z Z T S E W L H J Q Z X U A K N  
 N E T R E I O R N Y K B L  
 A D D P C C F U D E J O T Z C W N C  
 C L A U F A F N T I K C H A L T U S E

XI

a b c d e f g h i j k l m n o p q r s t u v w x y z.

V T W N S ~~T~~ R S N X ~~X~~ N X ~~H~~ R C J H K Q N W B D T L  
U K C K W ~~N~~ R T I A E R V E R B I N D U N A K K E

J H G J D ~~R~~ Y Q Q E F V W N W S A H K H G I C P B B  
I H O K H E ~~N~~ E R B I S F U N I N U L N L N U L N

~~E G K H T L E K A A E G N U E S N G V H Z J J V V~~  
~~F E N P A T A C N E N L O C I Y Y D R E I Y S E O H~~

A J D S ~~V~~ I C Y H H R A ~~P~~ K K W B H D O M C  
R M A ~~L~~ T U N A B Z W O N F X V I C R X N U

~~F~~ ~~F~~ ~~F~~ ~~K~~ ~~P~~ ~~L~~ ~~K~~ ~~A~~ ~~A~~ ~~B~~ ~~N~~ ~~U~~ ~~S~~ ~~M~~ ~~B~~ ~~V~~ ~~Z~~ ~~S~~ ~~J~~ ~~V~~ ~~V~~  
~~P~~ ~~A~~ ~~T~~ ~~O~~ ~~N~~ ~~E~~ ~~L~~ ~~C~~ ~~I~~ ~~Y~~ ~~Y~~ ~~D~~ ~~R~~ ~~I~~ ~~Y~~ ~~S~~ ~~E~~ ~~C~~

XII

F R ~~K~~ ~~K~~ U M B D N S A C K L U S R D F O G Q R P D.  
R Z ~~A~~ ~~H~~ N K K N I C H T M E H R N O F A L I E H X

J S M J A L M G A U V H V A G Q C K Y Z X R  
T R K P L S E I N D R E I Y Z W O X B A R O M

I C R ~~M~~ ~~R~~ ~~P~~ ~~D~~ ~~Y~~ ~~J~~ ~~X~~ ~~L~~ ~~O~~ ~~Q~~ ~~K~~ ~~U~~ ~~Z~~ ~~N~~ ~~H~~ ~~V~~ ~~P~~ ~~I~~ ~~L~~ ~~Z~~  
S K ~~L~~ ~~M~~ ~~R~~ ~~L~~ ~~Z~~ ~~E~~ ~~T~~ ~~K~~ ~~K~~ ~~S~~ ~~O~~ ~~F~~ ~~I~~ ~~X~~ ~~K~~ ~~A~~ ~~A~~ ~~X~~ ~~L~~ ~~U~~ ~~C~~

XIII

~~E~~ ~~Y~~ ~~P~~ ~~F~~ ~~S~~ ~~A~~ ~~O~~ ~~E~~ ~~B~~ ~~C~~ ~~A~~ ~~B~~ ~~V~~ ~~I~~ ~~V~~ ~~A~~ ~~E~~ ~~L~~ ~~Y~~ ~~C~~ ~~F~~ ~~B~~ ~~K~~ ~~W~~ ~~V~~  
T P J S ~~E~~ ~~C~~ ~~H~~ ~~I~~ ~~W~~ ~~F~~ ~~C~~ ~~L~~ ~~Y~~ ~~C~~ ~~J~~ ~~D~~ ~~K~~ ~~V~~ ~~V~~ ~~V~~  
L N U ~~A~~ ~~V~~ ~~L~~ ~~L~~ ~~L~~ ~~U~~ ~~H~~ ~~R~~ ~~D~~ ~~I~~ ~~L~~ ~~E~~ ~~I~~ ~~T~~ ~~N~~ ~~M~~

~~a b c d e f g h i j k l m n o p q r s t u v w x y z~~

~~x v s q o n y k e x r h t n j z a j b n t v h b d r~~

~~z i v n a r v i k a u s a v s c h a l f o t~~

v

c



a b c d e f g h i j k l m n o p q r s t u v w x y z

(C) Y Y T R V V B R U A X O P S G R T X J B U P A N Z  
 U N D T L U J A U W E R P E N T J U L V K G A A D S

J A A Y C Y D H T Y  
 T T Z U H R T X K L

X J R S J T P T U Z T Z Y P P T L A D D H  
 N P U R A R J J T T S A C K G U R U N L

J J S G R D I D X Y Q S K Z U E R I D T V I A  
 U L T X I K N O C T L T U P U N D N L T T N

XIV

V C N P B B S Z U L U C A I K B O V P P D H G X P  
 H I A R H R B R S T I N N T S G R R A T I N H

N J Y R Y N U H Z P J U U P A V K J X Y T V U O  
 E J S N U L U T N S U E U N Y W O N H B S I C

P Z D D G B K O W Y F X T A N R W S X A O T  
 X Y T V I R N U L J L U C I Y Y V I E R P U

K J O J P P W X K Y K I G G I Q Y V H V  
 N R U U N S L T I S S Z U O N U L L

XIV



A B C D E F G H I J K L M

A  
B  
C  
D  
E  
F  
G  
H  
I  
J  
K  
L  
M  
N  
O  
P  
Q  
R  
S  
T  
U  
V  
W  
X  
Y  
Z

A B C D E F G H I J K L M

F-735

d into j

$$NF = OZ - XV$$

$$NZ = OP$$

$$KF = VZ$$

$$KF - KZ = VZ$$

$$LA = SN - RH$$

$$LN = SA - AT$$

Page 36

d into j

$$PO - TX$$

$$CS - YU - EM - UG - DR$$

$$XS - MW$$

$$KF = VZ - SN = LA$$

$$RH$$

Page 37

e into i

$$LE = LF - PM$$

$$IS = LE - EV - AL$$

$$TN - NT = OW$$

$\begin{array}{c} ?x \\ | \\ TN - NT = OW \\ \swarrow \quad \searrow \\ SY \quad VC \\ | \\ XG \end{array}$

$$NV = OT - AC$$

$\begin{array}{c} LE \\ | \\ NV = OT - AC \\ \swarrow \quad \searrow \\ XY \quad QZ \\ \quad \quad | \\ \quad \quad 27 \end{array}$

Page 38

A blank 26x26 grid with numbers 1-26 along the top and left edges. The grid is used for graphing or plotting data.

b into h

5-8

$$LO - 4R = ET - KG - NS - 1A$$

$$14T = ER - \sqrt{N}$$

$$AG - 2H - RP - \sqrt{K} = U1$$

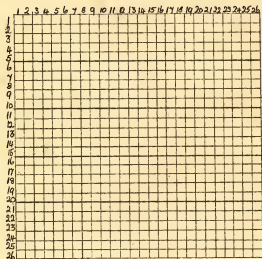
$$\sqrt{1} = VK - XX - 1-R$$

Fig 39

q into j

$$QA \xrightarrow{\sqrt{}} ER \xrightarrow{\sqrt{}} DE \xrightarrow{\sqrt{}} CY$$

Fig 40







### Finding new wheels. Steeper knock-out

So far we have been dealing with the problem of getting out the connections of an entirely new machine, or one for which we know no more than the diagonal. There is another problem, that of finding the connections of some newly introduced wheels, the old wheels, or at any rate some of them, remaining as well; this includes the case of a change of Umschaltungs.

The most hopeful case for getting out the new wheels is when ~~ix~~ one of the known wheels occurs in the R.H.W. position. If the machine has no Stecker there is no difficulty. We solve some messages by single wheel processes. This will be slightly more difficult than when we know the connections of the middle wheel, as we shall have to guess what is said in <sup>three or four</sup> ~~its~~ different turnovers. However ~~xxxxxxxxxxxx~~ when the R.H.W. rod start has been found from a guess in one turnover it does not take any time to test a not probable throughout the message (the rods on which the various letters of the message occur can be written down once for all, and the not probable punched out and run over the inverse oblong). For simplicity let us suppose that we have read the message right through. We then have the couplings in several ~~many~~ consecutive positions of the middle wheel, and can apply the method of p 28, 29 to find its upright.

In the case that the machine has Stecker we need rather more  
data, and very much more patience. The sort of data that one needs is  
a crib of length about 70, or else one of length <sup>26</sup> ~~18~~ and depth 2. The  
trouble about cribs without any depth is that one uses up the  
the great many of the constetations ~~after~~ between each turnover in  
determining the couplings.

An example is shown of ~~xxxx~~ a crib of length 18 and depth 2. This is to be regarded as one of ~~lengths 26 which can be made from~~ greater length which has been cut down to allow for turnover. The text of the crib is shown at the top of Fig 41. ~~The xxxxx~~  
~~xxxxxxxxxxxxxxxxxxxx~~ We are taking the worst case of 13 Stecker. There are several half-bombes in the crib, and we decide to work with TW. We have to make ~~xxxxx~~ 17576 different hypotheses, (epp) corresponding to the 26 possible different places on the R.H.W

ANTRANSPORTCHASSE  
FGNYFZJWLOWDUDLMHD

L I S T Y W E W A V O N U E W A Z E I N  
T A D J S B U P U L C H A D T E F D K M.

A	(F)	(F)	U	W	(E)	AG 3	✓
B	OW	(W)		OD		BT	✓
C				C		C	
D	S	U		F	(A)(H)	D	✓
E	(A)	(A)		L	Z	EW	
F	(N)			U	(S)	FOS	
G	A		O			GA 3	
H	T	(S)				HV 9	
I						I	
J						JL 11	
K	T		V	P	(S)	KM 6	
L	(G)(T)	Z		(N)(M)		LJ 11	
M			IR C			MK 6	
N						NS 4	
O						OF 8	
P		W				P	
Q						Q	
R						R	
S	Y	(J)	O		(M)	SN 4	
T	(N)(J)		(W)	(W)		TB 1	← ✓
U						U	
V		E A	H			VH 9	
W	(B)	P	T A	(I)		WE 1	← ✓
X						X	
Y	RS					Y	
Z		N			F	Z	

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18

MSFUOGLBFH(E)YITKRKMNNWBEFHEL 0  
ZACZGHNAPTJ(S)GWACZTDEFMFHIRW. 5

LQGKMTJDNOPWPRB(S)GJTKMJQRXKZT 1  
ACDMRUVXUBCIVKEFWBRVXUOYZAH 2

RHLNKEOPQXQSTCHKLNKRSLAUM 4  
QNHIRSTATVWFKNOQNUVBODXPUKO 2

GZBCLATUWATABHUIDVAQUWNTNXYZ 5  
OISTUBUWXGLOPROVNCPEYQVLP 4

Fig. 42. Inverse alphabet used by the German in the Enigma machine.

and the possible different 'Stecker values' of T and W. Any assumption as to the Stecker values of T and W ~~implies~~ implies two rod pairings, and when we have set these rods up we can look round and see if there are any other Stecker which are consequences of the rod pairings and the Stecker we have already. Any new Stecker we find may allow us to set up more pairs of rods. So we go on until either no new consequences can be drawn (this may be rather frequently the case), or there is a contradiction. If there is confirmation and afterwards we can draw no ~~more~~ further consequences it may be worth while bringing in extra hypotheses.

In the actual working it seems best to set the crib out as in Fig 41, so that the occurrence of any letter can be spotted at once. We write the Stecker values of the letters in pencil on the right, possibly on a separate sheet which slips underneath. In order to avoid bogus ~~constatations~~ <sup>for want</sup> constataions we cover up the constataions with shirt buttons as they are used. Fig 41 shows the working for the correct hypothesis W/E,T/B. The 'covered' letters are shown ringed. In order to show how the working was done the steps have been numbered, the number being put against the constataion used and also against the Stecker values or rod pairing which resulted. The work as shown is not quite complete. It is possible to go further and get the Stecker values of all letters except D,X. There are six or more confirmations.

There are a number of other possibilities besides working from a half-bombe. It depends largely on the number of Stecker expected which will be the most profitable. When the number of Stecker is low (say 6) it is probably best to try half-bombes as unsteckered and to look for clicks which have all four letters unsteckered.

It seems unlikely that this method will ever be applied, partly because of the difficulty of obtaining the right kind of data.

However, ~~the~~ <sup>to find a new Unkehr rule</sup> much the same method could be applied with data of the kind that arises with the air enigma. ~~Thereon~~ <sup>Thereon</sup> may find the Ringstellung by Herivelismus, and also have a certain number of constataions at known window positions arising from CILLI e

The wheel order may also be known from CILLIs more or less accurately. ~~if~~ We now make up rods giving, not the effect of going through the R.H.W. but through all three wheels, and with the columns not corresponding to ~~all~~ all possible positions, but to the positions where there are known constetations, and use them instead of the ordinary rods: there is no difficulty about T.O.

Identification of wheels

When one has found the connections of a wheel one naturally wants to verify that it is not one of the wheels used in some other known machine. A convenient way of doing this is to find the class of substitutions which transforms one column of the red square into the next (see p 194). Thus the class of the wheel found on p 26 was 13,8,3,2. ~~Further examination of this~~ This 'class' is independent of what point of the wheel square we take to be the top left hand corner, and so is an absolute characteristic of the wheel. It even remains the same if the wheel is used in a machine with a different diagonal. In the case of an Umkehrwheel we can form the class of the substitution consisting of going through the U.K.W. and then sliding one position backwards on the diagonal. A list of characteristics for the known machines is given below

Messiaen machine

- I. 19,7
- II. 14,12
- III. 10,8,5,3
- U.K.W. 15,9,1,1
- Service machine
- I. 13,6,4,3
- II. 16,10
- III. 7,7,6,6
- IV. 11,11,8,2
- V. 9,9,6,2
- VI. 24,2 two apart 9,8,6,3
- VII. 12,5,5,4
- VIII. 24,2 two apart 22,4
- U.K.W. A. 9,8,4,2,2,1
- B. 10,7,8,1
- C. 13,8,2,2

Railway machine

- I. 24,2 two apart 18,5,2,1
- II. 12,8,4,2
- III. 14,8,3,1
- U.K.W. 24,2
- Commercial
- I. 18,8
- II. 19,7
- III. 12,9,4,1
- U.K.W. 22,2,1,1



MANHESIM . ~~XXXXXXXXXXXXXXXXXXXX~~ In order to decode more of the message we ~~XXXXXXXXXXXX~~ can either try using the three couplings after the turnover to read a little more. This is shown in Fig 44. It is not possible to fill in the intermediate letters and we have to find some other method. One is to try decoding after the T.O. with various assumptions about ~~the~~ which wheel is in the middle position, and what rod position the M.W. is in. We shall not actually need to do the decoding for each such position, as a very large proportion of the possibilities is immediately eliminated by the ~~the~~ ~~known~~ known to occur after the T.O. In fact we have the seven couplings ku,so,fa,yn,ey,td,vh before the T.O. and ~~the~~ ~~known~~ ~~the~~ ~~ca~~, la after it ~~and possibly two more~~. We ~~might~~ could treat these couplings with respect to the middle wheel in the same way as we treated the original crib with respect to the right hand wheel. However it is not really necessary to get out the rods. It is easiest to work with the rod squares and for each possible position of the middle wheel look and see what coupling before the T.O. is a consequence of ca after the T.O. For example there are the bits of red red

12  
EA  
VO

and therefore if the message starts in rod position 1 for the middle wheel the coupling mv must have occurred before the T.O. in order that ca may occur after it. Consequently this <sup>(above)</sup> position for the middle wheel is impossible. That the middle wheel rods can be used in this way amounts to nothing more than that they can be used in decoding in the way described on p. 14,15. In this way we find that the only possible positions for the middle wheel <sup>is</sup> ~~if we find~~ ~~XXXXXXXXXXXX~~ ~~XXXX~~ ~~XXXXXXXXXXXX~~ red 11, and we have for couplings after the T.O. yg,uv,kt,wh,ws,em,el,oa ~~XXXXXXXXXXXX~~ ~~XXXXXXXXXXXX~~ ~~XXXXXXXXXXXX~~ and the part of the message from the first to the second T.O. reads

VKOUZK RZCOOV FKVLDZNNRDS  
EIM.CAL.A.MEETOTER.IT...E.

We can fill this in to read, for the whole message up to this point  
~~XXXXXXXXXXXXXXXXXXXX~~ DANZIGVONDIANNHEIDKIDANZARMAKETOTERBITTEREFEL.  
 The other couplings ~~xx~~ rf,jz,qi can now be read off the filled  
 altogether we now have  
 in letters, and ~~xxxxxxxx~~ the couplings of the M.W. rods  
 qc,er,eb,ex,ws,jm,pz,ff,yd,zl,hb. We can decode as described in  
 Chap II; the two remaining middle wheel couplings will soon be  
 found.

We might of course use either the middle wheel couplings or  
 the right hand wheel couplings to find the position of the  
 L.H.W. and U.K.W. and we could then do the decoding on a machine  
 instead of on the rods. Methods for doing this will be described in  
 the next chapter. The rest of this chapter will be devoted to  
 methods of brightening up the ~~xxx~~ first parts of the process.

The inverse rods

Instead of picking out the R.H.W. rods and laying them against  
 the crib as in Figs 43,44 we might write down the ~~xxxxxxxx~~ rod  
 couplings which are consequences of each of the constetations, thus  
 when testing pre-start 26

FKSJTQMY  
 DANZIGVON  
 omuq<sup>ikie</sup>  
 wjeom<sup>xy</sup>

The contradiction which we found before by setting up the pair ~~ow~~  
 now shows itself in the form of two contradictory couplings  
 ow,oq. In the case of pre-start 4 we have

FKSJTQMY  
 DANZIGVON  
 upl<sup>am</sup>  
 ked<sup>winkie</sup>

and our confirmations (clicks) show up as repetitions of the  
 couplings uk,qf. If we actually did this we should lose time  
 in comparison with the original process, but we can actually get  
 all the couplings in the different positions by a more mechanical  
 method.

We have the lines of the inverse square (p. 10) written out  
 on rods in double length, called 'inverse rods'. We



Handwritten note at bottom right corner.

...C N K P S D C A L T A N Q A I X H T P Y E O V Z T U C N Q S G V K E T A N L J U Y G H K P M I Q T K Z V D J S F V K E T A N L J U Y G H K P  
...N K D I S L V K E T A N Q A I X H T P Y E O V Z T U C N Q S G V K E T A N L J U Y G H K P

Q T R O G E T A N Q A I X H T P Y E O V Z T U C N Q S G V K E T A N L J U Y G H K P  
...K T C G O S K P Q B W U V N A T K J C V J L U C I P Q C I R O S E T K D N Y U V N A T K J C

Y V H E D C Y W X I B T S Z P E G O K Q U L E J T Y V A D C Y U V X E I D T S Z P E G O S  
...N Y T L Q A R G X F K E P Z O U D H I Z K Q U L E J T Y V A D C Y U V X E I D T S Z P E G O S

N U E I F W A T T S B P Y W V Z A T O C A N Q K U K L N U E I F W A R J S B P Y U V Z A T O  
...C H S L G K A T U T G D Y N E P X Z G F I O P Z V U C M S L A K R T U T G B Y N E P X

...M F J O R Z H P V T B W K Y T L D U I K S C Q K Z A T F T O R Z H P V T B W K Y  
...Y D U D P J M C L T X V I G H A S K S N A B A P U C Y D V O P J T C L T X V I G H

E T F S O A Z H F V P O U K Y T L J U I K S C Q K Z A T F T O R Z H P V T B W K Y  
...T G Z S L U G O C T K L X A R E U Y I V N P P H K V T Q Z S L U G O C T K L X A R

K S T X Z C O D T L W S O P G Y N F H G V I E R F U K O S X Z C O D T L W S O P G  
...N H I M P S H Q J K U C C B Z A P G Y I A V K E Z N I U D P S H Q J K U C

D G A C K E P Z B W B T I Z H S J C A D N Y T L Q A R G K P K E P Z B W B T I  
...T Y K L V A Z C I G O R J S D C H N T W V A E K I P I A Y K L V A Z C I G O R J S

L C V T H O I D X F G R B T X C P A Z S T T U Y G L C V T H O I D X F G  
...G X P K E P Z O W B M I Z H S J C A D N Y T L Q A R G K P K E P Z O W B M I

Fig. 4. Strip of Purple Paper

PKS5RPA X Y  
DANZ, SAVON

P K S J T Q N Y L B S C K V K X

D A N Z I G V O N

A U H E V W B P L C K R R B I T D Q Q Y X N O Q M J

B F K V T P O Z L I M H C J S V G A E B J Q N T K U

Fig 43. Tachy pre-start 26

P K S J T Q N Y L B S C K V K X

D A N Z I G V O N A N N H N

M B D T J F X O V Q V Z K I O U N C H S G R A T L K

S T P R X L C J Q Y H I W C Z E A N D K P V U M U

K C Y A U Q O L G V K B B H S R N K W P Z C I R C

P Z K G C A B V F U P A I M V S N H T S Z O J P

Q V U P C G I Q B O J T P E M V F T E X K S Y A f

J Q Y K I T P U N G X H A C F O W V H L B R V g

I Y S B D O A T O X Z A W N P L K V F I E G Z W X

T E I O A N F G T U V R C J K B U T A Z O Y B B A

U B E N Y R V P S T K O Z F W H Z C P A W L K G a

G U P H V S N C P I Y S O L B K P J Z J V W F T y

C D O K J U P N Y G P K F Z L Y R Q I B A H S X b

T L C Z K R F U L S I H R K V B Q D R W Y J A P d

S H L P J U R N Z B F Y G O X Z I Q C G M V C l

R W X P O Z L I M H C J S V G A E B J Q N T K U w

Fig 44.  
Pre-start 4

V K X U Z H E R J Z O Q V E T K V L D K S N R D B S

E I M G E O E R

Fig 45. Comparison of 470 after T.O.



pick out the xxix inverse rods named after the letters in the crib, and lay them down in pairs, staggering them backwards. This is best seen in Fig 46. The various columns in this set-up show us the various rod couplings which are consequences of the crib and various hypotheses as to the pre-start. In the figure the pre-starts have been written along the top, but this is not part of the normal routine. With this method we can easily see contradictions which are independent of where the T.O. occurs e.g. for pre-start 1 we have the couplings w1, w1, j1 arising from the crib in that order. There must be a T.O. between the w1 end and the w1 end and also between the w1 end and j1, which, apart from double T.O. is impossible.

## Make

There is another ~~Kasiski~~ method which gives essentially the same result as the <sup>and</sup> inverse rods and seems to be a little quicker~~y~~ to require rather less permanent apparatus. We need to have the inverse squares written out with part of the beginning of the square repeated <sup>and</sup> again at the beginning, and in rather small letters. In order to work a particular crib we take some ~~part~~ <sup>place</sup> in ~~place~~ with the inverse oblong and write the diagonal down the side of it, and ~~XX~~ ~~XXXXXX~~ write the crib along the bottom. Then for each letter of the crib (either code or decode) we punch a hole ~~straight~~ in the column in which it occurs, and in the line ~~next~~ after it (Fig 47). We then move this mask over the inverse oblong. Each position of the mask corresponds to a different start on the rods. The pair of letters showing through the two holes in a column give the coupling which is a consequence of the constetation written in that column (Fig 48).

Another advantage of this method is that we can test all colours with one mask. This advantage can however also be got by making inverse rods with all the colours on one rod.

# Charts.

When we want to try the same decode for ~~many~~ a great many different messages, and perhaps for many different places in the same message it may be worth while, <sup>to</sup> make special statistics for that crib. We can make statistics of the positions in which there will be 'clicks'. There is quite a problem as to the form in which the statistics ought to be presented. I will describe two forms which have actually been used; named after the principal cribs for which they were made. First however I must explain

## PERCOMMANTANTE charts.

This is the perfect form of chart for use when the position of the crib in the message is known. The chart ~~has several major divisions according to the different~~ terminology I shall use. Let us take for example the crib ~~IBNLESSELY~~ fitted onto a part of the message ~~AEIRCHTWBJ~~. There is a click as shown below

19	20	21	22	23	24	25	26	1	2	3	rod positions
A	E	I	R	C	M	T	W	B	Z	J	message
X	B	R	U	E	S	S	E	L	X	X	crib
N	V	Y	L	C	O	T	W	B	P	U	rod
D	G	G	K	W	C	U	E	L	A	B	rod

$\frac{1}{2} \times 9 = 9$

As the constatacion of the click are consecutive I shall say that the 'click distance' is 1. W is called the 'first cipher letter' and B the second cipher letter, E the first and L the second 'crib letters'. As the first letter of the crib comes at rod position 19 we will say that the 'rod start' is 19. As the first crib letter E is the eighth letter of the crib we say that the crib position of the click is 8.

## PERCOMMANTANTE charts.

This is the perfect form of chart for use when the position of the crib in the message is known exactly. The chart has several major divisions according to the different possible first crib letters. Each of these major divisions is further divided into lines labelled with the second crib letters, and columns labelled with the first cipher letters. In the ~~resulting~~ resulting small

rectangles are written the second cipher letter and the red start. Thus the eighth no. j or division of a PERCOMMANDEMENT type chart made out for XBRUESSELXX would look like this

A	B	C . . .	W . . .
B <sup>2</sup> <sub>1</sub>			B 19
B <sup>2</sup> <sub>2</sub>			
B <sup>2</sup> <sub>3</sub>			

all entries apart from the one corresponding to the click shown in Fig 49 having been omitted. The <sup>numbers</sup> letters written above on the right of the letters in the names of the rows distinguish between different occurrences of the same letter in the crib. By writing the message downward in gauge with the lines of the chart it is very easy to see the possible clicks. We note down the red starts, and, if we find one of them repeated try it out by the method described at the beginning of the chapter.

#### BRUESSEL type charts.

These have the advantage over the PERCOMMANDEMENT type charts that one can investigate all possible positions of the crib in the message without doing them all independently, but it has some ~~unpleasant~~ counterbalancing disadvantages. In the form in which they were made for the Railway traffic all three colours were put ~~together~~ together and there were separate sheets for the different click distances. I now think that it might be better to separate the colours and to have three or four click distances on a sheet. In any case the sheets are further divided into lines according to the different first cipher letters and the entries in the lines consist of the second cipher letter, the red start and the crib position of the click. Thus the click shown in Fig 44 would be represented on sheet I in line W by the entry B 19<sup>8</sup> in green. The chart is usually used one sheet at a time;

the message is written out with plenty of room for entries below it. Whilst using sheet I ~~xxxxxx~~ for each letter of the message we take the corresponding line of the sheet and look in it for the letter which comes next in the message. For each such entry that we find we ~~xxxxxxxxxxxxxxxx~~ enter the red start on the message under the letter which corresponds to the first ~~xxxxxx~~ letter of the crib. We know where this is because the entry on the chart gives the crib position. When we get ~~xxx~~ the same number twice in a column we try <sup>out</sup> the corresponding red position and position in the message.

A possible improvement of the lay out which might combine the advantages of the PERCOMMANDEANTE and BRUESSEL type charts would be to take a fairly wide column for each click distance, all the columns being the same width, ~~xxxxxxxx~~ instead of having separate sheets, and to make the lines fairly deep. The message could then be written out in gaps with the chart. However I am afraid that this might <sup>make</sup> ~~xxx~~ both chart and message unwieldy. ~~xxxxxxxxxxxxxxxx~~ <sup>improvement</sup> ~~xxxxxxxx~~ An alternative possible improvement would be to have separate columns for the different <sup>cipher</sup> second letters. This would also mean having rather large charts, because of the great variation of the number of letters that would have to go into a rectangle.

Click position  
(index)


4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95	96	97	98	99	100	101	102	103	104	105	106	107	108	109	110	111	112	113	114	115	116	117	118	119	120	121	122	123	124	125	126	127	128	129	130	131	132	133	134	135	136	137	138	139	140	141	142	143	144	145	146	147	148	149	150	151	152	153	154	155	156	157	158	159	160	161	162	163	164	165	166	167	168	169	170	171	172	173	174	175	176	177	178	179	180	181	182	183	184	185	186	187	188	189	190	191	192	193	194	195	196	197	198	199	200	201	202	203	204	205	206	207	208	209	210	211	212	213	214	215	216	217	218	219	220	221	222	223	224	225	226	227	228	229	230	231	232	233	234	235	236	237	238	239	240	241	242	243	244	245	246	247	248	249	250	251	252	253	254	255	256	257	258	259	260	261	262	263	264	265	266	267	268	269	270	271	272	273	274	275	276	277	278	279	280	281	282	283	284	285	286	287	288	289	290	291	292	293	294	295	296	297	298	299	300	301	302	303	304	305	306	307	308	309	310	311	312	313	314	315	316	317	318	319	320	321	322	323	324	325	326	327	328	329	330	331	332	333	334	335	336	337	338	339	340	341	342	343	344	345	346	347	348	349	350	351	352	353	354	355	356	357	358	359	360	361	362	363	364	365	366	367	368	369	370	371	372	373	374	375	376	377	378	379	380	381	382	383	384	385	386	387	388	389	390	391	392	393	394	395	396	397	398	399	400	401	402	403	404	405	406	407	408	409	410	411	412	413	414	415	416	417	418	419	420	421	422	423	424	425	426	427	428	429	430	431	432	433	434	435	436	437	438	439	440	441	442	443	444	445	446	447	448	449	450	451	452	453	454	455	456	457	458	459	460	461	462	463	464	465	466	467	468	469	470	471	472	473	474	475	476	477	478	479	480	481	482	483	484	485	486	487	488	489	490	491	492	493	494	495	496	497	498	499	500	501	502	503	504	505	506	507	508	509	510	511	512	513	514	515	516	517	518	519	520	521	522	523	524	525	526	527	528	529	530	531	532	533	534	535	536	537	538	539	540	541	542	543	544	545	546	547	548	549	550	551	552	553	554	555	556	557	558	559	560	561	562	563	564	565	566	567	568	569	570	571	572	573	574	575	576	577	578	579	580	581	582	583	584	585	586	587	588	589	590	591	592	593	594	595	596	597	598	599	600	601	602	603	604	605	606	607	608	609	610	611	612	613	614	615	616	617	618	619	620	621	622	623	624	625	626	627	628	629	630	631	632	633	634	635	636	637	638	639	640	641	642	643	644	645	646	647	648	649	650	651	652	653	654	655	656	657	658	659	660	661	662	663	664	665	666	667	668	669	670	671	672	673	674	675	676	677	678	679	680	681	682	683	684	685	686	687	688	689	690	691	692	693	694	695	696	697	698	699	700	701	702	703	704	705	706	707	708	709	710	711	712	713	714	715	716	717	718	719	720	721	722	723	724	725	726	727	728	729	730	731	732	733	734	735	736	737	738	739	740	741	742	743	744	745	746	747	748	749	750	751	752	753	754	755	756	757	758	759	760	761	762	763	764	765	766	767	768	769	770	771	772	773	774	775	776	777	778	779	780	781	782	783	784	785	786	787	788	789	790	791	792	793	794	795	796	797	798	799	800	801	802	803	804	805	806	807	808	809	810	811	812	813	814	815	816	817	818	819	820	821	822	823	824	825	826	827	828	829	830	831	832	833	834	835	836	837	838	839	840	841	842	843	844	845	846	847	848	849	850	851	852	853	854	855	856	857	858	859	860	861	862	863	864	865	866	867	868	869	870	871	872	873	874	875	876	877	878	879	880	881	882	883	884	885	886	887	888	889	890	891	892	893	894	895	896	897	898	899	900	901	902	903	904	905	906	907	908	909	910	911	912	913	914	915	916	917	918	919	920	921	922	923	924	925	926	927	928	929	930	931	932	933	934	935	936	937	938	939	940	941	942	943	944	945	946	947	948	949	950	951	952	953	954	955	956	957	958	959	960	961	962	963	964	965	966	967	968	969	970	971	972	973	974	975	976	977	978	979	980	981	982	983	984	985	986	987	988	989	990	991	992	993	994	995	996	997	998	999	1000	1001	1002	1003	1004	1005	1006	1007	1008	1009	1010	1011	1012	1013	1014	1015	1016	1017	1018	1019	1020	1021	1022	1023	1024	1025	1026	1027	1028	1029	1030	1031	1032	1033	1034	1035	1036	1037	1038	1039	1040	1041	1042	1043	1044	1045	1046	1047	1048	1049	1050	1051	1052	1053	1054	1055	1056	1057	1058	1059	1060	1061	1062	1063	1064	1065	1066	1067	1068	1069	1070	1071	1072	1073	1074	1075	1076	1077	1078	1079	1080	1081	1082	1083	1084	1085	1086	1087	1088	1089	1090	1091	1092	1093	1094	1095	1096	1097	1098	1099	1100	1101	1102	1103	1104	1105	1106	1107	1108	1109	1110	1111	1112	1113	1114	1115	1116	1117	1118	1119	1120	1121	1122	1123	1124	1125	1126	1127	1128	1129	1130	1131	1132	1133	1134	1135	1136	1137	1138	1139	1140	1141	1142	1143	1144	1145	1146	1147	1148	1149	1150	1151	1152	1153	1154	1155	1156	1157	1158	1159	1160	1161	1162	1163	1164	1165	1166	1167	1168	1169	1170	1171	1172	1173	1174	1175	1176	1177	1178	1179	1180	1181	1182	1183	1184	1185	1186	1187	1188	1189	1190	1191	1192	1193	1194	1195	1196	1197	1198	1199	1200	1201	1202	1203	1204	1205	1206	1207	1208	1209	1210	1211	1212	1213	1214	1215	1216	1217	1218	1219	1220	1221	1222	122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### Making of charts

Although there is so much room for variation in the form which a chart can take the manner in which they are made is fairly stereotyped. There are two kinds of click to be catalogued, called 'direct' and 'cross'. Direct clicks are those in which both letters of the crib occur on the same rod. Both clicks in Fig 44 are direct clicks. Cross clicks have one of the crib letters on one rod and the other on the other.

When cataloguing cross clicks we make 25 pictures like Fig 50, by writing the crib diagonally and filling up a square with rods, and finally copying the left lower half into the right upper half symmetrically across the diagonal. The different pictures

  
correspond to different rod starts. Each square above the diagonal gives us an entry for the chart. The lower letter is the first cipher letter, and the upper is the second cipher letter. The row gives the click position, i.e. with a BRUESSEL type chart the number in the 'index' position. The click distance (i.e. the sheet, with BRUESSEL type) is determined by how far the square is from the central diagonal; in the figure the squares corresponding to click distance III are ringed in pencil. With a PERCOMMANDANTE type chart we should not use the diagonals but the columns. Some of the

~~figures are crossed out in the original~~  
squares do not correspond to possible entries, as they could only arise from rods paired with themselves. These ~~xxx~~ squares have been crossed out in Fig 50.

For cataloguing direct clicks we have to find all cases in which a pair of letters on a rod can fit with a pair of letters of the crib, e.g.

I	B	R	U	E	S	S	E	L	I	X	I	X	crib
D	E	G	G	K	W	U	U	E	L	A	B		rod

Each such case will give us 25 different entries in the chart,

all with the same click distance, rod start and crib position. In cataloguing these either in a PERCOMMANDANT or a BRUSSEL chart it is sufficient if we put the second cipher letters all in similar positions and only once enter the remaining information, for each set of 25.

### X-charts

Sometimes one will find messages with about 30% of X's in the decode. These can be got out by a 'majority vote' method, looking for the R.H.W. starting position which gives the greatest number of clicks if we assume the message to say XXXXX all through. If there are actually 30% of X's there will be about 22 genuine clicks between X's per T.O. : there will also be an average of about 0.5 ~~click~~ apparent clicks arising from letters which are not X, giving altogether 2.7 clicks per T.O. with the correct start. With the wrong start we have one bogus click per T.O. If we do not know where the T.O. is these figures have to be modified. In the right place we have 3.7 clicks per length of 26, and in the wrong place 2.0.

~~There are two forms of X-chart, one favoured by Kendrick and one by Turing.~~ With X-charts there are less variables involved than with ordinary charts, as there is no question as to where the crib should be set against the message. The variables involved therefore are the first and second ~~click~~ cipher letters, the click distance, and the rod position of the first constatation of the click. There are two ways of setting the chart out, one favoured by Kendrick and one by Turing.

With Turing's form of chart ~~the first~~ there are 26 lines named after the first cipher letters and 26 columns corresponding to the possible click distances. The second cipher letter and the rod position are entered in the square. The chart can be used by writing the message out in gauge with the chart, and putting each letter in turn over the corresponding letter in the left-hand

column which names the lines, and looking for each letter ~~of~~  
among the next 26 of the message in the square of the chart  
directly below it. ~~With the~~ In noting the click down we ~~calculate~~  
the implied rod start of the message by subtracting the position  
in the message of the first cipher letter from the rod position  
of the first cipher letter, i.e. the number in the square. We  
enter against this rod start the position in the message of the  
first cipher letter. The rod start with the greatest number of  
entries against it is presumed to be the right one. To ~~read~~ read  
the message after we have found the R.E.W. rod start we can  
try setting up the rods giving the clicks and see if this results  
in any further identifications, but this hardly ever gives the  
solution. ~~The~~ The generally accepted method is to take ~~majority~~  
~~enter~~ the couplings giving the clicks and note down from a  
catalogue the places in which they could occur, and then take a  
'majority vote'.

In making an X-chart we can make a set-up like Fig 4<sup>b</sup>. This  
will measure 26 x 26 and ~~with~~ only one of them will be  
needed. It will simply consist of a rod-square rearranged with the  
X's down the diagonal. When making the entries for a particular  
value of rod position of the first constetation of the click (i.e. the  
entries where a particular number is written in the square) we copy  
in pencil  
down a line from the rearranged rod-square, starting immediately  
after the X, across the top of the rod square, and also the column  
starting at the same X. ~~The~~ The entry to be made in any column  
then  
can be seen by looking at the top. Having made these entries we  
rub out the lines at the top and replace them with others.

In Mendrick's type of XChart the ~~main~~ names of the lines  
 the first  
 give ~~the~~ ~~xxx~~ of the cipher letters/ The columns give the position  
 second cipher letter and  
 of the other cipher letter, and the entry in the square is the  
 position of the first cipher letter~~xxxxxxx~~. This form of chart is  
 particularly useful when we have a hunch about the red start.

15  
79

Consecutive tables.

In the second part of the process, where we are finding the position of the middle wheel we can speed up the work by the use of consecutive tables. These are of two kinds, forward and backward, and look very like rod squares. The letter in column 18, say and row R of the forward consecutive square is the letter which occurs in column 19 of the rod with R in column 18. The letter and row R in column 18 of the backward consecutive square is that which occurs in column 17 on the same rod. Like rod squares and inverse squares these consecutive squares 'have a diagonal' i.e. can be filled in from a single upright by writing 'the diagonal' diagonally downwards to the left. In our DANZIGVON example we could have used the backward consecutive as soon as we had found the couplings ku, ep, fr, qn, sy, td, vh, lw before the T.O. and sw, ce, le after it. We should have laid rulers against the lines o, s of the backward consecutive square, and read off the consequences before the T.O. of having ce after it, in the various possible positions of the middle wheel, and would have looked to see whether these consequences were consistent with our data. We could then have repeated with ~~ws xxxxxxxxxxxttttt~~ looking only at the positions consistent with ce. The forward consecutive can be used when the place has been found for reading off the couplings after the T.O. (although this is only a small advantage), or in a case where we have started from the end of the message and worked backwards.

## Chapter V. Coupling catalogues

When we have found the rod position of the R.H.W. and a few couplings for a message it is possible to find the positions of the other wheels from a suitable catalogue.

### Short catalogue

One method is to try independently all the possible positions for the middle wheel. We shall want to know the middle wheel couplings which are consequences of these various assumptions. This can be done by setting up inverse rods for the middle wheel. The rods are paired off according to the R.H.W. couplings, i.e. M.W. output, ~~xxxxxxxx~~. This has been done for the the couplings ku,fx,op which arose in the DANZIGVON crib in Fig 56, assuming the red wheel in the middle. The pairs in each column of this set up give possible M.W. couplings. We have ~~xxxx~~ now to find out whether these couplings are possible. Our procedure is rather different according as the U.K.W. does or does not rotate. In the case that the U.K.W. does not rotate it will be sufficient to have a (the rows and columns lettered preferably with the diagonal alphabet) Pose sheet ~~xxxx~~ in which, in the RW square ~~xxxx~~ are entered the positions of the left hand wheel at which the ~~xxxxxx~~ RW is one of the pairs in the L.H.W. output alphabet. This is known as the 'short catalogue' for this wheel. To use it in connection with the DANZIGVON crib we should take each column of Fig 56 in turn and look up the pairs in it on the short catalogue and see if all the squares had a number in common. If we found such a case the number in the square would give the L.H.W. rod position, and the column of ~~xxx~~ Fig 56 would give the M.W. position. Actually the U.K.W. rotates for our example so that we should have no success.

In the case that the U.K.W. rotates we need essentially the same short catalogue, but we arrange it slightly differently. Instead of the lines of the catalogue corresponding to fixed output letters they correspond to fixed distances on the diagonal, <sup>(say 4 letters)</sup> between the output letters. This may be seen from Figs 52, 53 which illustrate such a catalogue. The pairings are written above the ~~xxxxxxxxxxxxxxxxxxxx~~ figures giving the positions ~~of~~

Green

Q 31121-0496754-0496754 Q

X	W	E	R	T	Z	U	I	O	A	S	D	F	M	N	O	P	Q	R	Y	X	C	V	B	N	M	Y	Z
10. 32						2. 16	4. 3. 7. 9.	15. 18				5. 4				26. 212. 513. 35	6.									1. 7. 21	
10. 32	26	16				10. 14. 20.	18. 1. 2. 14.	2. 9. 24.				4.				4. 8.		7. 8. 23	15. 21.	5.					4.	1. 7. 21	
36			2. 9.	7.												4. 13.			2. 5. 10. 15. 20.	24. 9. 25	3. 16.	8.					
6. 20.			3. 9. 18.			10. 9. 18.	7. 20. 26	9. 12. 25.								14. 19. 20. 24.										14. 7.	
8			21. 6. 20.	26. 8.	4.	1. 2. 3. 7.	16. 18. 24.	2. 6. 9. 15.								24. 7. 12.										10. 18.	
3. 20. 20	12. 13. 7.	3. 9. 11.				26.	10. 20									2. 6. 9. 15.										1. 14.	15.
4. 8.	10. 20.					2. 15.	8.	9. 12.	5. 6. 1. 20.						7. 9.	13.	4. 23.		7. 3. 9.							24. 6.	
3. 1. 2. 25. 8.	5. 18. 6. 26.					5.		20.	7. 12. 16. 9.	21.						12. 18.			1. 13.	24. 25.	4.						
7. 8.			22.			2. 15. 5.		6.	7. 9.							10. 9.			13.	24.						1. 18.	15.
15. 18. 3. 4. 12. 7.	20. 10. 22.							5.	14.	14. 24.						20. 20. 16.			16. 13.							1. 18.	15.
24.	9. 20. 8.							5.	7. 14.	12. 9. 26. 23.						15. 16. 4. 9.			13. 20.							2. 14. 19.	
21.	26.	4.	9. 12. 20.	6.		7. 14.		10. 13. 12. 8.								16. 24.			2. 14. 19.							3. 5. 16. 17.	
1. 2. 4.	9. 13. 2. 9. 4.	5. 16. 7.				1. 10. 18. 7.		4. 9. 13. 19. 22.								10. 8.			20.							10. 14.	
4. 6. 20. 2. 5. 3.			7. 21.	20. 9. 10. 4.	19. 24. 23. 1. 22.											13.			15. 21. 2.							20. 20. 16.	
26. 18.	14. 10. 8. 2. 5.	21.				4. 2.		15. 13.								20. 2. 3. 10. 14.			2. 7.							1. 2. 3.	
2. 10. 14. 13.	9. 13. 5. 4. 6.	7. 8.				1. 26.	18.	11.								2. 5. 10. 14.			16. 17.							1. 2. 3.	
5. 18.	2. 20.	9. 18.				19.	2. 20. 16.	6. 18.								1. 2. 3. 10.			12. 7.							4. 11. 26.	
25.	3. 8.	1.	13.			22.	4. 9. 16. 24.	26. 2. 15.								16. 14. 1. 17.			19.							4. 11.	
6. 7. 1. 10. 20.	24. 26.	1. 2. 3. 4. 12.	7. 16.			8.	2. 25. 3. 11.									12.			9.	26. 5. 2.	19. 9.						
8. 23.	14. 2. 12. 14.	18. 24.				13. 20. 3.										7. 14. 7.			1. 5. 20. 24.							6. 11. 26.	
19. 15. 24.			7. 11. 12. 3. 25.	1. 6. 12.		20. 9. 14. 9.	7.	20. 2. 5. 10. 14.								2. 5. 10. 14.			8.		</						

Fig 51. Short catenulae.  
Refring 14.

III Barney Barney

49 49

Short catechism for use with children

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95	96	97	98	99	100
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63	64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95	96	97	98	99	100



relating to technology

Small sheet of short columns for

1	A	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
2	C	U	U	A	R	T	A	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
3	E	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
4	G	P	T	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L
5	K	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
6	M	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
7	O	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
8	Q	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
9	S	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
10	U	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
11	Y	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
12	X	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	
13	Z	Q	W	E	R	A	T	Z	Z	U	I	0	0	A	S	A	S	D	Z	T	A	A	A	H	T	A	K	T	P	Y	X	A	C	V	B	B	N	N	N	L	L	L	L	

R Z Z O X F M C P W H S I E A N T L G A U W B D J S w  
 C J R L B P F W Z Q X V V M A K H S I P T A N Y K G k  
 U A C G L V Y E J D O R W N X Q H S I E N F K D T w  
 G O W R R U P S B Y J L F O Z Q M V E N D I T L C A w  
 I S V H A B X R K F A T E L U W J D O R M G P F Z I y e  
 U M O V P H A I T D U C K G O R Y B P S F F J A X X i  
 D B J W N C T P G S Z R A I E K F A T L H Y G U O O a f  
 H D Y G W N C J S L R M P X X B I A F O X J E M V K b  
 N K E M V Z Y H D U T W O R P G S Z Q J X H I A A F t  
 T A N I C K G O U R S Z X J F I E P V K A D D H L Y a u  
 P R L B O X J F I Z K A T Y H D U W K C J O S S G M y f  
 J G S P F Q B X H A M E N K Y Y V U L D I T H W N C g  
 L A B Y J S O Z F I P P H A T X N Y D G G K W C U E t  
 B P U S Z M Q A R H I X J F B A Z E N Y Y L C O T W v  
 S N X K D A U G O B Y J B Z C M X F H H P E V I R a u  
 F X H E M V K D A T L Y C C N O W G A U K R L B P J y  
 M C P F S I H A N X K D U V L C G J J Y R B O T W D a  
 A Y K Z A T N L L E G U Y H D V O T W B C P M X I R c  
 V Y G D O J S M C P F I B A V H K K X T N A X E F L d  
 A E T T I Y D N X K A G A U W L B R M T O Z A Y S H l  
 H F A K D L V Y G O N W B J P P C I M S U R G A B f  
 E L C J T O R B M M W F Z P G S C I R A V X K N Y U x  
 K T A N Y G E U W C B L S R J D O Y Z W M X P H Y f  
 O F I A E E Z K A V P H M D I T L N C W B S U R M X n  
 Z W M X H A I E V N N A D T K F A X U E L C Y J B P y  
 Y I D U L W O T J O C K G N S U R M B X A V A Z E N b

Railway  
Couplings Wheel

Fig 54. Coupling of U.K.W. & L.W. to

anti of that catalogue.

of the L.H.W. in which these pairings occur, the U.K.W. <sup>being</sup> understood to be in the zero position. Either form of short catalogue may be made by setting up the L.H.W. rods paired according to the U.K.W. as in Fig 54, and analysing the resulting pairs.

To understand the use of the short catalogue when the U.K.W. rotates we must ~~remember~~ remember that if the U.K.W. and L.H.W. are rotated in step the effect is a ~~single~~ slide along the diagonal of the resulting pairs. If we are given actual pairs for which the U.K.W. was not in the zero position we can slide the pairs along the diagonal until we have pairs which would have occurred with the U.K.W. in the zero position. This will show up on the catalogue because there will be a ~~link~~ <sup>number</sup> in common in the squares under these pairs. For instance in the case of the DANZIGWON crib we found the middle wheel to be ~~xxxxx~~ red in position 14. This gives the middle wheel couplings <sup>pn,ve,hn,uy</sup> ~~xxxxx~~ as consequences of the R.H.W. couplings qn,uk,fx,ep. These can be read off from Fig 56, although of course we should only set up the M.W. inverse rods in a case where we did not know the M.W. position. If we slide ~~xxxxx~~ ~~xxxxx~~ pr,hn,ev,uy ten places forward along the diagonal we get wg,mi,zf,ke, and in each of the squares wg, mi,zf, ke on the green (L.H.W.) short catalogue we find the number 4, i.e. these pairs occur at U.K.W. 0 L.H.W. 4; consequently qn,... occur at U.K.W. 10, L.H.W. 14. The mechanical process would actually be to take pr on the small sheet of the catalogue and lay it against ve on the large sheet. This automatically results in wg and <sup>14</sup>mi being together and all other pairs of pairs resulting from sliding pr,ev along the diagonal. We look in the pairs of squares to see if there are numbers in common. When we find such a case we have to look in a third square resulting from sliding <sup>n</sup>hm. It is as well therefore to have rulers in gauge with the catalogue to measure off the distances. Having found the right amount of slide forward on the diagonal, i.e. to the right in the catalogue we calculate the positions of the wheels from the formulae



U.K.W. position = slide forward on diagonal  
L.H.W. position = number in square + slide

The Turing sheets

The short catalogue should work very well when the Umkehrwalz rotates, and there is no information ~~xxxx~~ connecting the position of the U.K.W. with the positions of the other wheels~~xxxxxxx~~. In the case of a fixed U.K.W. we can often make use of an analysis of R.H.W. couplings.

The lay out of the catalogue is largely determined by the special method ~~xxxxxxxx~~ by which they are made, but it seems to be reasonably convenient in use. The catalogue is divided into sheets numbered 1 to 13. Each of these sheets consists of a 26x26 square with margin at top and left hand side, preferably on 1/3" gauge. ~~The number of the sheet is given at the top of the sheet. The letters and numbers in the margin are the only ones concerned when the sheets are being used, the others being part of the construction, and left on to help in tracing errors. The entries 10,18,21 in the square in column 15 and the row with KV in the margin mean that the pair KV occurs when the M.W. is in position 15 and L.H.W. in any of the positions 10,18,21. In order to find the positions at which two couplings can occur we have only to find the corresponding lines of the catalogue against one another and compare the numbers in the adjacent squares. It is fairly easy to find the right sheet because the number of the sheet gives the distance along the diagonal of the two letters of the pair, e.g. K and V are at distance 5 along the diagonal (KPYXCV) and KV occurs on sheet 5.~~

One such sheet is shown in Fig 3/4 partly constructed. The letters and numbers in ink are the only ones concerned when the sheets are being used, the others being part of the construction, and left on to help in tracing errors. The entries 10,18,21 in the square in column 15 and the row with KV in the margin mean that the <sup>coupling</sup> pair KV occurs when the M.W. is in position 15 and L.H.W. in any of the positions 10,18,21. In order to find the positions at which two couplings can occur we have only to find the corresponding lines of the catalogue against one another and compare the numbers in the adjacent squares. It is fairly easy to find the right sheet because the number of the sheet gives the distance along the diagonal of the two letters of the pair, e.g. K and V are at distance 5 along the diagonal (KPYXCV) and KV occurs on sheet 5.

### Construction of the Turing sheets

The construction of the catalogue depends on making almost simultaneously all the entries corresponding to ~~pairwise~~ cases in which the current flows through the same two wires of the M.W. In the partially constructed sheet S in Fig 5b some of the diagonals have been filled in fully, and each of these corresponds to a pair of wires of the M.W. As the M.W. rotates the red points at the right hand ends of the wires move steadily backwards along 'the diagonal'. We see ~~that~~ also that as ~~we~~ move along the filled in diagonal the red position steadily increases, and the letters in the pairings move slide backwards along 'the diagonal'. Meanwhile the left hand ends of the wires are steadily rotating, so that the middle wheel couplings are sliding along 'the diagonal'. The entries in the squares are the positions of the L.H.W. where these M.W. couplings can occur, and the slide along the diagonal amounts to a diagonal movement along the short catalogue. Take for instance the ~~the~~ filled in diagonal on Fig 5b nearest to the central diagonal. The second entry on this diagonal is 2,5,16,26 which is the entry at HL in Fig 5i: next along the diagonal in Fig 5b is the entry 10 which occurs at GM in Fig 5i, and so on, the diagonal in Fig 5i being repeated backwards in Fig 5b.

This phenomenon may ~~also~~ be explained with reference to the red square, instead of the wheels: this is really more practical, as we have to make the catalogue up from the red square. A possible method for making up the catalogue would have been this. In each square on the sheets we write in, in pencil, the M.W. couplings which would be needed to produce the <sup>part of the beginning of the line as</sup> M.W. output ~~required~~ at the <sup>beginning of the</sup> M.W. position given by the <sup>the</sup> column in which the square occurs. To do this we should have to write down in each line the inverse rods named after the letters at the beginning of the line. This has been done in a part of Fig 5b (top R.H. corner). We should then have the square filled with one inverse (M.W.) square, with top and bottom reversed, and another such reversed square somewhat displaced upwards. The entries in green ink could be obtained by

replacing each pair of pencil letters by the corresponding entry on Fig 5<sup>1</sup>, i.e. by the position of the L.H.W. at which that pair of letters occurs as L.H.W. output. Now the whole of the pencil square can be obtained from its top line simply by filling in along diagonals. Translated into terms of the green ink entries this means to say that we only need to be given the positions at which <sup>the</sup> ~~the~~ start copying from the short catalogue.

Actually we copy out the diagonals of the short catalogue onto staircase shaped strips (known as 'Christmas decorations' or 'hand frills') in reversed order, with the position in the short catalogue written above each square. These hand frills are numbered by the (constant) distance apart on 'the diagonal' of the pair of letters on them; e.g. in the hand frill <sup>no 5</sup> shown in position for copying in Fig 5<sup>2</sup> I and F are at distance 5 on quertzu and so are D and K. Instead of actually filling in the whole square with pairs of pencil letters we take the entries which might have been made in the top line, and write them in the top margin, and also ~~write~~ put the entries which might have gone in the left hand column into the left hand margin. In order to find what hand frill to use for a particular diagonal the distances apart along quertzu of the letters along the top are calculated. This should be done quite independently, to give a check on incorrectly copied letters (see 'Mystic numbers').

The reason for having the imaginary grid squares implied in the construction inverted is in order that the writing of diagonals may be from left to right and downwards, which is considered easier than from right to left and downwards.





Solving a short crib

The ~~first~~ chief application of the Turing sheets is to the solution of cribs from a length of 2 to 6 letters. We set up the inverse rods as usual, but find that ~~xxxxxxxxxxxx~~ by no means <sup>incorrect</sup> all the ~~xxxxxxxx~~ positions are eliminated by coupling contradictions. We therefore look to see whether there is any position in which the couplings can occur. Take for example the crib ANK, with end wheel order I III II (red, green, purple), U.K.W. pos. 0 cipher ~~ANK~~. We set up the inverse rods as in Fig 57, and for each column of the resulting set up compare the lines of the catalogue named after the pairs in the column. For each pair we shall want to find quickly the right sheet on which to look, and this means subtracting the pairs on the diagonal (i.e. finding their distance apart on quortzu). To do this we can either have a table of differences or else use 'mystic number rods'

'Mystic numbers'

Fig 58 shows a table of 'mystic numbers' for the red wheel. The meaning of the table is this. Take the 8th line for example. It could be made by taking ~~rod~~ inverse rod Q and inverse rod Q, <sup>quortzu</sup> Q being eight places on along ~~the xxxxxxxx~~ from Q. We lay the two rods together and find the differences of the resulting pairs; e.g. the <sup>fifth</sup> ~~third~~ entry in line 8 is 6, and the fifth letter of the red inverse ~~xxxxxxxx~~ rod Q is Y, the fifth letter of inverse rod <sup>6</sup> Q is F, and Y and F are ~~xxxxx~~apart on quortzu(~~FGHJKPY~~). If then we had a set up of inverse rods including the pair QQ we could use the series of numbers of line 8 of the mystic numbers to <sup>tell</sup> ~~give~~ us on which sheets the various pairs should be looked up. However we can also use line 8 of this table on many other occasions. Suppose for example that the pair ES of inverse rods is up. The series of sheets on which we have to look is again given by line 8, but we have to start in the third column under E instead of at the beginning under Q. Quite a convenient arrangement is to have the lines of the table written out on rods in gauge with the inverse rods and of double length. (This was once done for the service machine wheel III. Three lines of the table were put onto ~~xxxx~~ three sides of Mr Knox's blank wooden inverse rods, and the fourth side occupied with the letters of the diagonal, in that case AEE A BOD... It was

[illegible][illegible][illegible]

Fig. 57. Set of waste rods for  
BRC  
ANX, and mychi no. rods for feeding water sheets.

	Q	W	E	R	T	Z	U	I	O	A	S	D	F	G	H	J	K	P	Y	X	C	V	B	N	M	L
1	3	13	2	9	3	2	10	2	12	4	3	12	4	2	4	2	9	8	12	9	10	12	1	9	2	3
10	11	7	6	5	8	8	10	8	7	9	10	2	6	6	11	9	6	5	1	2	13	8	11	1	6	
3	8	6	4	4	5	6	6	6	5	5	13	12	2	8	11	7	3	3	11	11	3	4	10	8	4	7
4	9	9	2	12	3	8	10	3	9	9	11	10	4	9	3	5	12	7	3	12	6	2	7	5	9	5
5	12	11	12	12	11	12	13	11	5	7	7	8	13	1	9	12	2	5	2	3	8	5	4	8	7	12
6	12	1	10	2	11	11	1	7	7	3	5	1	5	11	8	4	12	6	11	1	5	8	9	10	10	11
7	4	3	4	6	8	3	3	9	11	1	4	9	7	6	8	10	11	3	13	4	2	5	11	1	13	9
8	6	9	8	9	6	1	1	13	13	8	12	5	10	10	6	9	6	5	10	7	11	3	2	4	11	7
9	6	5	11	3	2	1	3	11	4	10	2	4	6	4	5	8	4	2	7	6	13	12	5	6	5	9
10	2	2	1	7	4	5	5	2	4	2	7	6	8	3	12	6	7	1	6	4	4	9	7	4	7	3
11	1	12	5	5	8	7	12	6	10	11	3	6	7	12	10	9	10	12	8	13	7	7	3	2	5	1
12	11	8	3	1	10	10	4	8	1	1	9	7	10	12	13	12	3	10	1	10	9	9	1	12	1	3
13	11	10	1	1	7	2	8	1	11	13	10	2	8	11	10	1	1	7	2	8	1	11	13	10	2	8

I  
Red

Fig. 5-8. Magic squares for Rectangular sheet I

not a success as the rods were incorrectly copied). For the crib  
 BRC  
 ANX these mystic number rods are shown in position over the  
 inverse rods in Fig 67. Every fifth letter from the top XXX of  
 the mystic number table is also shown.

Another use for the mystic number table is in the making of the  
 Turing sheets. The line of pencil numbers along the top of any  
 sheet is the line of mystic numbers with the sheet number as its  
 line number, and starting at column L. *Given Page 68, 66*

The mystic numbers can of course be used by actual subtraction  
 from the inverse rods. However it is actually easier to ~~subtract~~  
~~subtract~~ do the calculation in terms of the letters of an  
 upright. It turns out that one can manage with one upright, which  
 one subtracts from itself, staggered various amounts. One can  
~~therefore~~ transform these letters into numbers XXX to simplify the  
 subtraction. I shall not give the details of this.

#### EINS catalogue

In this chapter and the last we have not exhausted all the  
 possible methods of dealing with the Unsteckered enigma, and enigma  
 with known Stecker. When the Umkehrwalz does not rotate we can  
 catalogue the result of encoding a short word such as EINS at  
 every possible position. The details of this are explained in Chapter

### Jeffreys sheets

In cases where the wheel order is unknown it is useful to have ~~sixty~~ the positions ~~where~~ and wheel orders where a coupling occurs all catalogued together. In order to make comparison of couplings feasible one puts the catalogue into the form of punched sheets, which can be laid one on top of another. These are known as Jeffreys sheets.

The actual form of the Jeffreys sheets catalogue is this. There are 325 sheets labelled AB, AC, ... AZ, BC, ..., BZ, ..., ..., YZ. Each sheet measures  $26 \times 20\frac{4}{5}$  plus margins of about three inches. They are divided into ~~sixty~~ columns an inch wide, and lines  $\frac{4}{5}$ " deep. The whole is further subdivided into squares  $\frac{1}{5} \times \frac{1}{5}$ . The  $\frac{4}{5} \times 1$  rectangles correspond to the different possible rod positions of the L.H. and M.W. The subdivisions of the rectangles correspond to the twenty possible wheel orders for L.H.W. and M.W. with the five first wheels of the service machine.

### Jeffreys-Turing sheets

There is a possibility of speeding up this work with short cribs where the U.K.W. rotates by making the Turing sheets in punched form. Suppose we expand every square of the Turing sheets into a rectangle  $7\frac{5}{8} \times 4\frac{5}{8}$  divided into 26 small squares, numbered 1 to 26 with two unused, and for each entry on the Turing sheet punch a hole in the corresponding small square. Then this is a matter of laying two of these sheets on top of one another, in such a way say that the lines VM and CR coincided would be to give us the positions in which the two couplings VM and CR occur when the U.K.W. is in the zero position: we also get the positions in which the couplings ~~may~~ slid along quertzu occur: but these after making a correction for the amount of slides are just the positions at which VM and CR occur including all possible rotations of the U.K.W. One would presumably normally place three sheets on top of one another, and there would have to be four different layings (because one could not have the sheets in cylindrical form). For this reason it would be better to have the sheets in double depth, but this would probably be out of the question.

*Mr. Kewer has not found a useful compromise by having two copies of Turing sheets: one of them punched & obtained by photography. This makes P-5-10*

a scheme on foot for putting the catalogue into cards.

Chapter VI. The steckered enigma. Bombe and Spider.

When one has a steckered enigma to deal with one's problems naturally divide themselves into what is to be done to find the Stecker, and what is to be done afterwards. Unless the indicating system is very well designed there will be no problem at all when the Stecker have been found, and even with a good indicating system we shall be able to apply ~~xxxxxxxx~~ the methods of the last two chapters to the individual messages. The obvious example of a good indicating system is the German Naval enigma cipher, which is dealt with in Chapter VII. This chapter is devoted to methods of finding the Stecker. Naturally enough we never find the Stecker without at the same time finding much other information.

Cribs.

The most obvious kind of data for finding the keys is a 'crib', i.e. a message of which a part of the plaintext is known. We shall mostly assume that our data is a crib, although actually it may be a number of constations arising from another source, e.g. a number of CHILLs or a Naval Banterismus.

FOURKEYPTWEEPT methods.

It is sometimes possible to find the keys by pencil and paper methods when the number of Stecker is not very great, e.g. 3 to 5. One would have to hope that several of the constations of the crib were 'unsteckered'. The best chance would be if the same pair of letters occurred twice in the crib (a 'half-bombe'). In this case, assuming 6 or 7 Stecker there would be a 25% chance of both constations being unsteckered. The positions at which these constations occurred could be found by means of the Turing sheets (if there were three wheels) or the Jeffreys sheets. The positions at which this occurred could be separately tested. Another possibility is to set up the inverse rods for the crib and to look for clicks. There is quite a good chance of any apparent click being a real click arising from because all four letters involved are unsteckered. The position on the right hand

wheel is given by the column of the inverse red set-up, and we can find all possible positions where the click coupling occurs from the Turing sheets or the Jeffreys sheet. In some cases there will be other constataions which are made up from letters supposed to be unsteckered because they occur in the click, and these will further reduce the number of places to be tested.

These methods have both of them given successful results, but they are not practicable for cases where there are many Stecker, or even where there are few Stecker and many wheelorders.

#### A mechanical method. The Bombe.

Now let us turn to the case where there is a large number of Stecker so many that any attempt to make use of the ~~known~~ unsteckered letters is not likely to succeed. To fix our ideas let us take a particular crib.

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23
D	A	E	D	A	Q	O	Z	S	I	Q	M	M	K	B	I	L	G	M	P	W	H	A
K	E	I	N	E	Z	U	S	A	E	T	Z	E	Z	U	M	V	O	R	B	E	R	I
24	25																					
I	V																					
Q	T																					

Presumably the method of solution will depend on taking hypotheses about parts of the keys and drawing what conclusions one can, hoping to get either a confirmation or a contradiction. The parts of the keys involved are the ~~xxxxxxxxxxxx~~ wheel order, the red start of the crib, whether there are any turnovers in the crib and if so where, and the Stecker. As regards the wheel order one is almost bound to consider all of these separately. If the crib were of very great length one might make no assumption about what wheels were in the L.H.W. position and M.W. position, and apply the method we have called a 'Stecker knockout' (an attempt of this kind was made with the 'Feindseligkeiten' crib in Nov. '39), ~~xxxxxx~~ or one might sometimes make assumptions about the L.H.W. and M.W. but none, until a late stage about the R.H.W. In this case we have to work entirely with constataions where the R.H.W. has the same position. This method was used for the crib from the ~~indef~~ Schluesesetzettel of the Verpoetenboot, with success; however I shall assume that all



wheel orders are being treated separately. As regards the turnover one will normally take several different hypotheses, e.g.

- 1) turnover between positions 1 and 5
- 2) " " " 5 and 10
- 3) " " " 10 and 15
- 4) " " " 15 and 20
- 5) " " " 20 and 25

*2nd 10*

With the first of these hypotheses one would have to leave<sup>out</sup> the constatactions in positions 2 to 4 ~~out~~, and similarly in all the other hypotheses four constatactions would have to be omitted. One could of course manage without leaving out any constatactions at all if one took 25 different hypotheses, and there will always be a problem as to what constatactions can best be dispensed with. In what follows I shall assume we are working the T.O. hypothesis numbered 5) above. We have not yet made sufficiently many hypotheses to be able to draw any immediate conclusions, and must therefore either assume something about the Stecker or about the rod start. If we were to assume something about the Stecker our best chance would be to assume the Stecker values of A and E, or of E and I, as we should then have ~~two~~ two constatactions corrected for Stecker, with only two Stecker assumptions. With Turing sheets one could find all possible places where these constatactions occurred, of which we should, on the average, find about 28.1. As there would be ~~28.1~~ <sup>28.1</sup> hypotheses of this kind to be worked we should gain very little in comparison with separate examination of all rod starts. If there had not been any half-bombes in the crib we should have fared even worse. We therefore work all possible hypotheses as to the rod start, and to simplify this we try to find characteristics of the crib which are independent of the Stecker. Such characteristics can be seen most easily if the crib is put in to the form of a picture

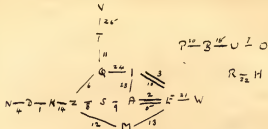


Fig 59. Picture from KOTKE ZUSATZU

Ant. Antennas 18 m. high mounted to allow for clearance.

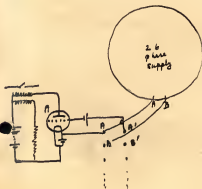
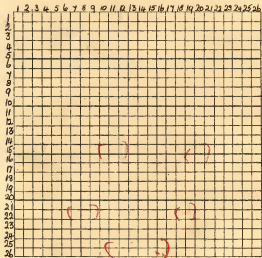


Fig 60. Circuit for 26 phase simultaneous scanning.



<sup>2 5</sup> <u>EAE</u>	<sup>3 10</sup> <u>EIE</u>	<sup>2 13 3</sup> <u>EAE</u>	<sup>2 9 8 6 14 3</sup> <u>EASZQIE</u>	<sup>13 12 9 3 5</sup> <u>EMZSAE</u>
	XHI			
	IQW			
	WAZ			
	ZRU			
	UBT			
	TAM			
	NDF			
	FCV			
	VON			
	<u>NEX</u>			

A /  
 B /  
 C /  
 D /  
 E /  
 F /  
 G /  
 H /  
 I /  
 J /  
 K /  
 L /  
 M /  
 N /  
 O /  
 P /  
 Q /  
 R /  
 S /  
 T /  
 U /  
 V /  
 W /  
 X /  
 Y /  
 Z /

XNW  
 WIO  
 OUF  
 FTK  
 KUT  
 PAS  
 SMD  
 DBN

OVC  
 CID  
 DNG  
 GTB  
 BUR  
 AZA  
 AWQ  
 QIH  
 HXE  
 ENO  
 KPK  
 PKP  
 SLS

APJY  
 YJY  
 PAYJ  
 JYJ

Fig 61. Stecker deductions with crib as p, with correct crib that is correct alphabet, but starting from an incorrect Stecker hypothesis 8/X. All other incorrect Stecker values of 12 are deduced.

<sup>2 5</sup> <u>EAE</u>	<sup>3 10</sup> <u>EIE</u>	<sup>2 13 3</sup> <u>EAE</u>	<sup>2 9 8 6 14 3</sup> <u>EASZQIE</u>	<sup>13 12 9 3 5</sup> <u>EMZSAE</u>
LHL	LSL	LHSL	LHIRVSL	LMRINL

Fig 62. Stecker deductions with same alphabets as Fig 61, but from correct Stecker hypothesis 47/L.

1 E A

2 25 3  
E A I E

W A P

Q K

T F

U O

I W

N X

B D

M S

V Q

L H

X E Y

N O U

F T

K U

P A

S M

J B

X H

W I

A Z

R U

V Q

M S

B D

U R

E V C

2 25 3  
S M O V

B A Y J

A T J Y

S F L

L H S L

A /  
B /  
C /  
D /  
E /  
F /  
G /  
H /  
I /  
J /  
K /  
L /  
M /  
N /  
O /  
P /  
Q /  
R /  
S /  
T /  
U /  
V /  
W /  
X /  
Y /  
Z /

like Fig 1<sup>st</sup>. From this picture we see that one characteristic which is independent of the Stecker is that there must be a letter which enciphered at either position 2 or position 5 of the crib gives the same result. This may also be expressed by saying that there must be a letter ~~which~~ such that, if it is enciphered at position 2, and the result reenciphered at position 5 the final result will be the original letter. Another such condition is that ~~the same~~ letter ~~which~~ enciphered successively at the positions 3, 10 must lead back to the original letter. ~~Three~~ other conditions of this kind are that the successive encipherments at positions 2, 23, 3 or at 2, 9, 8, 6, 24, 3 or at 13, 12, 8, 9, 5 starting from the same letter as before must lead back to it. There are other such series, e.g. 13, 12, 8, 24, 3 but these do not give conditions independent of the others. ~~There are other conditions independent of these~~ The letter to which all these multiple encipherments are applied is, of course, the Stecker value of E. We shall call E the 'central letter'. Any letter can of course be chosen as 'central letter', but the choice affects the series of positions or 'chaine' for the multiple encipherments. There are other conditions, as well as these that involve the multiple encipherments. For instance the Stecker values of the letters in Fig must all be different. ~~There are other conditions independent of these~~ The Stecker values for E, I, M, Z, Q, S, A are the letters that arise at the various stages in the multiple encipherments and the values for W, T, V, N, D, K can be found similarly. There is also the condition that the Stecker must be self-reciprocal, and the other parts of Fig 1<sup>st</sup>, P-B-U-O and R-H will also restrict the possibilities somewhat. Of these conditions the multiple encipherment one is obviously the easiest to apply, and with the crib as long as the one above it will be quite sufficient



If one has two of these 'Letchworth enigmas' one can connect the output points of the one to the input points of the other and then the connections through the two enigmas between the two sets of contacts left over will give the effect of successive encipherments at the positions occupied by the two enigmas. Naturally this can be extended to the case of longer series of enigmas, ~~xxxxxxxx~~ the output of each being connected to the input of the next.

Now let us return to our crib and see how we could use these Letchworth enigmas. For each of our 'chains' we could set up a series of enigmas. We should in fact use 18 enigmas which we will name as follows

A1,A2	with the respective positions	2,5
B1,B2		3,10
C1,C2,C3		2,23,3
D1,D2,D3,D4,D5,D6		2,9,8,6,24,3
E1,E2,E3,E4,E5,		13,12,8,9,5

By 'position 8' we here mean 'the position at which the combination numbered 8 in the crib, is, under the hypothesis we are testing, supposed to be enciphered'. The enigmas are connected up in this way: output of A1 to input of A2: output of B1 to input of B2: output of C1 to input of C2, output of C2 to input of C3: etc. This gives us five 'chains of enigmas' which we may call A,B,C,D,E, and there must be some letter, which enciphered with each chain gives itself. We could easily arrange to have all five chains controlled by one keyboard, and to have five lampboards shewing the results of the five multiple encipherments of the letter on the depressed key. ~~xxxxxxxxxxxxxxxx~~ After one hypothesis as to the rod start had been tested one would go on to the next, and this would usually involve simply moving the ~~xxxxxxxx~~ R.H.W. of each enigma forward one place. When 26 positions of the R.H.W. have been tested the M.W. must be made to move forward too. This movement of the wheels in step can be very easily done mechanically, the right hand wheels all being driven continuously from one shaft, and the motion of the other wheels being controlled by a carry mechanism.



registering

It now only remains to find a mechanical method of ~~determining~~ whether the multiple encipherment condition is fulfilled. This can be done most simply if we are willing to test each Stecker value of the central letter throughout all red starts before trying the next Stecker value. ~~In this case the machine would have to be set up so that it could register the results of the tests.~~

Suppose we are investigating the case where the Stecker value of the central letter is K. We let an current enter all of the chains of enigmes at their K input points, and at the K output points of the chains we put relays. The 'on' points of the five relays are put in series with a battery (say), and ~~another relay~~. A current flows through this last relay if and only if a current flows through all the other five relays, i.e. if the five multiple encipherments applied to K all give K. When this happens the effect is, essentially, to stop the machine, and such an occurrence is known at Letchworth as a 'set-right'. An alternative possibility is to have a quickly rotating 'scanner' which, during a revolution, would first connect the <sup>points A</sup> inputs of the chains to the current supply, and the output points A to the relays, and then would connect the input and output points B to the supply and relays. In a revolution of the scanner the output and input points A to Z would all have their turn, and the right hand wheel would then move on. This last possible evolution was called 'serial scanning' and led to all the possible forms of registration being known as different kinds of 'scanning'. The simple possibility that we first mentioned was called 'single line scanning'. Naturally there was much research into possible alternatives to these two kinds of scanning, which would enable all 26 possible Stecker values of the central letter to be tested simultaneously without any parts of the machine moving. Any device to do this was described as 'simultaneous scanning'.

The solution which was eventually found for this problem was more along mathematical than along electrical engineering lines, and would really not have been a solution of the problem as it was put <sup>to</sup> the electricians, ~~for~~ when we gave, as we thought, just the essentials of the problem. It turned out in the end that we had given them rather less than the essentials, and they therefore cannot be blamed for not having found the best solution. They did find a solution of the problem as it was put to them, which would probably have worked if they had had a few more months experimenting. As it was the mathematical solution was found before they had finished.

#### Pyre simultaneous scanning

The problem as given to the electricians was this. There are 52 contacts labelled  $A...Z, A', ..., Z'$ . At any moment each one of  $A, ..., Z$  is connected to one and only one of  $A', ..., Z'$ : the connections are changing all the time very quickly. For each letter of the alphabet there is a relay, and we want to arrange that the relay for the letter  $R$  will only close if contact  $R$  is connected to contact  $R'$ .

#### XXXXXXXXXXXXXXXXXXXX

The latest solution proposed for this problem depended on having current at 26 equidistant phases corresponding to the 26 different letters. There is also a thyatron valve\* for each letter. The filaments of the thyatrons are given potentials corresponding to their letters, and the grids are connected to the corresponding points  $A', ..., Z'$ . The points  $A, ..., Z$  are also

---

\*A thyatron valve has the property that no current flows in the anode circuit until the grid potential ~~XXXXXXXXXXXX~~ becomes more negative than a certain critical amount, after which the current continues to flow, regardless of the grid potential, until the anode <sup>potential</sup> ~~XXXXXXXX~~ is switched off.

given potentials with the phase of the letter concerned. The result is that the difference of potential of the filament and of thyrettron A the grid oscillates with an amplitude of at least  $2\pi \sqrt{L/C}$ ,  $E$  being the amplitude of the original supply ~~xxxxxx~~, unless A and A' are connected through the chain, in which case the potentials remain the same or differ only by whatever grid bias has been put into the grid circuit. The thyrettrons are so adjusted that an oscillation of amplitude  $2\pi \sqrt{L/C}$  will bring the potential of the grid to the critical value and the valve will 'fire'. The valve is coupled with a relay which only trips if the thyrettron fails to fire. This relay is actually a 'differential relay', with two sets of windings, one carrying a constant current and the other carrying the current from the anode circuit of the thyrettron. Fig 60 shows a possible form of circuit. It is ~~xxxxxxxxxx~~ probably not the exact form of circuit used in the Pye experiments, but is given to illustrate the theoretical possibility.

### The Spider

We can look at the Bombe in a slightly different way as a machine for making deductions about Stecker when the read start is assumed. Suppose ~~for~~ we were to put lamp-boards in between the enigmas of the machine, and label the lamp-boards with the appropriate letters off figure . For example in chain C the lampboard between C1 and C2 would be labelled A. If we were using ~~the machine with~~ <sup>one</sup> <sup>the</sup> key-board ~~we~~ could be labelled with the 'central letter'. Now when we depress a letter of the key-board we can read off from the lamp-boards some of the ~~xxxxxxxxxx~~ Stecker consequences of the hypothesis that the depressed letter is steckered to the central letter; ~~xxx~~ for one such consequence could be read off each lampboard, namely that the letter lighting is steckered to the name of the lamp-board.



In Fig 61 at the top are the chains, with the positions <sup>are</sup> and the letters of the chain. In each column ~~is~~ written some of the letters ~~in~~ which ~~may~~ can be inferred to be Stecker values of the letters at the heads of their columns from the hypothesis that X is a Stecker value of the central letter E. By no means all possible inferences of this kind are made in the figure, but among those that are made are all possible Stecker values for E except the right one, I. If we had taken a rod start that was wrong we should almost certainly have found that all of the Stecker values of E could be deduced from any one of them, and this will hold for any cribs with two or more chains. Remembering now that with our ~~first~~ Bombe one Stecker is deducible from another if the corresponding points on the lamp boards are connected through the enigmes, a correct rod start can only be one for which not all the input points of the chains are connected together; the positions at which this happens are almost exactly those at which ~~the~~ a Bombe with simultaneous scanning would have stopped.

#### ~~When several chains are used with the Bombe~~

This is roughly the idea of the 'spider'. It has been described in this section as a way of getting simultaneous scanning on the Bombe, and has been made to look as much like the Bombe as possible. In the next section another description of the spider is given.

#### The spider. A second description. Actual form.

In our original description of the Bombe we thought of it as a method of looking for characteristics of a crib which are independent of Stecker, but in the last section we thought of it more as a machine for making Stecker deductions. This last way of looking at it has obviously great possibilities, and so we will start afresh with this idea.

In the last section various points of the circuit were regarded as ~~corresponding~~ having certain Stecker corresponding to them. ~~Now~~ We are now going to carry this idea further and



13

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Stecker, e.g. if one knows that the letters which were ~~xxxxxxxx~~  
~~xxxxxxxx~~ unstecked on one day are invariably stecked on the  
next then, having ~~xxxxxx~~ found the keys for one days traffic one  
could when looking for the keys for the next day, connect together  
all points of the diagonal board which correspond to non-  
steckers which had occurred on the previous day. This would of  
course not entirely eliminate the inadmissible solutions, but  
would enormously reduce their number, the only solutions which  
would not be eliminated being those which were inadmissible on  
every ~~many~~ count.

One difference ~~xxxxxxxxxxxxxxxx~~ between this arrangement  
and the Bombe, or the spider as we described it in the last section  
is that we need only one enigma for each constetation.

Our machine is still not complete, as we have not put in any  
mechanism for distinguishing correct from incorrect positions. In  
the case of a crib giving a picture like Fig 59 where most of  
the letters are connected together into one 'web' it is sufficient  
at some point on  
to let current into the diagonal board ~~xx~~ some line with named  
after a letter on the main web, e.g. at the Ee point in the case of  
the crib we have been considering. In this case the only possible  
positions will be ones in which the current fails to reach all  
the other points of the E line of the diagonal board. We can  
detect whether this happens by connecting the points of the E line  
through differential relays to the other pole of our current  
parallel with one another and in series with the stop mech  
supply, and putting the 'on' points of the relays in ~~xxxxxx~~. Normally  
current will flow through all the differential relays, and they  
will not move. When one reaches a position which might be correct  
the current fails to reach one of these relays, and the current  
permanently flowing in the other wiring coil of the relay causes  
it to close, and bring the stopping mechanism into play. ~~xxxx~~  
~~xxxxxxxxxxxxxxxx~~ with Most likely what will happen is that there will  
be just one relay which closes, and this will be one connected to  
a point of the diagonal board which corresponds to a Stecker  
which is possibly correct: more accurately, if this Stecker is  
not correct the position is not correct. Another possibility is

that all relays close except the one connected to the point at which the current enters the diagonal board, and this point ~~is then~~ corresponds to the only possible Stecker. In cases where there is rather scanty, and the stops therefore very frequent, other things may happen, e.g. we might find four relays closing simultaneously, all of them connected together through the enigme and the cross connections of the diagonal board, and therefore none of them corresponding to possible Stecker.

zuckner  
The national affairs last night were

In order for it to be possible to make the necessary connections between the enigmas, the diagonal board and the relays there has to be a good deal of additional gear. The input and output rows of the enigmas are brought to rows of 26 contacts called 'female jacks'. The rows of the diagonal board are also brought. The 26 relays and the current supply are also brought to a jack. to female jacks. Any two female jacks can be connected with 'pleated jacks' consisting of 26 wires pleated together and ending in male jacks which can be plugged into the female jacks. In order to make it possible to connect ~~the~~ three or more rows of contacts together one is also provided with ~~an~~ 'commons' consisting of four ~~xxx~~ female jacks with corresponding points connected together. There is also a device for connecting together the ~~input~~ output jack of one ~~enigma~~ ~~xxx~~ and the input of the next, both being connected to another female jack, which can be used for connecting them ~~anywhere~~ ~~xxx~~ to anywhere else one wishes.

On the first spider made there were 30 engines, and three diagonal boards and 'inputs' i.e., sets of relays and stepping devices. There were also 18 sets of cams.





5int, D	5int, E	5int, F	5int, A	5int, B	5int, C
Input jack Fa	Fb	Fc	Fd	Fe	Ff
5in, A	5in, B	5in, C	5in, D	5in, E	5in, F

Engine 5

5int, A	5int, B	5int, C	5int, D	5int, E	5int, F
Output jack Ca	Cb	Cc	Cd	Ce	Cf
5in, D	5in, E	5in, F	5in, A	5in, B	5in, C

1int, F	1int, D	1int, E	1int, B	1int, C	1int, A
Input jack Fa	Fb	Fc	Fd	Fe	Ff
1in, A	1in, B	1in, C	1in, D	1in, E	1in, F

(missing from A)

Engine 1

1int, A	1int, B	1int, C	1int, D	1int, E	1int, F
Output jack Aa	Ab	Ac	Ad	Ae	Af
1in, F	1in, D	1in, E	1in, B	1in, C	1in, A

2int, B	2int, A	2int, D	2int, C	2int, F	2int, E
Input jack Aa	Ab	Ac	Ad	Ae	Af
2in, A	2in, B	2in, C	2in, D	2in, E	2in, F

Engine 2

(missing from C)

2int, A	2int, B	2int, C	2int, D	2int, E	2int, F
Output jack Ca	Cb	Cc	Cd	Ce	Cf
2in, B	2in, A	2in, D	2in, C	2in, F	2in, E

3int, F	3int, D	3int, E	3int, B	3int, C	3int, A
Input jack Ca	Cb	Cc	Cd	Ce	Cf
3in, A	3in, B	3in, C	3in, D	3in, E	3in, F

Engine 3

(missing from E)

3int, A	3int, B	3int, C	3int, D	3int, E	3int, F
Output jack Ea	Eb	Ec	Ed	Ee	Ef
3in, F	3in, D	3in, E	3in, B	3in, C	3in, A

4int, C	4int, E	4int, A	4int, F	4int, B	4int, D
Input jack Ea	Eb	Ec	Ed	Ee	Ef
4in, A	4in, B	4in, C	4in, D	4in, E	4in, F

(missing from F)

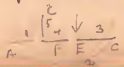
Engine 4

4int, A	4int, B	4int, C	4int, D	4int, E	4int, F
Output jack Fa	Fb	Fc	Fd	Fe	Ff
4in, C	4in, E	4in, A	4in, F	4in, B	4in, D

Fig 63 Spider connection with engine for 6 letters

	1	2	3	4	5
alphabet and into	A C E F F	A B	A A	A C	B C
	F A C E C	B D	C D	B D	B E
		B C	E A	C E	D A

Names of contacts are given in purple ink, contacts to which they are connected in green. Connection of diagonal bond to engine Fig 64.



Aa, A 1st, A 2nd, A	Aa Ba 1st, B 2nd, B	Aa Ca 1st, C 2nd, C	Aa Da 1st, D 2nd, D	Aa Ea 1st, E 2nd, E	Aa Fa 1st, F 2nd, F
Ba, Ab	Bb	Ba Cb	Ba Db	Ba Eb	Ba Fb
Ca, Ac 2nd, A 3rd, A 5th, A	Ca Bc 2nd, B 3rd, B 5th, B	Ca Cc 2nd, C 3rd, C 5th, C	Ca Dc 2nd, D 3rd, D 5th, D	Ca Ec 2nd, E 3rd, E 5th, E	Ca Fc 2nd, F 3rd, F 5th, F
Da, Ad	Da Bd	Da Cd	Da	Da Ed	Da Fd
Ea, Ae 3rd, A 4th, A Input A, output	Ea Be 3rd, B 4th, B Input B, output	Ea Ce 3rd, C 4th, C Input C	Ea De 3rd, D 4th, D Input D	Ea Ee 3rd, E 4th, E Input E	Ea Fe 3rd, F 4th, F Input F
Fa, Af 4th, A 1st, A 5th, A	Fa Bf 4th, B 1st, B 5th, B	Fa Cf 4th, C 1st, C 5th, C	Fa Df 4th, D 1st, D 5th, D	Fa Ef 4th, E 1st, E 5th, E	Fa Ff 4th, F 1st, F 5th, F

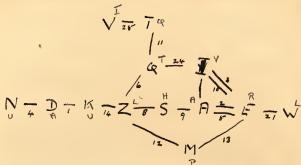
Fig 63. Connection of diagonal board. See fig 63.

'Input' is at E. Great hypothesis E/A. The squares in this figure represent contacts. As in fig. The purple letters are names and roman letters show the contacts to which they are connected.

entering at a correct Stecker if the position is correct. Let us further simplify the problem by supposing that there is only one 'web', i.e. that the 'picture' formed from the part of the crib that is being used forms one connected piece, e.g. with the crib on p we should have one web if we omit the constetations B, U, O, H. <sup>P B U R</sup> Clearly ~~sufficient condition for~~ a step is that the 'multiple encipherment' conditions should hold. Supposing that the number of independent chains or 'closures' is  $c$  then the number of positions where the multiple encipherment conditions hold will be about  $26^{4-c}$ . Some of the constetations of the web could still be omitted without any of the letters becoming disconnected from the rest. Let us choose some set of such constetations, ~~which cannot be omitted without breaking the web~~ in such a way that we cannot omit any more constetations without the web breaking up. When the constetations are omitted there will of course be no 'chains' or 'closures'. This set of constetations may be called the 'chain-closing constetations', and the others will be called the 'web-forming constetations'. At any position we may imagine that the web-forming constetations are brought into play first, and only if the position is possible for these are the chain-closing constetations used. Now the Stecker value of the input letter and the web-forming constetations will completely determine the Stecker values of the letters occurring in the web. When the chain closing constetations are brought in ~~it~~ it will already be completely determined what are the corresponding 'unstackered' constetations, so that if there are  $c$  chain-closing constetations the final number of steps will be a proportion  $26^{-c}$  of the steps which occur if they are omitted. Our problem reduced therefore to the case in which there are no closures. It is, I hope, also fairly clear that the number of steps will



To see what kind of contradictions are detected by the machine we can take the picture, Fig ~~67~~<sup>68</sup>, and on it write against each letter any Stecker values of that letter which can be deduced from the Stecker hypothesis which is read off the spider when it stops. This has been done in Fig ~~69~~<sup>70</sup> for a case where the input was on letter E of the diagonal board, and the relay R closed when the machine stopped; if the position of the stop were correct at all the correct Stecker would be given by the points of the diagonal board which were connected to Er, and they will be the direct consequences of the Stecker hypothesis E/R. Ifxxxxxxxxxxxxxxx  
xxxxxxdaxxixxxxxxxxxxxxxxxxSteckxxxxxxxxxxxxxxat  
xxxxxxxxxxxxxxlixxxxxxxxxxxxxxxlikeFigxxxxxxxxxxxxx  
xxxxxxxtaxdaxxxlixkexkxxxxxxserxxxxxxx As we are assuming that R was the only relay to close xxxxxxxxx, this relay cannot have been connected to any xix of the others, or it would have behaved similarly. We cannot therefore deduce any other Stecker value for E than R, and this explains why on the 'main web' in Fig ~~69~~<sup>70</sup> there is only one pencil letter against each ink letter. Wherever any pencil letter is the same as an ink letter we are able to xxxx write down another pencil letter corresponding tothe reciprocal Stecker or to the diagonal connections of the board. In one or two cases we find that the letter we might write down is there already. In others the new letter is written against xxx xx a letter of one of the minor webs; in such a case we clearly have a contradiction, but as it does not result in a second set of pencil letters on the main web the machine is not prevented from stopping. There are other contradictions; e.g. we have Z/L,W/L, but as L does not occur inthe crib this has no effect.



$\begin{matrix} P \\ D \\ D \\ P \end{matrix} \xrightarrow{20} \begin{matrix} B \\ L \\ R \\ Q \end{matrix} \xrightarrow{18} \begin{matrix} U \\ C \\ K \end{matrix} \xrightarrow{7} \begin{matrix} O \\ X \\ Y \\ B \end{matrix}$   
 $R \xrightarrow{11} H$   
 $\begin{matrix} B \\ L \end{matrix}$

### Relevant parts of alphabets

1 AU	2 AR	3 RV	4 ND	5 AR	6 QL	7 KN UV UB	8 LH	9 AH
10 RV	11 QT	12 LP	13 NK	14 LU	15 CK RN OU	16	17	18
19 DC FR ND	20 RL	21 BL	22 AV	23 VT	24 QI	25	26	

Fig 65, illustrating the kind of position at which the Spoker will stop. Here the input letter may be supposed to be Q and the relay which reads R. The sticker values of the letters, which are consequences of the hypothesis Q/R are written against the letters. There are unbroken lines such as Z/L, W/L: P/D, P/R, P/M which are not explained by the Spoker.

A blank 26x26 grid paper. The top edge is numbered 1 through 26, and the left edge is numbered 1 through 26. The grid is composed of small squares, with a slightly larger square at the top-left corner (row 1, column 1).



The machine gun

When using the spider there is a great deal of work in taking down data about stops from the machine and in testing these out afterwards, making it hardly feasible to run cribs which ~~run~~ give more than 5 stops per wheel order. ~~if it were possible to make the machine do the testing itself.~~ As the complete data about the direct consequences ~~of any~~ of any Stecker hypothesis at any position are already contained in the connections of the points of the diagonal board it seems ~~that~~ that it should be possible to make the machine do the testing itself. It would not be necessary to improve on the stopping arrangement of the spider itself, as one could ~~use the spider as already described,~~ use the spider ~~as already described,~~ as already described, and have an arrangement by which, whenever it stopped a further mechanism is brought in to play, which looks more closely into the Stecker. Such a mechanism will be described as a machine gun, regardless of what its construction may be.

With almost any crib the <sup>spider stops</sup> proportion of ~~machine gun~~ ~~machine gun~~ that would be passed by ~~the~~ machine gun as possible would be <sup>ratio</sup> higher than the ~~exercise~~ of ~~total~~ spider stops to total possible hypotheses. Consequently the amount of time that can economically be allowed to the machine gun for examining a position is vastly greater than can be allowed to the spider. We might for instance run a crib which gives 100 spider stops per wheel order, and the time for running, apart from time spent during stops might be 25 minutes. If the machine gun were allowed 5 seconds per position, as compared with the spider's 1/10 second only 8 minutes would be added to the time for the run.



which are wired up in such a way that we can distinguish  
 whether or not two or more of them are energized. When we are  
 testing the Stecker values of A we have the 26 contacts of the  
 A line of the diagonal board connected to the corresponding  
 relays in this set. ~~xxxxxxxxxxxx~~ What is principally  
 lacking is some device for connecting the rows of the diagonal  
 board successively to the set of relays. This fortunately was  
 found in post-office standard equipment; ~~and~~ the clicking noise  
 that this gadget makes when in operation gives the whole  
 apparatus its name. If we ~~xxxxxxxx~~ find no contradictions in  
 the Steckers of any letter the whole position is deemed as good.  
 The machine is designed to print the position and the Stecker  
 in such a case. Here again I do not know the exact method used,  
 but the following simple arrangement seems to give much the  
 same effect, although perhaps it could not be made to work  
 quite fast enough. The Steckers are given by ~~xxxxxxxx~~  
 typing one letter in a column headed by the other. ~~xxxxxx~~  
 When any  
 letter is tested for Stecker contradictions the relays  
 corresponding to the Stecker values of the letter close. We can  
 arrange that these relays operate keys of the typewriter, but  
 that in the case that there is a contradiction this is prevented  
 special typed instead  
 and some symbol is ~~xxxxxxxx~~ showing that the whole is wrong. When  
 no relay closes nothing is typed. The ~~xxxx~~ of the typewriter ~~xxxx~~  
 not operated by these keys but only by the space bar, and this is  
 there is a change of  
 moved whenever the letter ~~xxxxxxxxxxxx~~ whose Stecker are  
 being examined ~~xxxxxxx~~.

Additional gadgets

Besides the spider and machine-gun a number of other improvements ~~xxxxxxxxxxxx~~ of the Bombe are now being planned. We have already mentioned that it is possible to use additional data about Stecker by connecting up points of the diagonal board. It is planned to make this more straightforward by leading the points of the diagonal board to 325 points of a plug board; the plug board also has a great many points all connected together, and any Stecker which one believes to be false one simply connects to this set.

Another gadget is designed to deal with ~~xxxxxxxxxxxxxxxxxxxx~~ ~~xxxx~~ cases such as that in which there are two 'webs' with six and no chains letters on each. A ~~xx~~ little experiment will show that in the great majority of cases with such data, when the solution is found, the Stecker value of a letter on either web will imply the whole set of Steckers for the letters of both webs: in ~~as~~ the current terminology, "In the right place we can nearly always get from one web onto the other". If however we try to run such data on the spider, even with the machine gun attachment, there will be an enormous number of stops, and the vast majority of these will be cases in which "we have not got onto the second web". If we are prepared to reject these possibilities without testing them we shall not very greatly decrease the probability of our finding the right solution, but very greatly reduce the amount of testing to be done. If in addition the spider can be persuaded not to stop in these positions, the spider time saved will be enormous. Some arrangement of this ~~xxx~~ kind is being made but I will not attempt to describe how it works.

~~xx~~

With some of the ciphers there ~~xxxxxx~~ is information about the Ringstellung (Meriveliennue) which makes certain stopping

pieces wrong in virtue of their position, and not of the alphabets produced at these positions. There is an arrangement, known as a 'Ringstellung cut-out' which will prevent the machine from stopping in such positions. The design of such a cut-out clearly presents no difficulties of principle.

There are also plans for "majority vote" gadgets which will enable one to make use of data which is not very reliable. A hypothesis will only be regarded as rejected if it contradicts three (say) of the unreliable pieces of data. This method may be applied to the case of unreliable data about Stecker.

Handwritten notes at top right corner.

Handwritten text block A, containing various letters and symbols arranged in rows.

Handwritten letter 'B' in red ink.

Handwritten text block B, containing various letters and symbols arranged in rows.

Handwritten letter 'C' in red ink.

Handwritten text block C, containing various letters and symbols arranged in rows.

Handwritten text block D, containing various letters and symbols arranged in rows.

Handwritten letter 'E' in red ink.

Handwritten text block E, containing various letters and symbols arranged in rows.

Handwritten letter 'F' in red ink.

Handwritten text block F, containing various letters and symbols arranged in rows.

Handwritten letter 'G' in red ink.

Handwritten notes at the top of the page, including circled numbers 2, 5, and 10, and various letters and symbols.

Handwritten notes in the first section, including the word "H" and various letters and symbols.

Handwritten notes in the second section, including the word "H" and various letters and symbols.

Handwritten notes in the third section, including the word "H" and various letters and symbols.

Handwritten notes in the fourth section, including the word "H" and various letters and symbols.

Handwritten notes in the fifth section, including the word "H" and various letters and symbols.





Handwritten notes at the top of the page, including a circled 'O' and various letters and numbers.

Section marked with a red 'T' on the left. Contains handwritten text and numbers, including a circled 'U'.

Section marked with a red 'X' on the left. Contains handwritten text and numbers, including circled 'R' and 'K'.

Section marked with a red '3' on the left. Contains handwritten text and numbers, including a circled 'P'.

Section marked with a red 'V' on the left. Contains handwritten text and numbers, including a circled '1'.

Section marked with a red 'C' on the left. Contains handwritten text and numbers, including a circled 'E'.

Section marked with a red 'T' on the left. Contains handwritten text and numbers, including a circled 'L'.

>  
 AB    A    KEB    FKA    WMA.  
 CD  
 EF    REB    FKB    WMB.

AOB  
 BC  
 CD  
 DE  
 EF  
 F  
 G  
 H  
 I  
 J  
 K

if alphabet REA is AD

then BKE will be postalphabet REB

But the word line if however has more between A & B

suppose AID

+  
 3 = 100%

19/5 1736/92 • 44

10-9

ST 90 95 100 105 110 115 120 125 130 135 140 145 150 155 160 165 170 175 180 185  
 A R E O N C F E E S U R J L J I U G L T F N E O N E E U B M I O R T N I H L I S C U J B B L A X 1 A X 2 g M S B S S J N V I T O D S g M T E D K E D V R E P

## Chapter VII. The German Naval Enigma Cipher

### Historical

In the ~~period~~<sup>much</sup> from about 1931 to April 30, 1937 the Naval cipher ~~machine~~ used the same indicating system as the other German service enigma ciphers, viz. the 'boxing' method recommended by the firm that sold the 'commercial' enigma. With this ~~in~~ system as well as the set up of the machine consisting of wheel order, Ringstellung, and Stecker, there was a window position fixed for the day, and known as the 'Grundstellung'. When it was desired to encipher a message from a list of about 1700 trigrams e.g. ZLE one first chose three letters at ~~random~~<sup>random</sup>. One then set the machine to the Grundstellung and enciphered ZLEZLE. The resulting six letters were put at the beginning of the message, and the remainder of the message consisted of the result of enciphering the plaintext with pre-start window position ZLE. (This differs from the other boxing indicating systems in that ~~the~~ most of these allow the trigrams such as ZLE to be chosen at random instead of from ~~their~~ a restricted list.

The weakness of this indicating system is that a great deal of information is given away about the 'Grundstellung'. If and a known diagonal there were no Stecker, and the traffic amounted to 100 messages per diem it would be possible to find the connections of the machine, and if there were Stecker but the connections of the machine were known it would be possible to find the keys every ~~from~~ day with the same amount of traffic. ~~Suppose for example that~~  
To explain the possibility of finding the keys let us suppose that the following were a set of indicators for one day's traffic:

UJOOEL	AFIJVI	TIQOHL	RSICAI	JNZSUG
VYIYIM	MIDRWZ	INVZUV	APTJNA	UANOER
ALAJMB	KMPPFI	CZNYOR	GJWLBS	HALUXK
XDVBKV	MUZLIG	RLFCMN	CMUYSQ	UOG OHT
QLYAMM	IIFQWN	LMZKFC	IBWQIV	ZFUMVQ
GRYLMN	KHAPRB	MIDWNU	OOFVCN	JWSBAS
JIPBWN	SMLEPK	AULJJJ	RZHCOC	ZDIXKI
YELFTI	RPKCVJ	UWIOGI	AVRJJK	KPBPMY
EMARFB	LKZERP	FXRDEK	GILNWI	FTGD-T
TENGOR	BANRER	EYYKPM	UAMODC	HEKUTP
AYIJPI	GKJLEP	ZCDWLU	SOCKLE	UHSOIS
UWNOBR	HKCITY	KZSKOH	DLYNOM	
OQBVCI	TONGER	VZPRON	EYKCPV	
EMIDUJ	OIKWNP			

In the two indicators UJOOBL and UAKODR the repetition of the first letters is followed by a repetition of the fourth letters. That this must always happen is clear from the fact that the fourth letter arises from the first by enciphering at the position directly after the Grundstellung and re-enciphering three places further on. ~~xxxxxxx~~  
~~xxxxxxxx~~ This phenomenon enables us to tell very quickly with any cipher whether the boxing form of indication is being used. From the indicators we can find the effects of the three repeated encipherments. In Fig 46 we have entered ~~xxxxxxxxxxxx~~ in one of the columns against each letter the effect of enciphering it first at the position immediately after the Grundstellung and then at the position four places after the Grundstellung: thus we have the entry J against A, with five dots. This means that A enciphered at the first and fourth positions gives J, and that this information has been given us from six indicators, which are actually ALAJMB,AYIYPI,AFIYVI,APTINA,AVRJEK,AXUJJJ. The other two columns give us the results of the encipherments at the second and then the fifth position, and at the third and then the sixth. We get for the results of these double encipherments

$$G_4G_1$$

...KIDAJKXP...TGLEW...{DSEBUVEKBRCTT}

$$G_5G_2$$

{PNUJBQCIAFPVKY}{CHDSANKESTORZ}

$$G_6G_3$$

(V)(I){JFMREKFWSHOLX}...DUQEZOTABYMC...

$G_4G_1$  here means the encipherment with the first alphabet and then with the fourth, the reversal of the natural order being in agreement with mathematical tradition. There can be no doubt as to how the substitution  $G_6G_5$  is to be completed, but at first sight it might appear that there are two possibilities for  $G_4G_1$ . However if we remember what we found out in the section 'alphabets and boxes' we see that it must be possible to pair off the cycles of  $G_4G_1$  into ones of equal length.

There are various things which could be done now. Of course one might put the whole data onto the epider, but at the time that this system was in force no such machine had been thought of. Another method, which was that principally used by the Poles is to have a permanent catalogue of the box shapes ~~xxxx~~ for  $G_4G_1$ ,  $G_5G_2$ ,  $G_6G_3$  for every Grundstellung, assuming that there is not a T.O. between the first and last of the six alphabets. ~~xxxxxxx xxxxxxxxxx xxxxxxxxxx xxxxxxxxxx xxxxxxxxxx~~  
or numbers  
If we give some standard order to the box shapes, we can also put the possible series of three box shapes into an order, and can enter against each set of three box shapes the Grundstellungen for which this set is realized. To use this catalogue with our problem we should work out the box shapes viz.  $G_4G_1$  is 26,  $G_5G_2$  is 26,  $G_6G_3$  is 24,2. These box shapes actually ~~xxxxxxx~~ ~~xxxxxxx~~ 26 actually has the number 1 and 24,2 the number 2: they are the two commonest shapes as can be seen from the table p 194. We then look up 1.1.2 in our catalogue, and find about 120 entries against it for each wheel order. Each of these will have to be tested out in some way or other. The most satisfactory method seems to be this: We form the permutation ~~xxxx~~ ~~xxxx~~  $G_5G_3G_5G_2$ . It is (PRGKLMQDJYWHISBEAUFV)(GXZL)(O)(T) so that this permutation is of the class 20,4,1,1. For each possible Grundstellung it is possible to calculate the corresponding class for the unscrambled alphabets. This can fortunately be done mechanically by means of a form of 'cyclo-meter'. It would be as well to enter against each position the class of this permutation, and this might have been done at the time of construction of the catalogue. ~~xxxxxxx~~  
In the case in question the right Grundstellung is found to have the position 1,1,26 with wheel order III I,II,III (service machine, Unkelrweitz A). The corresponding boxes are

DT PE CL  
 NO HN TU  
 HL UG XU  
 UP JL JD  
 CK BE FO  
 AI QX RM  
 EJ OD NY  
 XV TI KB  
 BQ MI PV  
 NS FW WL  
 OR AS IC  
 YW KH HZ  
FL YC AB

We ~~extensively~~ must have V/A or V/S. If V/A we can identify the cycle (CHWSADKRETBZ) of the G<sub>5</sub>G<sub>2</sub> with stacker, with the half compartment of the second box in this way

(CHWSADKRETBZ)  
 (CHWIVCKLONZ)

i.e. we ~~xxxxx~~ have to assume the stacker O/C,I/S,A/V,T/L,R/N, and that HYN,D,X,K,G,Z are unstackered. This large number of unstackered letters is a strong confirmation, and the repetition of the Stacker I/S is further confirmation. When we fit the pos of the boxes together we find that these five are all the stacker.

There are other methods that can be applied, depending on the number of Stacker being small. The number of Stacker used in the Naval was 6 from 1931 to Nov. 1938 and possibly later. We might for instance have assumed that A and S were both unstackered and therefore assumed that the constetation S occurred in both the alphabets G<sub>3</sub> and G<sub>5</sub>. With the Turing sheets we could find the possible positions for this, and then use a cyclometer to test ~~xxxx~~ the box shapes in those positions. This is naturally only worth while if we have no box-shape catalogue. Another possibility is to 'equates' the boxes, i.e. to find out from the permutations G<sub>4</sub>G<sub>1</sub> etc what the original ~~xxxxxxxx~~ alphabets G<sub>1</sub> and G<sub>4</sub> were. In our case there are actually 13 different possibilities for G<sub>1</sub>, 13 for G<sub>2</sub> and 12 for G<sub>3</sub>. There are two things we can do to distinguish between the correct and the incorrect possibilities. We can use known statistics about the list of admissible message settings, choosing that combination of alphabets that gives the ~~xxxx~~ greatest number of ~~message~~ settings that have ~~xx~~ occurred ~~pre~~.

repetitions between the message settings for the day in question and message settings of previously solved days. We might also do a 'Benburismue' i.e. we might make use of the fact that if two messages are written out with letters that were enciphered at the same position written in the same column, then the number of ~~xxxx~~ repetitions of letters in a column will be the same as if the messages had not been enciphered, and therefore will be <sup>on average</sup> greater than if the messages had been otherwise placed. Actually this effect was very small for the Naval traffic in 1937 and earlier. The repetition frequency was 1/20, as compared with 1/16.5 for the 1940 Naval traffic and the Air traffic, ~~was~~ 1/12 for plain language German, and 1/26 for incorrectly placed enciphered messages (the repetition frequency is the ratio of the number of identical pairs in a column to the number of pairs in columns, identical or not). With so low a repetition frequency it is ~~xxxxxxxxxxxx~~ extremely difficult to equate the boxes unless the traffic is rather heavy. This method however applies quite well with the Air traffic up to Sept 14 1938, but there there are better methods of equating. Once the boxes have been equated by one means or another we shall have many more cases of ~~xxx~~ half-bombes which we can assume to have been unsteckered. This method will nearly always get the result, if the equating can be done.

After we have found the row position of the Grundstellung and the Stecker it only remains to find the Ringstellung. Usually this would be known already, as, at this period, the wheel order and Ringstellung were only changed about once a fortnight. However if these have just been changed it is necessary to read one message. This could always be done, as a great many messages were sent in two or more parts. In such cases the call signs and signatures of the parts were essentially the same, and the ~~xxxx~~ parts after the first began by saying that they were continuations, giving the last part of the time group of the ~~previous~~ message as a reference. ~~Thisxxxxxxx~~ ~~xxxx~~ The method of giving numbers at that time was to use



the top row of the key board and P thus

QWERTZUIOP  
1234567890

The number was put between Y's to show that it was a number, and the whole repeated as a check. The continuation of a message whose time group was 2330 would ~~begin~~ begin FORTWEEFYWEEFY. We could then find the position where the message started by single wheel processes, and as we already know the window position of the start, we can calculate the Ringstellung.

On the 1st May 1937 a new indicating system was introduced. The first two groups (four letters each) of the message were repeated at the end. This clearly showed that these two groups formed the indicator. The repetition also showed that no check could be expected within the first two groups themselves. This was discouraging, as the essential weakness of the boxing method was that the something was enciphered twice with the machine. With the new method of indicating, whatever it is, the best one can hope is that either it will enable us to set the message, or that we from some information about the setting of the message obtained ~~from~~ elsewhere we may be able to deduce something about the ~~affair~~ machine setting. However the first thing to be done was to find out how the indicators worked, and if necessary therefore to try and read some messages with which the new system was being used. To do this one can use the FORTWEEFY messages, and apply one of the methods described at the beginning of the last chapter. In this way the Poles found the keys for the 8th of May 1937, and as they found that the wheel order and the turnovers were the same as for the end of April they rightly assumed that the wheel order and Ringstellung had remained the same during the end of April and the beginning of May. This made it easier for them to

find the keys for other days at the beginning of May and they actually found the Stecker for ~~about~~ the 2nd, 3rd, 4th, 5th and 8th, and read about 100 messages. ~~XXXXXXXXXXXXXXXXXXXX~~  
~~XXXX~~ The indicators and window positions of four (selected) messages for the 5th were

Indicator	Window start
KFXJ EWTW	P C V
SYLQ EWUF	B Z V
JMHO UVQG	M E M
JMFE FEVC	MY X

The repetition of the EW combined with the repetition of V suggests that the ~~third~~~~fourth~~ fifth and sixth letters describe the third letter of the window position, and similarly one is led to believe that the first two letters of the indicator represent the first letter of the window position, and that the third and fourth represent the second. ~~XXXXXXXXXXXX~~  
~~XXXX~~ Presumably this effect is somehow produced by means of a table of bigramme equivalents of letters, but it cannot be done simply by replacing the letters of the window position with one of their bigramme equivalents, and then putting in a dummy bigramme, for in this case the window position corresponding to JMFE FEVC would have to be say MYI instead of MYX. Probably some encipherment is involved somewhere. The two most natural alternatives are . i) The letters of the window position are replaced by some bigramme equivalents and then the whole enciphered at some 'Grundstellung', or ii) The window position is enciphered at the Grundstellung, and the resulting letters replaced by bigramme equivalents. The second of these alternatives was made far more probable by the following indicators occurring on the 2nd May

KEDP IVJO	V C P
XXXX JXJY	V U E
ROXX JXMA	N U M

With this second alternative we can ~~now~~ deduce from the

first two indicators that the bigrammes XX and XX have the same value, and this is confirmed from the second and third, where XX and XX occur in the second position instead of the first.

It so happened that the change of indicating system had not been very well made, and a certain torpedo boat, with the call sign AFA: had not been provided with the bigramme tables. This boat sent a message in another cipher explaining this on the 1st May, and it was arranged that traffic with AFA: was to take place according to the old system until May 4, when the bigramme tables would be supplied. Sufficient traffic passed on May 2,3 ~~for~~ to and from AFA: for the Grundstellung used to be found, the Stecker having already been found from the FORTYSEVEN messages. It was natural to assume that the Grundstellung used by AFA: was the Grundstellung to be used with the correct method of indication, and as soon as we noticed the two indicators mentioned above we tried this out and found it to be the case.

There actually turned out to be some more complications, <sup>at least</sup> There were two Grundstellungen instead of one. One of them was called the Allgemeine and the other the Offiziers Grundstellung. This made it extremely difficult to find either Grundstellung. The Poles pointed out another possibility, viz that the trigrammes were still probably not chosen at random. They suggested that probably the window positions enciphered at the Grundstellung, rather than the window positions themselves were taken off the restricted list.

In Nov. 1939 a prisoner told us that the German Navy had now given up writing numbers with Y...YY...Y and that ~~the~~ the digits of the numbers were spelt out in full. When we heard this we examined the messages toward the end of 1937 which were expected to be continuations and wrote the expected beginnings under them. The proportion of 'crashes' i.e. of letters apparently left unaltered by encipherment, ~~then~~ then shows how nearly correct our guesses were. Assuming that the change ~~was~~

mentioned by the prisoner had already taken place we found th at  
about 70% of these cribs must have been right. Further 'crash  
analyses' were made for other periods up to Aug 1939, all with  
fairly favourable results. At the same time there had been some  
chen gas in the machine, known to ave taken place because ofthe  
corresponding changes in the machine used by the army and air  
whose traffic h ad been read. In the summer of 1937 the Umkehrwalze  
had been changed from A to B, and in Dec/1938 two new wheels  
~~in IV~~<sup>x</sup> and V had been introduced. After  
the beginning of the war (Sept 1939) the FORTYSEVEN messages  
were no longer traceable, because there were no more call signs.  
However there had been some traffic at various times during  
manoeuvres and crises since the occupation of Austria, Therax  
~~xxxxxxxtxxxxxxxxxxxmighttexasixtofindxxxxx~~  
~~keyxxxxxxxxxxxxxx~~ and there were a few days where there  
was both traffic with and without call signs. We hoped that  
we might be able to find the keys for some such days and~~s~~ so  
to find the kind of thing that was said in the traffic without  
call signs. There seemed to be  
some doubt as to the feasibility of this plan, as itappeared  
the call sign traffic on any day was always either the whole of  
the Baltic traffic or the whole of the non-Baltic traffic, and  
the Baltic traf- fic in 1937 used to be on a different key from the  
rest. Following this programme we found the keys for Nov. 28  
1938 and for a number of days near there. The number of Stecker  
was 6. The wheel order and Ringstellung seemed to remain constant  
for about a week; at en y rate they did not change between Nov 28  
and Nov 29. The Stecker ~~were~~ were not batted; the  
same letter was never steckered on t o consecutive days. This  
of course might be extremely valuable. If the traffic had been  
heavier it would have enabled us to find the keys so long as  
this lasted, and there were many cribs. Actually we got no  
further than this, as at this point a good deal of date was.

'pinched' from a German boat, enabling us set the keys for April 22-27 1940. At the same time we pinched a book of instructions telling us the precise form of the indicating system.

To encipher a message the operator chooses two trigrammes out of a book. The first of these trigrammes is called the 'Schluesselkenngruppe'. The choice of this is partly determined by the nature of the message: e.g. all 'dummy' messages have the Schlueselkenngruppe taken from one part of the book and genuine messages have then taken from elsewhere: we do not know very much about this. The second trigramme is called the Verfahrenkenngruppe. Suppose the Schlueselkenngruppe is CIV and the Verfahrenkenngruppe is TOD then the operator chooses two dummy letters, Q and X say, and writes this down

Q O I V  
T O D X

From the Verfahrenkenngruppe is obtained the window position for the start of the message, by enciphering at the Grundstellung. From the eight letters above one also obtains the indicator for the message, by substitution from a table which gives bigrammes for bigrammes. The ~~matrix~~ substitution is done by replacing the vertical pairs above with bigrammes, e.g. ~~ix~~ in this case, if the substitute for QT were BA, and TH for CO, PO for ID, and CN for VX then the indicator for the message is DATH POCH. Apart from the Schlueselkenngruppe feature this is the method we had inferred was being used. ~~This extra~~ This extra feature accounts for the bigrammes in the indicators being almost perfectly hit. Also the fact that it is never the message setting itself which is chosen at random by the operator eliminates any remaining hope that one might use 'operator's psychology' to help in finding out the alphabets. From our point of view of course the Schlueselkenngruppen might as well not exist, and the 'bigramme lists' to us remain letter entered. Pages sheets with one ~~matrix~~ in each square, and not two. There is however the restriction that there must be exactly 26 occurrences of each letter.

\* This has gone out of date

### Methods of reading the individual messages

With the system of indication that has been used since May 1937 we are not able to read all the messages as soon as we have read one. A few may be read by single wheel processes, starting from a short crib, but we cannot hope to read the whole traffic in this way. Also, when we have found the Grundstellung ~~xxxxxxxxxxxx~~, and if there is plenty of traffic, we may be able to make use of ~~xxx~~ some big-rams which occurred in messages already read. These methods are not enough by themselves. In the 1937 traffic there was no 'not probable', and we had planned a method for finding the right starting position, making use of the fact that the correct decode would probably have more letters E in it than any of the others. It was intended to have a long punched paper roll, the punching showing the effect of enciphering E in the various positions. This paper was to move under a series of about 200 brushes whose position was determined by the letters of the enciphered message. The number of brushes which poked through the holes at any moment was the number of letters E in the decode of the message, the window position ~~xxxxx~~ being determined by the position of the roll. All positions giving more than a certain number of letters E were to be recorded and these positions intensively tested. This machine was called 'the rack'.

It was never necessary to make a rack because when the 1938 messages were read it was found that the word EINS occurred very frequently. We therefore made a catalogue of the encoded values of EINS at every possible starting position, and arranged the encoded values in alphabetical order. The unanalysed catalogue was made by enciphering first E at every possible position, then I, N and S. This was done with the automatic typewriting machine. The values of I were stuck below the values of E with a stagger; the values of N and S were underneath these again, with suitable staggers. The result was that the effect of enciphering EINS appeared in vertical columns.

This unanalysed catalogue was known to the girls as 'corsets'. In analysing the catalogue we took 25 sheets named A to Z, with Z omitted: each sheet had 25 lines, named A to Z with I omitted. Supposing on sheet L and line 4 of the corsets we found LVOM as a value of EINS we would enter <sup>OM</sup> 13.4 on line V of sheet L. ~~Furthermore~~ In a later form of the catalogue we also made 'existence sheets'. In the existence sheets we would enter M in line V and column O of sheet L. To use the catalogue we first analysed the tetragrammes in the messages ~~xxxxxx~~ according to their first letters. One would then take the existence sheet and go through all the tetragrammes marking the tetragrammes which occurred on the existence sheet, and marking against them the entry (e.g. 13.4) from the catalogue. Afterwards one would have to go back to the corsets, and search in the right line for the tetragramme, and work out its position: this was done with a cardboard strip and known as 'snaking'. Having found the position one would have to set up the machine, decipher the tetragramme, verifying that it gave EINS and then continue to decipher and see if one continued to get sense.

This process has since been greatly improved. Instead of making the corsets off the 'X-machines' we have a machine called the 'test-plate' or 'baby' which typed out the results of enciphering EINS in all positions in a much more convenient form. Also we ~~make~~ no longer analyse the groups by hand, but have <sup>together with their position</sup> ~~xxxxxxxxxxxxxxxxxxxxxxxx~~ them punched on cards, which are then sorted into alphabetical order, and listed. A further improvement is that the test-plate is now made to punch the cards directly.

Roughly, our programme when the wheel order, Ringstellung, and Stecker for a day have been found, is as follows. We make an EINS catalogue, and use <sup>it</sup> to get out pairs of messages in which the second indicator bigramme of one is the same as the third indicator bigramme of the other. If we have four such cases we have sufficient data about the Grundstellung to be able to find it by means of the Bombe, provided that we have

found the double T.O. ~~we~~ we then continue to get messages out with the EINS catalogue; each message gives us some ~~hi~~ values of bigrammes, which are entered on a Poes sheet. From time to time we go through the messages substituting for the bigrammes the values that have been found from the messages. With messages for which we know the values of two of the bigrammes we apply the method known as 'twiddling' or 'bonking'. We have to decipher the first few letters of the message at all of the 26 places consistent with our knowledge of the bigrammes. This is usually done in columns, one column at a time, each column corresponding to a letter of the message. The twiddling is best done on the Letchworth enigmas, as they have no automatic T.O. Some more messages can be solved ~~by~~ when one bigramme is known, preferably that corresponding to the L.H.W. on the test-plate by deciphering a few letters at every one of the 676 places. But this method is rather difficult to work in practice. It seems much more difficult to spot the right answer when one has to look through so many possibilities. The right answer is hardly ever noticed unless it is one of the obvious ones such as ~~BIDNEY~~ <sup>BRIDE</sup> WESPE, MUECKE, MOSKITO, HOENISSE, KREER, ANAN, ADAM, GRUPPE, ~~WESPE~~. The case where the R.H.W. bigramme is known cannot be done on the test-plate at all. One ~~may~~ can of course use the K-machine in much the same way as was done with the original form of EINS catalogue. This has never been a success. One can also use 'hand methods'. One can go through the message looking for places where two consecutive letters occur on the same rod. The deciphered values also occur on the same rod, and we can examine the rods for possible bigrammes. Combining this with the Turing sheet, Kendrick has solved quite a number of messages. This method is known as 'clicks on the rods'.

We now have the EINS catalogue collated with the messages of the day so that the only remaining work in EINS is the task of such possible EINS with machine, working out the



### Identification of bigramme lists and evaluation of unknown bigrammes.

The Vekfabrenkengruppe (V.K.G. or trigramme) is as we have explained not chosen at random, but from a list of  $n$  out 11,000, and within this list the choices are not made at random uniformly. This fact enables us to identify which bigramme lists are being used, for if we choose the right bigramme list and work out the V.K.G. we shall find that a comparatively large proportion of the  $n$  have occurred before, and if we choose the wrong one, a comparatively small proportion.

The more precise theory of this identification is as follows. Let us suppose that  $n$  the  $24^3$  different trigrammes  $\mu$ , have thus been used <sup>once</sup> before  $n$  times,  $\mu_1$  twice etc. Let us call a trigramme which has occurred before  $t$  times a "trigramme of the  $t$ -class". We can then express our information in the form:

Of the occurrences of trigrammes there have been  $\mu_1$  in the 1-class,  $2\mu_2$  in the 2-class,  $3\mu_3$  in the 3-class etc:

Now take a random sample of these occurrences, forming a proportion  $\alpha$  of the whole, and let us imagine that this random sample consists of the last of the trigrammes which were found. There will be  $\alpha\mu_1$  in the 1 class,  $2\alpha\mu_2$  in the 2 class, etc. Now the ones in the 1 class would have been, when they were found, ones which had not occurred before, and those which in the 2 class ones which had occurred before once, and so on. Hence we can say that for the last ~~trigrammes found~~ occurrences of trigrammes entered, the numbers which had occurred before <sup>once, twice, thrice, ...</sup> are in the ratios of  $\mu_1, 2\mu_2, 3\mu_3, \dots$ . We must expect these ratios to hold also of the next few occurrences to be entered. The process of finding new occurrences of trigrammes and looking up the numbers of previous occurrences can therefore be regarded as like having an urn containing cards, each of which bears a trigramme and a number, and making draws from the urn. The number of cards bearing the number  $t$  is to be proportional to  $(t+1)\mu_{t+1}$ . On the other hand we have to consider the process of choosing trigrammes at random. This is to be ~~regarded as~~ compared with drawing cards from an urn containing cards in different proportions,

This process worked well initially. The popular trigraphs were at the top of columns on the centre pages of the K-book, but the German instructions were to mark any trigraph as "it was used, and not to re-use it". Thus the repeat rate of the new trigraphs with those known to have been used gradually dropped.

The <sup>put it in</sup> K-book <sup>and for this trigraph adding</sup> contained <sup>all</sup> trigraphs, in <sup>the</sup> <sup>2nd</sup> <sup>1941</sup> <sup>order</sup>. I. J. Good, <sup>at the</sup> <sup>free choice</sup> <sup>dummy letter</sup> devised a quicker method, using the non-randomness of the German code. SELM 2/2/70

q. 14.263

in code

Each trigram must occur equally often in this urn, and must of course have with it the number of previous occurrences of this trigram. Now imagine that we have worked out a certain number of V.K.G. using a given bigramme table, and that we have found out how many times each of them had occurred before. This can be compared with being given one of the urns, and told it is Q:1 on this being the random urn, and then drawing a certain number of cards from the urn. After the draw we have a new idea of the odds that the urn is the random urn, and we should have a corresponding modified idea of the odds that the bigramme list is the right one. Let us suppose that the trigramme, in the order as they were found, had the numbers  $r_1, r_2, \dots, r_s$  of previous occurrences, and that correspondingly the cards drawn from the urn bore the numbers  $r_1, r_2, \dots, r_s$ . The proportion of cases of draws of  $s$  cards from the urn, giving these results with the same order, is  $u_{r_1}, u_{r_2}, \dots, u_{r_s}$  where  $u_r$  is the proportion of  $r$ -cards in the urn. Likewise the proportion of cases where this happens with the other urn is  $u'_{r_1}, \dots, u'_{r_s}$  with a corresponding meaning for  $u'_r$ . Then the odds on the urn not being the random one after the draw experiment are

$$\frac{u_{r_1}}{u'_{r_1}} \cdot \frac{u_{r_2}}{u'_{r_2}} \dots \frac{u_{r_s}}{u'_{r_s}} : Q$$

In other words the drawing of a card with the number  $x = m$  improves the odds by a factor of  $\frac{u_m}{u'_m}$ , which is equal to  $\frac{2^{s-1} (m+1) r_{m+1}}{(\sum_{n=1}^m (n+1) r_{n+1}) r_m}$  except in the case  $m=0$  when it is  $\frac{2^{s-1} r_1}{(\sum_{n=1}^m r_n) \delta \frac{1}{2} r_1}$

The same method may be applied for the identification of some unknown bigramme. By taking into account a number of days traffic all using the same bigramme table we may find a number of indicators whose V.K.G. would be completely known if we knew the value of a certain bigramme. If we make the right hypothesis as to the value, we should get trigramme agreeing with the statistics as before. In this sort of case, as the data is liable to be very scanty, it is essential to use the accurate theory as described above.